CS 352 – Compilers: Principles and Practice Mid-term Examination, 10/26/05

Instructions: Read carefully through the whole exam first and plan your time. Note the relative weight of each question and part (as a percentage of the score for the whole exam). The total points is 100 (your grade will be the percentage of your answers that are correct).

This exam is **closed book**, **closed notes**. You may *not* refer to any book or other materials.

You have **75 minutes** to complete both (2) questions. Write your answers on this paper (use both sides if necessary).

Name:

Student Number:

Signature

- 1. [Regular expressions, scanning; 40%]
 - (a) (15%) As described in class, construct an NFA that recognizes the same language as defined by the following regular expression:

$$(ab^*c) \mid (abc^*)$$

Answer:



(b) (25%) Using the subset construction, convert the NFA into a DFA. Optimize the resulting DFA by merging equivalent states (if any).Answer:



2. [Context Free Grammars, parsing; 60%] Consider the following simple grammar (with numbered productions) and the language it describes:

(a) (5%) Write a regular expression for this *language*.

Answer:

ab | ba

(b) (10%) Show the steps of a top-down parser as it parses the input *ab*, showing the input as it is consumed and the parse stack (left is top, right is bottom) at each step of the parse.

Answer:

| Stack | Input | |
|-------------|-------|--|
| <i>S</i> \$ | ab\$ | |
| AaAb\$ | ab\$ | |
| aAb\$ | ab\$ | |
| Ab\$ | b\$ | |
| b\$ | b\$ | |
| \$ | \$ | |

(c) (5%) Is this language LL(1)? Why or why not?

Answer:

Yes, the language is regular and all regular languages are LL(1).

(d) (10%) Construct the (top-down) LL(1) parse table for this grammar. Is the grammar LL(1)? Why or why not?

Answer:

| | FIRST | FOLLOW | а | b | \$ |
|---|-------|--------|--------------------------|--------------------------|----|
| S | ab | \$ | $S \rightarrow AaAb$ | $S \rightarrow BbBa$ | |
| Α | 3 | ab | $A \rightarrow \epsilon$ | $A \rightarrow \epsilon$ | |
| B | 8 | ab | $B \rightarrow \epsilon$ | $B \rightarrow \epsilon$ | |

(e) (10%) Show the steps of a bottom-up parser as it parses the input *ab*, showing the input as it is consumed, the parse stack (left is bottom, right is top) at each step of the parse, and the action applied at each step.

Answer:

| Stack | Input | Action |
|-------------|-------|--------|
| \$ | ab\$ | reduce |
| \$A | ab\$ | shift |
| \$Aa | b\$ | reduce |
| \$AaA | b\$ | shift |
| \$AaAb | \$ | reduce |
| \$ <i>S</i> | \$ | accept |

(f) (10%) Construct the (bottom-up) SLR(1) parse table for this grammar. Is the grammar SLR(1)? Why or why not?

Answer:

| | | 0 | al | α¢ | | | |
|--------|-----------------|-----|-------------------------|----------------|----|---|---|
| | | 0: | $S' \rightarrow$ | •5\$ | _ | | |
| | | | $S \rightarrow$ | •AaA | b | | |
| | | | $S \rightarrow$ | •BbE | Ba | | |
| | | | $A \rightarrow$ | • | | | |
| | | | $B \rightarrow$ | • | | | |
| | | 1: | $S \rightarrow$ | $A \bullet a$ | Ab | | |
| | | 2: | $S \rightarrow$ | $B \bullet b$ | Ba | | |
| | | 3: | $S \rightarrow$ | $Aa \bullet A$ | Ab | | |
| | | | $\tilde{A} \rightarrow$ | • | | | |
| | | 4: | $S \rightarrow$ | $Bb \bullet I$ | Ba | | |
| | | | $\tilde{B} \rightarrow$ | • | | | |
| | | 5: | $S \rightarrow$ | AaA | •b | | |
| | | 6: | $S \rightarrow$ | BbB • | •a | | |
| | | 7: | $S \rightarrow$ | AaAk | • | | |
| | | 8. | $\tilde{S} \rightarrow$ | RhRa | | | |
| | | g · | S'→ | 5.5 | | | |
| | |). | 5 / | ψΨQ | | | |
| | a | | b | \$ | S | A | В |
| 0 | r3/ | r4 | r3/r4 | | 9 | 1 | 2 |
| 1 | s3 [′] | | , | | | | |
| 2 | | | <i>s</i> 4 | | | | |
| 3 | <i>r</i> 3 | | r3 | | | 5 | |
| 4 | r4 | | r4 | | | | 6 |
| 5 | | | <i>s</i> 7 | | | | |
| 6 | <i>s</i> 8 | | | | | | |
| 7 | ~~ | | | r1 | | | |
| , 8 | | | | r^{\prime} | | | |
| 0 | | | | 12 | | | |
| 9 | | | | а | | | |

(g) (10%) Here is the LR(1) parse table for this grammar:

| | а | b | \$ | S | Α | B |
|---|------------|------------|------------|---|---|---|
| 0 | <i>r</i> 3 | r4 | | 9 | 1 | 2 |
| 1 | <i>s</i> 3 | | | | | |
| 2 | | <i>s</i> 4 | | | | |
| 3 | | <i>r</i> 3 | | | 5 | |
| 4 | r4 | | | | | 6 |
| 5 | | <i>s</i> 7 | | | | |
| 6 | <i>s</i> 8 | | | | | |
| 7 | | | r1 | | | |
| 8 | | | <i>r</i> 2 | | | |
| 9 | | | а | | | |

Show the steps of this LR parser as it uses the table to parse the input *ba*, showing the input as it is consumed, the state stack (left is bottom, right is top) at each step of the parse.

Answer:

| Stack | Input | Action |
|--------------------------|-------|------------|
| 0\$ | ba\$ | r4 |
| $0_{\$} 2_B$ | ba\$ | <i>s</i> 4 |
| $0_{\$} 2_B 4_b$ | a\$ | r4 |
| $0_{\$} 2_B 4_b 6_B$ | a\$ | <i>s</i> 8 |
| $0_{\$} 2_B 4_b 6_B 8_a$ | \$ | r2 |
| $0_{\$}9_{S}$ | \$ | а |