

Orthogonal Polynomials (in Matlab)

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- Recurrence coefficients

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- Modified Chebyshev algorithm

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- Discrete Stieltjes and Lanczos algorithm

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- Discretization methods

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Sobolev Orthogonal Polynomials

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Sobolev Orthogonal Polynomials

- Moment-based algorithm

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Sobolev Orthogonal Polynomials

- Moment-based algorithm
- Discretization algorithm
- Zeros

Reference W. Gautschi, "Orthogonal Polynomials: Computation and Approximation", Clarendon Press, Oxford, 2004

<http://www.cs.purdue.edu/homes/wxg/papers.html>

click on Madrid.ps

also in: Lecture Notes in Mathematics, 2007(?)

<http://www.cs.purdue.edu/archives/2002/wxg/codes>

Background

inner product

$$(p, q)_{\mathbf{d}\lambda} = \int_{\mathbb{R}} p(t)q(t)\mathbf{d}\lambda(t), \quad \mathbf{d}\lambda \geq 0$$

orthogonal polynomials

$$\pi_k(\cdot) = \pi_k(\cdot; \mathbf{d}\lambda) : \quad (\pi_k, \pi_\ell)_{\mathbf{d}\lambda} \begin{cases} = 0, & k \neq \ell \\ > 0, & k = \ell \end{cases}$$

three-term recurrence relation

$$\pi_{k+1}(t) = (t - \alpha_k)\pi_k(t) - \beta_k\pi_{k-1}(t), \quad k = 0, 1, \dots, n-1$$

$$\alpha_k = \alpha_k(\mathbf{d}\lambda) \in \mathbb{R}, \quad \beta_k = \beta_k(\mathbf{d}\lambda) > 0 \quad (\beta_0 = \int_{\mathbb{R}} \mathbf{d}\lambda(t))$$

Jacobi Matrix

of infinite order

$$\mathbf{J}(\mathbf{d}\lambda) = \begin{bmatrix} \alpha_0 & \sqrt{\beta_1} & & & 0 \\ \sqrt{\beta_1} & \alpha_1 & \sqrt{\beta_2} & & \\ & \sqrt{\beta_2} & \alpha_2 & \ddots & \\ & & \ddots & \ddots & \\ 0 & & & & \end{bmatrix}$$

of order n

$$\mathbf{J}_n(\mathbf{d}\lambda) = \mathbf{J}(\mathbf{d}\lambda)_{[1:n,1:n]}$$

"Classical" Weight Functions

$$d\lambda(t) = w(t)dt$$

name	$w(t)$	on
Jacobi	$(1-t)^\alpha(1+t)^\beta$	$[-1, 1]$
Laguerre	$t^\alpha e^{-t}$	$[0, \infty]$
Hermite	$ t ^{2\alpha} e^{-t^2}$	$[-\infty, \infty]$

Matlab

example: `ab=r_jacobi(N,a,b)`

α_0	β_0
α_1	β_1
\vdots	\vdots
α_{N-1}	β_{N-1}

ab

$$N \in \mathbb{N}, a > -1, b > -1$$

Demo #1

```
N=10;
```

```
ab=r_jacobi(N,-.5,1.5);
```

Demo #1

```
N=10;
```

```
ab=r_jacobi(N,-.5,1.5);
```

k	α_k	β_k
0	0.6666666666666667	4.71238898038469
1	0.1333333333333333	0.1388888888888889
2	0.05714285714286	0.2100000000000000
3	0.03174603174603	0.22959183673469
4	0.02020202020202	0.23765432098765
5	0.01398601398601	0.24173553719008
6	0.01025641025641	0.24408284023669
7	0.00784313725490	0.245555555555556
8	0.00619195046440	0.24653979238754
9	0.00501253132832	0.24722991689751

Modified Chebyshev Algorithm

Modified Chebyshev Algorithm

modified moments

$$m_k = \int_{\mathbb{R}} p_k(t) d\lambda(t), \quad k = 0, 1, 2, \dots$$

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modified moment map

$$\mathbb{R}^{2n} \mapsto \mathbb{R}^{2n} : \quad [m_k]_{k=0}^{2n-1} \mapsto [\alpha_k, \beta_k]_{k=0}^{n-1}$$

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algorithm: Sack et al., 1971; Wheeler, 1974

conditioning: G., 1968; 1982

Matlab

`ab=chebyshev(N, mom, abm)`

Demo #2

elliptic orthogonal polynomials

$$d\lambda(t) = [(1 - \omega^2 t^2)(1 - t^2)]^{-1/2} dt \quad \text{on } [-1, 1],$$
$$0 \leq \omega < 1$$

Demo #2

elliptic orthogonal polynomials

$$d\lambda(t) = [(1 - \omega^2 t^2)(1 - t^2)]^{-1/2} dt \quad \text{on } [-1, 1], \\ 0 \leq \omega < 1$$

Chebyshev moments

$$m_0 = \int_{-1}^1 d\lambda(t), \quad m_k = \frac{1}{2^{k-1}} \int_{-1}^1 T_k(t) d\lambda(t), \quad k \geq 1$$

Demo #2

elliptic orthogonal polynomials

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Chebyshev moments

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Matlab

```
function ab=r_elliptic(N,om2)
abm=r_jacobi(2*N-1,-1/2);
mom=mm_elliptic(N,om2);
ab=chebyshev(N,mom,abm);
```

Demo #2 (cont')

$$\omega^2 = .999, \quad N = 40$$

Demo #2 (cont')

$$\omega^2 = .999, \quad N = 40$$

k	β_k	k	β_k	k	β_k	k	β_k
0	9.68226512	10	0.24936494	20	0.24992062	30	0.24998062
1	0.79378214	11	0.24951641	21	0.24993230	31	0.24998288
2	0.11986767	12	0.24962381	22	0.24994197	32	0.24998485
3	0.22704012	13	0.24970218	23	0.24995003	33	0.24998657
4	0.24106088	14	0.24976074	24	0.24995679	34	0.24998806
5	0.24542853	15	0.24980537	25	0.24996249	35	0.24998937
6	0.24730165	16	0.24983998	26	0.24996732	36	0.24999052
7	0.24825871	17	0.24986721	27	0.24997145	37	0.24999154
8	0.24880566	18	0.24988890	28	0.24997497	38	0.24999243
9	0.24914365	19	0.24990639	29	0.24997800	39	0.24999322



Discrete Measure

Discrete Measure

$$d\lambda_N(t) = \sum_{k=1}^N w_k \delta(t - x_k)$$

inner product $(p, q)_N = \sum_{k=1}^N w_k p(x_k) q(x_k)$

Algorithm #1: Stieltjes, 1884

Discrete Measure

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Algorithm #1: Stieltjes, 1884

Darboux's formulae

$$(D) \quad \begin{cases} \alpha_k = \frac{(t\pi_k, \pi_k)_N}{(\pi_k, \pi_k)_N}, & k = 0, 1, \dots, n-1, \\ \beta_k = \frac{(\pi_k, \pi_k)_N}{(\pi_{k-1}, \pi_{k-1})_N}, & k = 1, 2, \dots, n-1 \end{cases}$$

Discrete Measure

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Darboux's formulae

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plus recurrence relation (R)

$$\begin{aligned} \pi_0 = 1 &\xrightarrow{(D)} \alpha_0, \beta_0 \xrightarrow{(R)} \pi_1 \xrightarrow{(D)} \alpha_1, \beta_1 \xrightarrow{(R)} \dots \\ &\xrightarrow{(D)} \alpha_{n-1}, \beta_{n-1} \end{aligned}$$

Algorithm #2 Lanczos, 1950

$$Q^T \begin{bmatrix} 1 & \sqrt{w_1} & \sqrt{w_2} & \cdots & \sqrt{w_N} \\ \sqrt{w_1} & x_1 & 0 & \cdots & 0 \\ \sqrt{w_2} & 0 & x_2 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \sqrt{w_N} & 0 & 0 & \cdots & x_N \end{bmatrix} Q$$

$$= \begin{bmatrix} 1 & \sqrt{\beta_0} & 0 & \cdots & 0 \\ \sqrt{\beta_0} & \alpha_0 & \sqrt{\beta_1} & \cdots & 0 \\ 0 & \sqrt{\beta_1} & \alpha_1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & \alpha_{N-1} \end{bmatrix}$$

stable version of Lanczos algorithm

Rutishauser, 1963; Gragg and Harrod, 1984

Matlab

$$\left. \begin{array}{l} \text{ab=stieltjes}(n, \text{xw}) \\ \text{ab=lanczos}(n, \text{xw}) \end{array} \right\} n \leq N$$

x_1	w_1
x_2	w_2
\vdots	\vdots
x_N	w_N

XW



Discretization Methods G.; 1968, 1982

Discretization Methods G.; 1968, 1982

basic idea

$$d\lambda(t) \approx d\lambda_N(t), \quad \begin{cases} \alpha_k(d\lambda) \approx \alpha_k(d\lambda_N) \\ \beta_k(d\lambda) \approx \beta_k(d\lambda_N) \end{cases}$$

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example

$$w(t) = (1 - t^2)^{-1/2} + c \text{ on } [-1, 1], \quad c > 0$$

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example

$$w(t) = (1 - t^2)^{-1/2} + c \text{ on } [-1, 1], \quad c > 0$$

discretization

$$(p, q) =$$

$$\int_{-1}^1 p(t)q(t)(1 - t^2)^{-1/2} dt + c \int_{-1}^1 p(t)q(t) dt$$

$$\approx \sum_{k=1}^N w_k^{Ch} p(x_k^{Ch})q(x_k^{Ch}) + c \sum_{k=1}^N w_k^L p(x_k^L)q(x_k^L)$$

in general

discretize $d\lambda$ on

$$\bigcup_{j=1}^s [a_j, b_j]$$

using

taylor-made quadratures

general-purpose quadratures

} on each $[a_j, b_j]$

Matlab

`ab=mcdis(n, eps0, quad, Nmax)`

structure (global variables)

mc, mp, iq

$$AB = \begin{array}{|c|c|} \hline a_1 & b_1 \\ \hline a_2 & b_2 \\ \hline \vdots & \vdots \\ \hline a_{mc} & b_{mc} \\ \hline \end{array}, DM = \begin{array}{|c|c|} \hline x_1 & y_1 \\ \hline x_2 & y_2 \\ \hline \vdots & \vdots \\ \hline x_{mp} & y_{mp} \\ \hline \end{array}$$

example $w(t) = t^\alpha K_0(t)$ on $[0, \infty]$, $\alpha > -1$

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$$K_0(t) = \begin{cases} R(t) + I_0(t) \ln(1/t) & \text{if } 0 < t \leq 1, \\ t^{-1/2} e^{-t} S(t) & \text{if } 1 \leq t < \infty, \end{cases}$$

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$$\begin{aligned} & \int_0^\infty f(t)w(t)dt \\ &= \int_0^1 [R(t)f(t)]t^\alpha dt + \int_0^1 [I_0(t)f(t)]t^\alpha \ln(1/t)dt \\ & \quad + e^{-1} \int_0^\infty [(1+t)^{\alpha-1/2} S(1+t)f(1+t)]e^{-t} dt \end{aligned}$$

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$$AB = \begin{array}{|c|c|} \hline 0 & 1 \\ \hline 0 & 1 \\ \hline 0 & \infty \\ \hline \end{array}$$

$$w_1(t) = t^\alpha$$

$$w_2(t) = t^\alpha \ln(1/t)$$

$$w_3(t) = e^{-t}$$

Modification Algorithms G., 1982

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problem: given the recurrence coefficients of $d\lambda$, generate those of

$$d\lambda_{\text{mod}}(t) = r(t)d\lambda(t), \quad r \geq 0 \text{ on } \text{supp}(d\lambda), \text{ rational}$$

Modification Algorithms G., 1982

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example: Galant, 1971

$$r(t) = s(t - c), \quad c \in \mathbb{R} \setminus \text{supp}(d\lambda), \quad s = \pm 1$$

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example: Galant, 1971

$$r(t) = s(t - c), \quad c \in \mathbb{R} \setminus \text{supp}(d\lambda), \quad s = \pm 1$$

one step of (symmetric) LR algorithm:

$$s[\mathbf{J}_{n+1}(d\lambda) - c\mathbf{I}] = \mathbf{L}\mathbf{L}^T$$

$$\mathbf{J}_n(d\lambda_{\text{mod}}) = (\mathbf{L}^T \mathbf{L} + c\mathbf{I})_{[1:n, 1:n]}$$

nonlinear recurrence algorithm

Matlab

`ab=chri1(N,ab0,c), chri2,...,chri8`

application: induced orthogonal polynomials
(G. and Li, 1993)

$$d\lambda_{\text{mod}}(t) = [\pi_m(t; d\lambda)]^2 d\lambda(t)$$

$$[\pi_m(t; d\lambda)]^2 = \prod_{\mu=1}^m (t - \tau_{\mu})^2$$

m consecutive modifications by quadratic factors

Sobolev Orthogonal Polynomials

Sobolev Orthogonal Polynomials

Sobolev inner product

$$(p, q)_S = \int_{\mathbb{R}} p(t)q(t)d\lambda_0(t) + \int_{\mathbb{R}} p'(t)q'(t)d\lambda_1(t) \\ + \cdots + \int_{\mathbb{R}} p^{(s)}(t)q^{(s)}(t)d\lambda_s(t)$$

Sobolev Orthogonal Polynomials

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Sobolev orthogonal polynomials $\{\pi_k(\cdot; S)\}$

$$(\pi_k, \pi_\ell)_S \begin{cases} = 0, & k \neq \ell \\ > 0, & k = \ell \end{cases}$$

Sobolev Orthogonal Polynomials

Sobolev inner product

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Sobolev orthogonal polynomials $\{\pi_k(\cdot; S)\}$

$$(\pi_k, \pi_\ell)_S \begin{cases} = 0, & k \neq \ell \\ > 0, & k = \ell \end{cases}$$

recurrence relation

$$\pi_{k+1}(t) = t\pi_k(t) - \sum_{j=0}^k \beta_j^k \pi_{k-j}(t), \quad k = 0, 1, 2, \dots$$

Recurrence Matrix

recurrence relation

$$\pi_{k+1}(t) = t\pi_k(t) - \sum_{j=0}^k \beta_j^k \pi_{k-j}(t), \quad k = 0, 1, 2, \dots$$

matrix of recurrence coefficients

$$B_n = \begin{bmatrix} \beta_0^0 & \beta_1^1 & \beta_2^2 & \cdots & \beta_{n-2}^{n-2} & \beta_{n-1}^{n-1} \\ 1 & \beta_0^1 & \beta_1^2 & \cdots & \beta_{n-3}^{n-2} & \beta_{n-2}^{n-1} \\ 0 & 1 & \beta_0^2 & \cdots & \beta_{n-4}^{n-2} & \beta_{n-3}^{n-1} \\ \cdots & \cdots & \cdots & \cdots & \cdots & \cdots \\ 0 & 0 & 0 & \cdots & \beta_0^{n-2} & \beta_1^{n-1} \\ 0 & 0 & 0 & \cdots & 1 & \beta_0^{n-1} \end{bmatrix}$$

Modified Chebyshev Algorithm

Modified Chebyshev Algorithm

modified moments

$$m_k^{(\sigma)} = \int_{\mathbb{R}} p_k(t) d\lambda_{\sigma},$$

$$k = 0, 1, 2, \dots, ; \quad \sigma = 0, 1, \dots, s$$

Modified Chebyshev Algorithm

modified moments

$$m_k^{(\sigma)} = \int_{\mathbb{R}} p_k(t) d\lambda_{\sigma},$$

$$k = 0, 1, 2, \dots, ; \quad \sigma = 0, 1, \dots, s$$

modified moment map

$$\left[m_k^{(\sigma)} \right]_{k=0}^{2n-1}, \quad \sigma = 0, 1, \dots, s \mapsto \mathbf{B}_n$$

Modified Chebyshev Algorithm

modified moments

$$m_k^{(\sigma)} = \int_{\mathbb{R}} p_k(t) d\lambda_{\sigma},$$

$$k = 0, 1, 2, \dots, ; \quad \sigma = 0, 1, \dots, s$$

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$$\left[m_k^{(\sigma)} \right]_{k=0}^{2n-1}, \quad \sigma = 0, 1, \dots, s \mapsto \mathbf{B}_n$$

algorithm G. and Zhang, 1995

conditioning Zhang, 1994

Matlab (for $s = 1$)

`B=chebyshev_sob(N, mom, abm)`

example $d\lambda_0(t) = dt$, $d\lambda_1(t) = \gamma dt$ on $[-1, 1]$
(Althammer, 1962)

modified moments

$p_k(t)$ = monic Legendre

$$m_0^{(0)} = 2, m_0^{(1)} = 2\gamma; \quad m_k^{(0)} = m_k^{(1)} = 0, k > 0$$

Matlab

```
mom=zeros(2,2*N);  
mom(1,1)=2;    mom(2,1)=2*g;  
abm=r_jacobi(2*N-1);  
B=chebyshev_sob(N,mom,abm);
```

Discretized Stieltjes Algorithm

Discretized Stieltjes Algorithm

(G. and Zhang, 1995)

$$\beta_j^k = \frac{(t\pi_k, \pi_{k-j})_S}{(\pi_{k-j}, \pi_{k-j})_S}, \quad j = 0, 1, \dots, k$$

Discretized Stieltjes Algorithm

(G. and Zhang, 1995)

$$\beta_j^k = \frac{(t\pi_k, \pi_{k-j})_S}{(\pi_{k-j}, \pi_{k-j})_S}, \quad j = 0, 1, \dots, k$$

Gauss quadrature discretization

$$(p, q)_{d\lambda_\sigma} \approx \sum_{k=1}^{n_\sigma} w_k^{(\sigma)} p(x_k^{(\sigma)}) q(x_k^{(\sigma)}), \quad \sigma = 0, 1, \dots, s$$

Matlab

```
B=stieltjes_sob(N,s,nd,xw,a0,same)
md=max(nd), a0= $\alpha_0(d\lambda_0)$ 
```

Matlab (cont')

XW=

$x_1^{(0)}$	\dots	$x_1^{(s)}$	$w_1^{(0)}$	\dots	$w_1^{(s)}$
$x_2^{(0)}$	\dots	$x_2^{(s)}$	$w_2^{(0)}$	\dots	$w_2^{(s)}$
\vdots		\vdots	\vdots		\vdots
$x_{md}^{(0)}$	\dots	$x_{md}^{(s)}$	$w_{md}^{(0)}$	\dots	$w_{md}^{(s)}$

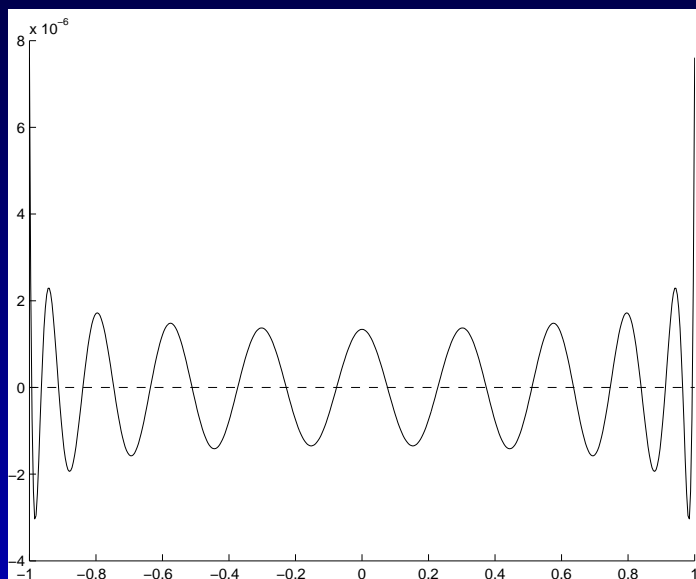
example Althammer's polynomials

Matlab

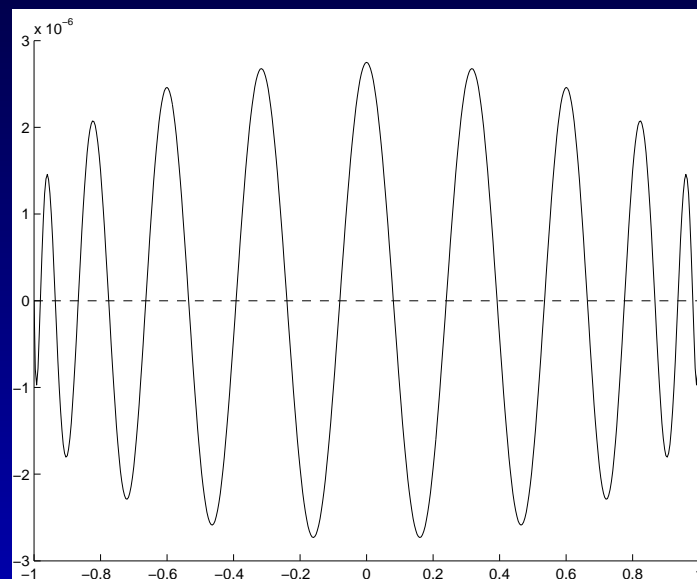
```
s=1; nd=[N N];  
a0=0; same=1;  
ab=r_jacobi(N);  
zw=gauss(N,ab);  
xw=[zw(:,1) zw(:,1) ...  
      zw(:,2) g*zw(:,2)];  
B=stieltjes_sob(N,s,nd,xw,a0,same);
```

Demo #4

Legendre vs Althammer polynomials



Legendre, $n = 20$



Althammer, $n = 20$

Zeros

if

$$\boldsymbol{\pi}^T(t) = [\pi_0(t), \pi_1(t), \dots, \pi_{n-1}(t)]$$

then

$$t\boldsymbol{\pi}^T(t) = \boldsymbol{\pi}^T(t)\mathbf{B}_n + \pi_n(t)\mathbf{e}_n^T$$

Theorem *The zeros τ_ν of π_n are the eigenvalues of \mathbf{B}_n and $\boldsymbol{\pi}^T(\tau_\nu)$ corresponding (left) eigenvectors.*

Matlab

```
z=sobzeros(n,N,B)
```