

# Subset-Sum Approximation

①

Given a set  $S = \{x_1, x_2, \dots, x_n\}$  of positive integers, and a,  $t \geq 1$ , positive integer, find a subset s.t. its sum is largest possible not exceeding  $t$ .

We design a polynomial time approximation scheme (PTAS). {It is fully polynomial since polynomial in  $n$  &  $\frac{1}{\epsilon}$ .}

## An exact solution.

1. compute all subsets  $S' \subseteq S$
2. return the subset with largest sum  $\leq t$ .

Compute the above sums iteratively.

$L_i = \{ \text{sums of all subsets of } \{x_1, x_2, \dots, x_i\} \text{ not exceeding } t \}$ .

Then return the maximum in  $L_n$ .

(2)  
 $L$ : list of positive integers.

$L+x$ : list of positive integers derived from  $L$  by adding  $x$  to each element in  $L$ .

$$L = \{1, 2, 4, 7\}, \quad L+2 = \{3, 4, 6, 9\}.$$

Similarly,

$$S+x : \{s+x \mid s \in S\}.$$

Exact-Subset-Sum( $S, t$ )

$n := |S|;$

$L_0 := \{0\};$

for  $i := 1$  to  $n$  do

$L_i := \text{Merge-Lists}(L_{i-1}, L_{i-1}+x_i)$

remove from  $L_i$  each element  $> t$

endfor

return the largest in  $L_n$ .

Since length of  $L_i$  can be as much as  $2^i$ , the algorithm above runs in exponential time and space.

FPTAS: Trim the list  $L_i$ . Let  $0 < \delta < 1$  be a parameter.

Let  $L'$  be the resulting list from a list  $L$ . For every element  $y$  removed from  $L$ , there is an element  $z \in L'$  s.t.

$$\frac{y}{1+\delta} \leq z \leq y.$$

$z$  "represents"  $y$  in  $L'$ .

Ex.  $L = \{10, 11, 12, 15, 20, 21, 22, 23, 24, 29\}$

$L' = \{10, 12, 15, 20, 23, 29\}$  with  $\delta = 0.1$ .

Trim( $L, \delta$ ) :  $L = \{y_1, \dots, y_m\}$  sorted increasing order.

$m := |L|$

$L' := \{y_1\}$

last :=  $y_1$

for  $i := 2$  to  $m$  do  
  if  $y_i > \text{last} \cdot (1+\delta)$  then  
    append  $y_i$  to  $L'$   
    last :=  $y_i$   
  endif

endfor  
return  $L'$

With  $0 < \epsilon < 1$ , we write the approximation <sup>(4)</sup>  
Scheme

Approx-Subset-Sum ( $S, t, \epsilon$ )

$n := |S|;$

$L_0 := \{0\};$

for  $i := 1$  to  $n$  do

$L_i := \text{Merge-list}(L_{i-1}, L_{i-1} + x_i)$

$L_i := \text{Trim}(L_i, t/2n)$

remove from  $L_i$  each element  $> t$

endfor

return  $z^*$ , the maximum in  $L_n$ .

Ex.  $S = \{104, 102, 201, 101\}$

$t = 308, \epsilon = 0.4$

$\delta = \epsilon/8 = 0.05$

$L_0 = \{0\};$

$L_1 = \{0, 104\}$

$L_2 = \{0, 102, 104, 206\}$

$L_2 = \{0, 102, 206\}$

$L_2 = \{0, 102, 206\}.$

$$L_3 = \{0, 102, 201, 206, 303, 407\}$$

$$L_3 = \{0, 102, 201, 303, 407\}$$

$$L_3 = \{0, 102, 201, 303\}$$

$$L_4 = \{0, 101, 102, 201, 203, 302, 303, 404\}$$

$$L_4 = \{0, 101, 201, 302, 404\}$$

$$L_4 = \{0, 101, 201, 302\}$$

$z^* = 302$  is within  $\epsilon = 40\%$  of the optimal  
 $307 = 104 + 102 + 101$ .

Theorem Approx-Subset-Sum is a FPTAS  
for subset-sum problem.

Proof. Let  $P_i$  be the set of values that can  
be obtained by selecting subsets of  
 $\{x_1, x_2, \dots, x_i\}$  and summing its members.  
Let  $y^* \in P_n$  denote the optimal value.

⑥

$z^* \leq y^*$  (follows from definition).

need to show

$$\frac{y^*}{z^*} \leq 1 + \epsilon.$$

By induction on  $i$ , it can be shown for each  $y \in P_i$ ,  $y \leq t$ , there is a  $z \in L_i$  s.t.

$$\textcircled{1} \quad \frac{y}{(1 + \epsilon/2^n)^i} \leq z \leq y.$$

In particular,  $\frac{y^*}{(1 + \epsilon/2^n)^n} \leq z \leq y^*$

$$\frac{y^*}{z} \leq (1 + \frac{\epsilon}{2^n})^n$$

$$\textcircled{2} \quad \frac{y^*}{z^*} \leq (1 + \frac{\epsilon}{2^n})^n$$

Since  $\frac{d}{dn} (1 + \frac{\epsilon}{2^n})^n > 0$ , the function

$(1 + \frac{\epsilon}{2^n})^n$  increases with  $n$  as it  
its limit  $\lim_{n \rightarrow \infty} (1 + \frac{\epsilon}{2^n})^n = e^{\epsilon/2}$ .

We have

(7)

$$\begin{aligned} \left(1 + \frac{\epsilon}{2n}\right)^n &\leq e^{\epsilon/2} \\ &\leq 1 + \frac{\epsilon}{2} + \left(\frac{\epsilon}{2}\right)^2 \\ &\leq 1 + \epsilon \end{aligned}$$

(3)

From (2) & (3)

$$y^* \leq (1 + \epsilon) z^*$$

To show fully polynomial behavior, we derive a bound on  $|L_i|$ .

Due to trimming two successive elements  $z$  and  $z'$  of  $L_i$  should satisfy

$$z'/z > 1 + \epsilon/2n.$$

Thus,  $L_i$  contains 0, possibly 1, and upto  $\lfloor \log_{1+\epsilon/2n} t \rfloor$  additional values.

$$\text{Thus, } |L_i| \leq \log_{1+\epsilon/2n} t + 2$$

$$= \frac{\ln t}{\ln\left(1 + \frac{\epsilon}{2n}\right)} + 2$$

$$\leq \frac{2n\left(1 + \frac{\epsilon}{2n}\right) \ln t}{\epsilon} + 2 \leq \frac{4n \ln t}{\epsilon} + 2$$

(polynomial in  $n$  &  $\frac{1}{\epsilon}, t$ )