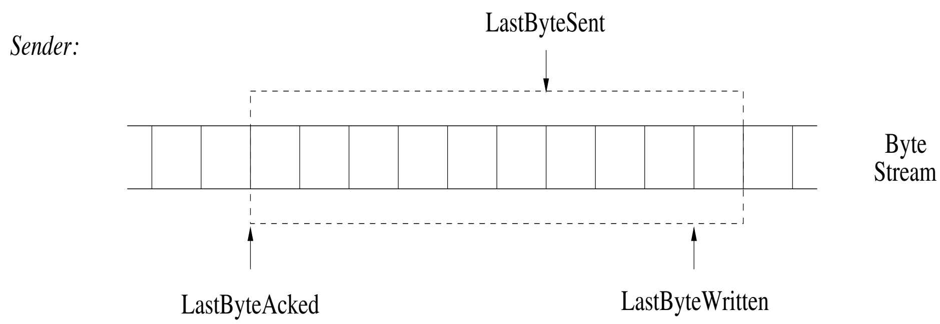
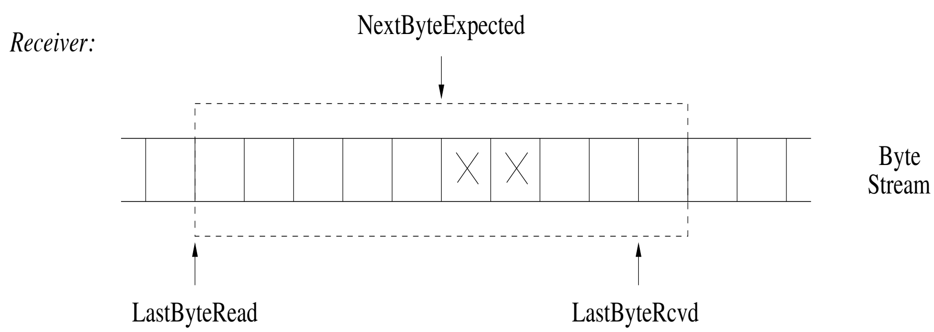
TCP's sliding window protocol:





- sender, receiver maintain buffers `MaxSendBuffer`, `MaxRcvBuffer`

Same as generic sliding window

$\longrightarrow$ data unit: byte, not packet

Sender side: maintain invariants

- `LastByteAcked` $\leq$ `LastByteSent` $\leq$ `LastByteWritten`
- `LastByteWritten`$-$`LastByteAcked` $<$ `MaxSendBuffer`

$\longrightarrow$ buffer flushing (advance window)

$\longrightarrow$ application blocking

- `LastByteSent`$-$`LastByteAcked` $\leq$ `AdvertisedWindow`

$\longrightarrow$ `AdvertisedWindow`: receiver side free space

$\longrightarrow$ upper bounded by receiver window

$\longrightarrow$ throttling effect

How much sender can still send:

$$\texttt{EffectiveWindow} = \texttt{AdvertisedWindow} - (\texttt{LastByteSent} - \texttt{LastByteAcked})$$

$\longrightarrow$ upper bound

$\longrightarrow$ sender may choose to send less

$\longrightarrow$ self-throttling

Affected through sender side variable

$\longrightarrow$ `CongestionWindow`

`EffectiveWindow` update procedure:

$$\texttt{EffectiveWindow} = \texttt{MaxWindow} - (\texttt{LastByteSent} - \texttt{LastByteAcked})$$

where

$$\texttt{MaxWindow} = \min\{\, \texttt{AdvertisedWindow}, \texttt{CongestionWindow} \,\}$$

How to set `CongestionWindow`.

$\longrightarrow$ TCP congestion control

Receiver side: maintain invariants

- $\texttt{LastByteRead} < \texttt{NextByteExpected} \leq \texttt{LastByteRcvd} + 1$
- $\texttt{LastByteRcvd} - \texttt{NextByteRead} < \texttt{MaxRcvBuffer}$

$\longrightarrow$ buffer flushing (advance window)

$\longrightarrow$ application blocking

Thus,

$$\texttt{AdvertisedWindow} = \texttt{MaxRcvBuffer} - (\texttt{LastByteRcvd} - \texttt{LastByteRead})$$

Issues:

How to let sender know of change in receiver window size after `AdvertisedWindow` becomes 0?

- trigger ACK event on receiver side when `AdvertisedWindow` becomes positive
- sender periodically sends 1-byte probing packet

→ design choice: smart sender/dumb receiver

→ same situation for congestion control

Silly window syndrome: Assuming receiver buffer is full, what if application reads one byte at a time with long pauses?

- can cause excessive 1-byte traffic
- if `AdvertisedWindow` $<$ MSS then set `AdvertisedWindow` $\leftarrow 0$

Do not want to send too many 1 B payload packets.

Nagle's method:

- rule: connection can have only one such unacknowledged packet outstanding
- while waiting for ACK, incoming bytes are accumulated (i.e., buffered)

→ compromise between real-time constraints and efficiency

→ facilitates interactive applications

RTT estimation:

$\rightarrow$ important to not underestimate nor overestimate

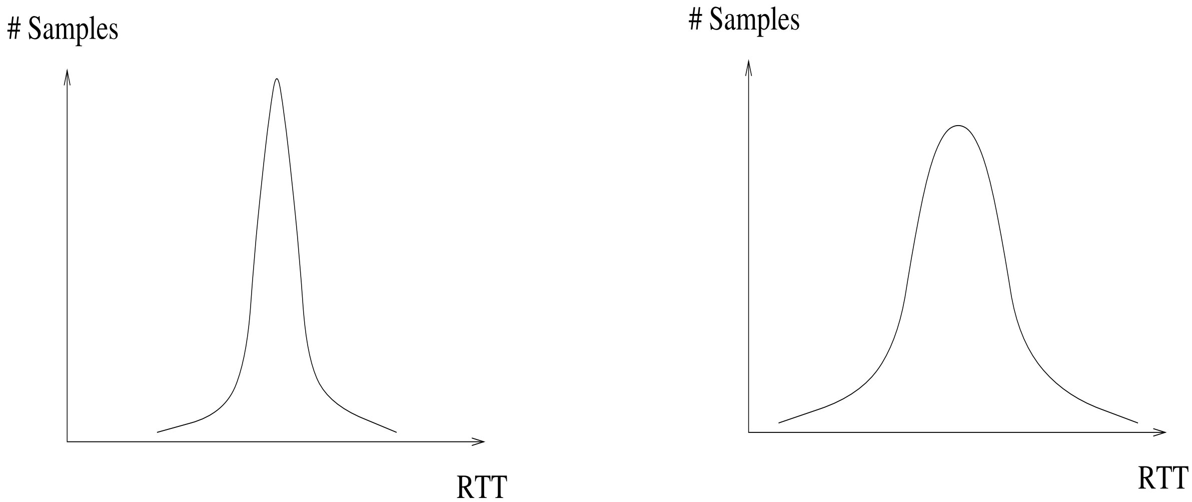
Karn/Partridge: maintain running average with precautions

$$\texttt{EstimateRTT} \leftarrow \alpha \cdot \texttt{EstimateRTT} + \beta \cdot \texttt{SampleRTT}$$

- `SampleRTT` computed by sender using timer
- $\alpha + \beta = 1$; $0.8 \leq \alpha \leq 0.9$, $0.1 \leq \beta \leq 0.2$
- $\texttt{TimeOut} \leftarrow 2 \times \texttt{EstimateRTT}$ or
  $\texttt{TimeOut} \leftarrow 2 \times \texttt{TimeOut}$ (if retransmit)

$\rightarrow$ need to be careful when estimating `SampleRTT`

Issue of variance:



$\longrightarrow$ need to account for variance

$\longrightarrow$ real-world: more messy

Jacobson/Karels:

- $\texttt{Difference} = \texttt{SampleRTT} - \texttt{EstimatedRTT}$
- $\texttt{EstimatedRTT} = \texttt{EstimatedRTT} + \delta \times \texttt{Difference}$
- $\texttt{Deviation} = \texttt{Deviation} + \delta \times (|\texttt{Difference}| - \texttt{Deviation})$

Here $0 < \delta < 1$.

Then

- $\texttt{TimeOut} = \mu \times \texttt{EstimatedRTT} + \phi \times \texttt{Deviation}$

where $\mu = 1$, $\phi = 4$.