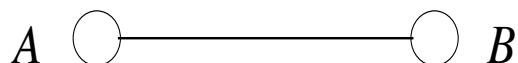


## FUNDAMENTALS OF INFORMATION TRANSMISSION

- applies to both wired and wireless networks
- wireless-specific features discussed separately

### **Sending bits using physical signals**

Simplest case: hosts  $A$  and  $B$  are connected by point-to-point link



→ e.g.,  $A$  wants to send bits 011001 to  $B$

Choices for physical signals

- sound waves: air pressure changes
- underwater sonar: water pressure changes
- light: electromagnetic waves
- what else?

Preferred mode for data communication:

- electromagnetic (EM) waves
- low latency (SOL) and large bandwidth (bps)
- some undesirable properties too

What is an electromagnetic wave?

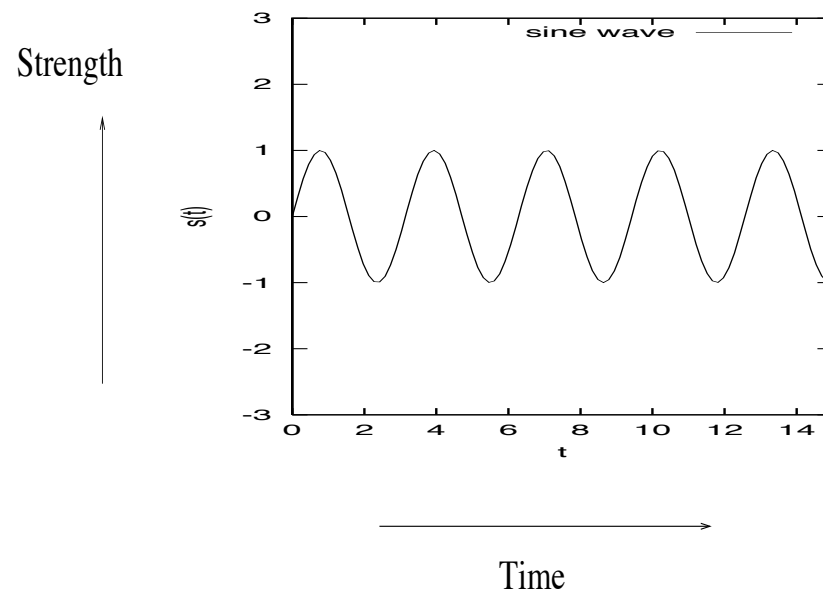
- in principle: a complicated question involving quantum mechanics
- still part of physics and engineering research

In today's systems: only straightforward EM features are exploited

View EM as a physical phenomenon/object which has a strength (or magnitude) that may vary over time.

In simple form, a measurable quantity (or magnitude, amplitude, power, energy) varies in a regular fashion.

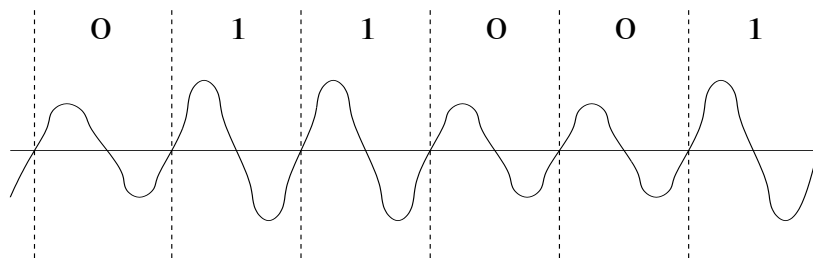
→ i.e., oscillating sine curve



Back to original problem:  $A$  wants to send  $B$  six bits  
011001

→ use magnitude of sine waves

high amplitude represents 1, low amplitude 0 (or vice versa)



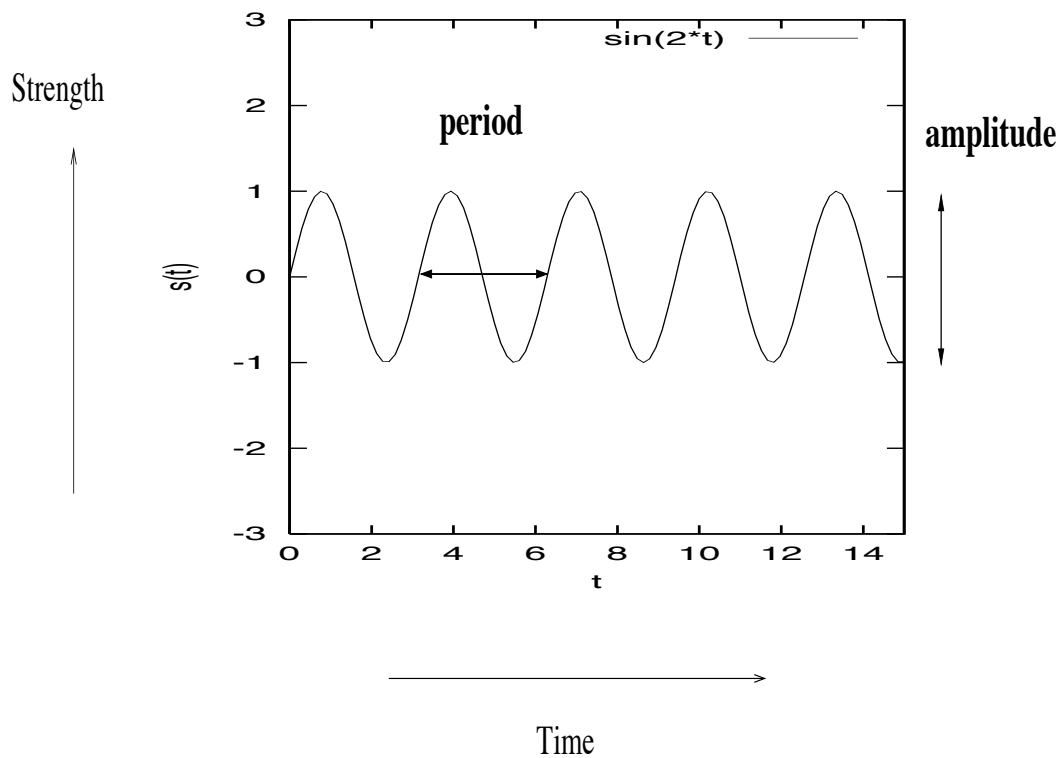
called amplitude modulation (AM)

→ i.e., modulate/manipulate amplitude to send bits

→ same concept as AM radio

→ difference?

Three features of EM as sinusoid:



→ period (also called cycle):  $T$

→ strength: amplitude

→ phase: shift in time

How to utilize these features to communicate bits ...

→ beyond binary AM

Basic properties of sinusoids:

How many periods can we squeeze in per second?

→ frequency:  $1/T$

→ e.g., if period is 1 msec then frequency is 1000 cycles/sec

→ unit called Hertz (Hz)

Another unit: length (m)

→ distance

→ how long is a period

→ i.e., footprint in space

→ empty space: e.g., 1 GHz EM sinusoid about 11.8 inches long

→ fiber optic cable?

In computer networks, by default, frequency is used to specify EM

→ sometimes period is used (esp. high frequency, e.g., 100's of GHz plus)

Example: benefit of using frequency for AM to calculate bps

→ bandwidth (bps) of point-to-point link

→ if frequency is 1 Hz then bandwidth 1 bps

→ if 1 MHz then 1 Mbps

→ if 1 GHz then 1 Gbps

→ if 1 THz then 1 Tbps

Networking problem solved!

→ not quite

Issues with increasing frequency:

One: increasing frequency requires increase in clock rate and processing speed

→ high cost

→ computing systems that control hardware operate at lower speeds

→ heavy lifting: computation

Two: wireless propagation

→ above 10 GHz requires line-of-sight (LOS)

→ complications due to multi-path propagation

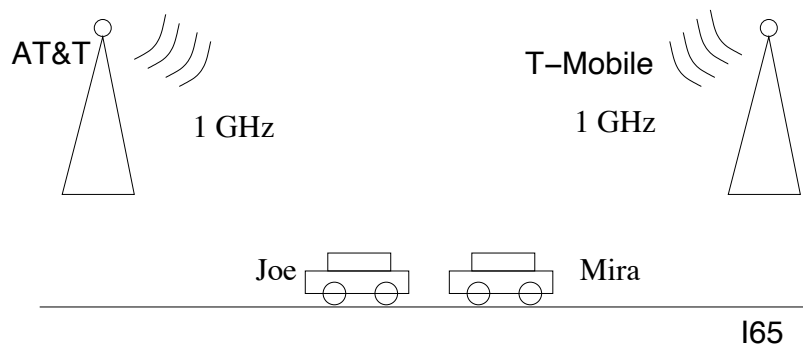
→ echos can be bad (and sometimes good)



Three: multi-user communication

→ not just point-to-point links connecting two parties

Example: wireless interference



Joe receives bits from AT&T's cell tower, Mira from T-Mobile.

→ Joe also hears T-Mobile's signal, Mira hears AT&T's signal

→ interference

→ What does Joe's smart phone actually hear?

Joe's device hears the sum of the two signals

→ property of electromagnetic waves

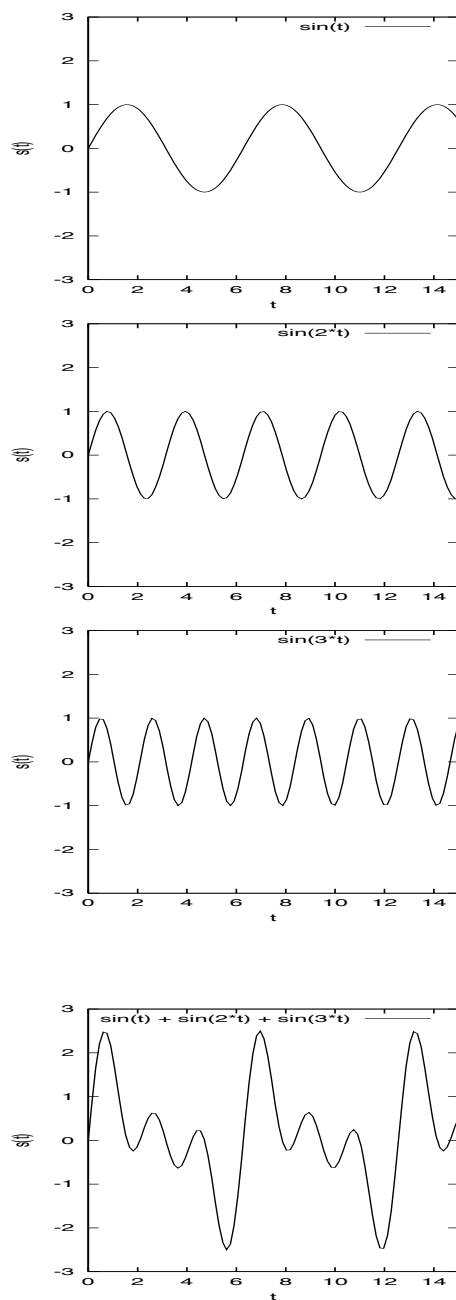
→ i.e., superposition

→ fundamental physics: linear

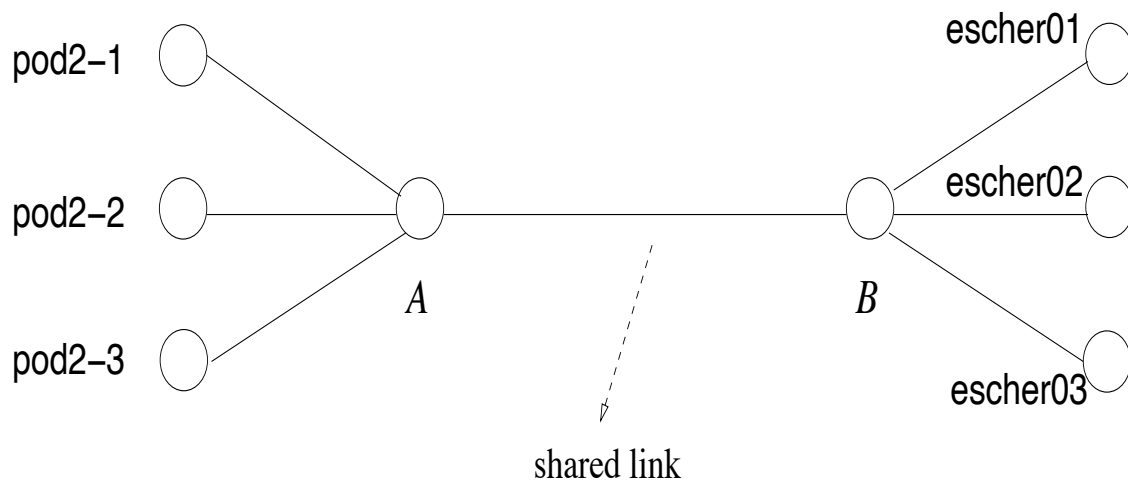
→ amenable to analysis and manipulation

→ basis for modern computer networks

Superposition of three sine waves:



Example: multiplexing (i.e., intentional sharing of resources)



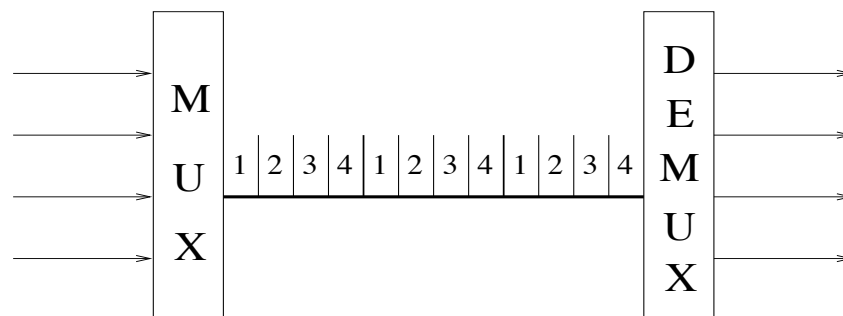
- LWSN B148/HAAS G56 machines:  $A$  and  $B$  are Ethernet switches
- $A$  and  $B$  are routers/switches that forward multiple traffic streams
- structured, orderly access

Splitting time based on AM method of sending bits using sine waves:

→ time-division multiplexing (TDM)

Ex.: four bit streams sharing same link

→ reserve time slots for each bit stream



→ user 1 gets slots 1, 5, 9, etc.

→ user 2 gets slots 2, 6, 10, etc.

→ router *A*: acts as multiplexer (MUX) or combiner

→ router *B*: acts as demultiplexer (DEMUX) or splitter

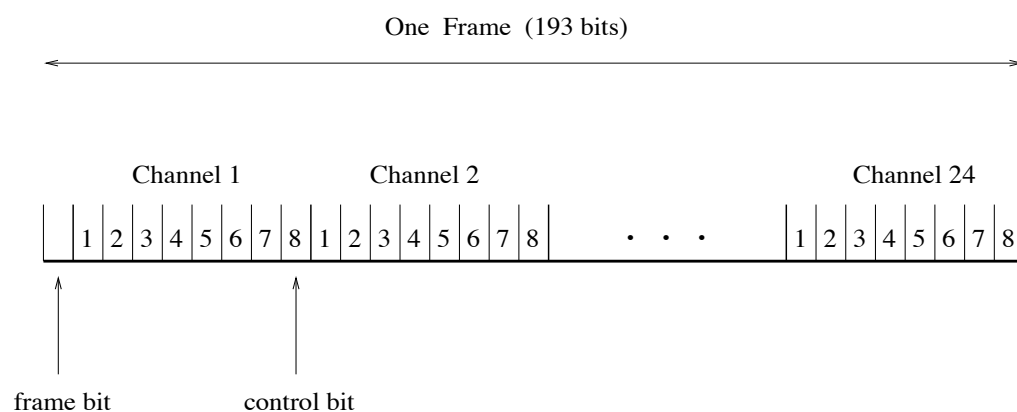
TDM, or TDMA (time-division multiple access) when slots belong to multiple users, is popular in cellular systems and traditional landline telephone systems.

→ simple but fundamental sharing technique

Real-world TDMA example from wired world:

→ T1 carrier (1.544 Mbps)

→ goal: support 24 simultaneous users (“channels”)



Specs of T1 carrier:

- 24 channels (i.e., users)
- time slot: 8-bit block (each user sends 8 consecutive bits)
- $24 \times 8 = 192$  bits of payload
- plus 1 control bit: total 193 bits in a frame (unit of packaged data)
- squeeze 8000 frames into 1 second time interval  
→ frame duration:  $125 \mu\text{sec}$
- total bandwidth (bps):  $8000 \times 193 = 1.544 \text{ Mbps}$
- per channel bandwidth (bps):  $8000 \times 8 = 64 \text{ Kbps}$   
→ landline quality telephony service

At one time, popular service sold by ISPs (mainly) to companies

- 20+ years back, Purdue leased about 6–7 T1 lines for the entire WL campus
- next level T3 line: 44.736 Mbps

Today: residential subscriber can get 1 Gbps or faster download speed

- uplink: significantly slower
- bandwidth asymmetry
- reflects client/server environment



TDMA: important multi-user link transmission technology

- works well if resource (e.g., frequency) is managed by central authority
- single provider
- otherwise: complications

What we want: parallel lanes where multiple bit streams are transmitted simultaneously

- essence of modern high-speed networks
- key technology: use multiple frequencies
- e.g., 1 GHz and 2 GHz for two parallel lanes

How does using multiple frequencies for multiple lanes work?

- classical method
- improvements: our goal
- modern broadband networks

## Roadmap:

- start with CDMA
  - focus on coding: symbol processing
  - conceptual basis for analog methods
- move on to FDMA
  - use analog signals (sinusoid) to send parallel bit streams
  - classical method
  - limitations
- arrive at OFDM (orthogonal frequency division multiplexing)
  - extend FDMA to squeeze in more parallel lanes
  - increase bandwidth (bps)

CDMA motivation: linear algebra approach for sending multiple bit streams

Example: three users Alice, Bob, Mira

→ simplest case: cell tower wants to send each user 1 bit

→ but not TDMA

Assign each user a 3-D vector: called code

→  $(1,0,0)$  for Alice

→  $(0,1,0)$  for Bob

→  $(0,0,1)$  for Mira

To send bit value 1 to Alice, 0 to Bob, 1 to Mira:

→ broadcast vector  $(1,0,1)$  to everyone

→ trivial: not much gained

Allow negative values:

→ send  $(1, -1, 1)$ : 1 means 1, -1 means 0

In general: let positive value means 1, negative value means 0

Example: assign Alice, Bob, Mira code vectors

→ Alice:  $(1, -2, 1)$

→ Bob:  $(3, 5, 7)$

→ Mira:  $(19, 4, -11)$

The code vectors are stored in their smart phones.

Cell tower transmits via broadcast:  $(17, -3, -17)$

→ ignore how the cell tower transmits  $(17, -3, -17)$  via electromagnetic waves

→ upon receiving  $(17, -3, -17)$ , how does Alice know what bit was sent?

Solution: Alice calculates dot product of received vector  $(17, -3, -17)$  with her code vector  $(1, -2, 1)$ .

Definition of dot product: Given two 3-D vectors  $x = (x_1, x_2, x_3)$  and  $y = (y_1, y_2, y_3)$ , their dot (or inner) product is

$$x \circ y = x_1y_1 + x_2y_2 + x_3y_3$$

For Alice:

$$(17, -3, -17) \circ (1, -2, 1) = 17 + 6 - 17 = 6 > 0$$

$\rightarrow$  positive means bit 1

For Bob:  $(17, -3, -17) \circ (3, 5, 7) = 51 - 15 - 119 = -83 < 0$

→ negative means bit 0

For Mira:  $(17, -3, -17) \circ (19, 4, -11) = 323 - 12 + 187 = 498 > 0$

→ positive means bit 1

Why does this work?

→ what is special about  $(1, -2, 1)$ ,  $(3, 5, 7)$ ,  $(19, 4, -11)$

→ where did  $(17, -3, -17)$  come from

The three code vectors are mutually orthogonal:  $x \circ y = 0$

$$\rightarrow (1, -2, 1) \circ (3, 5, 7) = 3 - 10 + 7 = 0$$

$$\rightarrow (1, -2, 1) \circ (19, 4, -11) = 19 - 8 - 11 = 0$$

$$\rightarrow (3, 5, 7) \circ (19, 4, -11) = 57 + 20 - 77 = 0$$

Cell tower's job: send 1 to Alice, 0 to Bob, 1 to Mira

Cell tower computes  $(17, -3, 17)$  to broadcast where

$$\begin{aligned} (+1) \cdot (1, -2, 1) + (-1) \cdot (3, 5, 7) + (+1) \cdot (19, 4, -11) \\ = (17, -3, 17) \end{aligned}$$



When Alice performs dot product of received vector  $(17, -3, -17)$  with her code vector  $(1, -2, 1)$ , it is equivalent to

$$\{(+1) \cdot (1, -2, 1) + (-1) \cdot (3, 5, 7) + (+1) \cdot (19, 4, -11)\} \\ \circ (1, -2, 1)$$

By orthogonality, the second and third terms vanish and what is left is

$$\rightarrow (+1)(1, -2, 1) \circ (1, -2, 1) = 1 + 4 + 1 = 6 > 0$$

$\rightarrow$  taking the dot product with oneself is always positive

For Bob:

$$\rightarrow (17, -3, -17) \circ (3, 5, 7) = 51 - 15 - 119 = -83 < 0$$

$\rightarrow$  negative means bit 0

For Mira:

$$\rightarrow (17, -3, -17) \circ (19, 4, -11) = 323 - 12 + 187 = 498 > 0$$

If we wanted the dot product for Alice to yield +1, Bob -1, Mira, +1, what to do?

Why might we not want the result to be 1, -1, 1, but 6, -83, 498?

CDMA (code division multiple access): using linear algebra, hide the bits to send in the coefficients of the code vectors.

→ in TDMA we divide time to transmit multiple bits in time slots

→ in CDMA, we “divide” code to transmit multiple bits as coefficients

→ coefficients are called spectrum

In CDMA, coding is used to encode multiple bits before transmission using electromagnetic waves occurs.

→ separate (analog) stage

→ omit here: will cover FDMA and OFDMA

→ driver in 90s-20s: QUALCOMM

→ e.g., Verizon, Sprint used CDMA in 3G cellular networks

→ retired in '22 and '23

Origin: military context (long history)

→ if code vectors are chosen to be random, additional feature of security (confidentiality)

Generalize:

To communicate  $n$  bits belonging to  $n$  users

- Set-up: assign  $n$  orthogonal code vectors in  $n$ -dimensional vector space

$$\rightarrow \mathbf{x}^1, \mathbf{x}^2, \dots, \mathbf{x}^n$$

- Sender: to encode  $n$  data bits  $a_1, a_2, \dots, a_n$  (+1 for 1, -1 for 0), compute

$$\rightarrow \mathbf{z} = a_1\mathbf{x}^1 + a_2\mathbf{x}^2 + \dots + a_n\mathbf{x}^n$$

$\rightarrow \mathbf{z}$  is an  $n$ -dimensional vector that hides  $n$  bits in its coefficients (spectra)

$\rightarrow$  convert  $\mathbf{z}$  into analog signal and transmit to all receivers

- Receiver: to decode user  $i$ 'th bit  $a_i$ , receiver computes dot product
  - $\mathbf{z} \circ \mathbf{x}^i = a_i(\mathbf{x}^i \circ \mathbf{x}^i) = a_i \times \text{positive constant}$
  - by orthogonality

Next: borrow the conceptual framework from linear algebra for hiding bits in electromagnetic waves

→ FDMA and OFDMA

→ replace  $n$ -dimensional vectors with continuous complex sinusoids

→ good news: much of the conceptual framework carries over