Shannon showed that there is a fundamental limitation to reliable data transmission.

- ther wider the bandwidth (Hz) the higher the reliable throughput
- the noisier the channel, the smaller the reliable throughput
  - $\rightarrow$  overhead spent dealing with corrupted bits

Channel Coding Theorem (Shannon): Given bandwidth W, signal power  $P_S$ , noise power  $P_N$ , channel subject to white noise,

$$C = W \log \left(1 + \frac{P_S}{P_N}\right)$$
 bps.

 $\rightarrow P_S/P_N$ : signal-to-noise ratio (SNR)

- Increase bandwidth W (Hz) to proportionally increase reliable throughput
  - $\rightarrow$  e.g., FDM, OFDM
  - $\rightarrow$  best possible way
  - $\rightarrow$  wireless bandwidth: scarce resource
- Power control (e.g., handheld wireless devices)
  - $\rightarrow$  trade-off w.r.t. battery power
  - $\rightarrow$  trade-off w.r.t. multi-user interference: doesn't work if everyone increases power
  - $\rightarrow$  signal-to-interference ratio (SIR)

Signal-to-noise ratio (SNR) is expressed as  $dB = 10 \log_{10}(P_S/P_N).$ 

Ex.: Assuming a decibel level of 30, what is the channel capacity of a telephone line?

First, W = 3000 Hz,  $P_S/P_N = 1000$ . Using Channel Coding Theorem,

 $C = 3000 \log 1001 \approx 30$  Kbps.

 $\longrightarrow$  compare against 28.8 Kbps modems

- $\longrightarrow$  what about 56 Kbps modems?
- $\longrightarrow$  xDSL lines?

- $\rightarrow$  modern communication: mainly for digitizing analog audio (music and voice)
- $\rightarrow$  key issue: digitizing time
- $\rightarrow$  digitizing amplitude: less critical due to log-response of auditory system

Sampling Theorem (Nyquist): Given continuous bandlimited signal s(t) with bandwidth W (Hz), s(t) can be reconstructed from its samples if

$$\nu > 2W$$

where  $\nu$  is the sampling rate.

 $\longrightarrow \nu$ : samples per second

- $\rightarrow$  sensitivity: 20 Hz–20 KHz range (roughly 20 KHz)
- $\rightarrow$  voice: 300 Hz–3.3 KHz (roughly 4 KHz)
- T1 TDM line: 1.544 Mbps
- $\rightarrow$  frame size 193 (24 users, 8 bits-per-user, 1 preamble bit)
- $\rightarrow 8000$  samples per second
- $\rightarrow 193 \times 8000 = 1.544$  Mbps
- CD quality audio: 44100 samples per second
- $\rightarrow$  also denoted Hz (44.1 KHz)
- DVD quality audio: 96 samples per second (and higher)