Two types of CDMA systems deployed in practice:

- $\rightarrow$  direct sequence spread spectrum (DSSS)
- $\rightarrow$  frequency hopping spread spectrum (FHSS)

Direct sequence spread spectrum (DSSS):

 $\rightarrow$  what we studied using linear algebra

Each user gets their own vector

- $\rightarrow$  called code vector (i.e., private key)
- $\rightarrow$  code vectors between users: orthogonal
- $\rightarrow$  in practice: also look random (pseudo-random)
- $\rightarrow$  prevents easy eaves dropping

Additional features/variations of DSSS:

Replication: replicate each data bit r-fold  $\rightarrow$  ex.: if r = 3 and data is 1001, then 111000000111  $\rightarrow$  why?

Scramble data bits: one-time pad idea

Ex.:

- data bits 1000100000
- $\bullet$  pseudo-random bits 1010011010
  - $\rightarrow$  private key or one-time pad
  - $\rightarrow$  called chipping sequence
- compute: data bits XOR chipping sequence
  - $\rightarrow 1000100000 \oplus 1010011010 = 0010111010$
  - $\rightarrow$  achieves one-time pad encryption to prevent eavesdropping

- sender transmits XOR'ed bit sequence 0010111010
  - $\rightarrow$  e.g., use amplitude (and/or phase) modulation of carrier wave
  - $\rightarrow$  separate issue
- how does receiver decode sender's data bits?

Other benefit of scrambling with pseudo-random key?  $\rightarrow$  hint: why it's called "spread spectrum" Single-user DSSS: used in 802.11b WLAN

- $\rightarrow$  11-bit chip sequence
- $\rightarrow$  single-user means: two laptops do not use orthogonal code vectors for simultaneous bit transmission
- $\rightarrow$  even use same chip sequence
- $\rightarrow$  multi-user communication: a different method called CSMA (carrier sense multiple access)
- $\rightarrow$  similar to Ethernet's method
- $\rightarrow$  discussed under link layer protocols

Second type of CDMA: frequency hopping spread spectrum (FHSS)

Select *m* carrier frequencies  $f_1, f_2, \ldots, f_m$ .

$$\rightarrow$$
 e.g.,  $m = 5$  with  $f_1 = 101$  MHz,  $f_2 = 102$  MHz, ...,  
 $f_5 = 105$  MHz

To send k bits, select pseudo-random sequence from 1, 2,  $\ldots, m$  of length k

 $\rightarrow$  e.g., if k=10 then 3 5 2 1 4 2 5 3 4 1

Send first bit on frequency  $f_3$ , second bit on  $f_5$ , ..., 10th bit  $f_1$ .

 $\rightarrow$  hop around

 $\rightarrow$  pseudo-random sequence is like private key

Benefits:

- $\rightarrow$  prevents eaves dropping
- $\rightarrow$  resistant to jamming ("spread spectrum")

Drawback?

Used in old IEEE 802.11 WLAN (standards specify both DSSS and FHSS)

Used in old IEEE 802.11 Bluetooth

- $\rightarrow$  79 frequency hopping sequence
- $\rightarrow$  now: part of 802.15
- $\rightarrow$  wireless PAN (personal area network)

- $\rightarrow$  key idea: use carrier waves that are orthogonal
- $\rightarrow$  spectra of carrier frequencies can overlap without causing interference
- $\rightarrow$  old guard band spacing not necessary

What does it mean for sine waves to be orthogonal to each other?

Dot product of two vectors  $x = (x_1, \ldots, x_n)$  and  $y = (y_1, \ldots, y_n)$ 

$$x \circ y = \sum_{i=1}^{n} x_i y_i$$

 $\longrightarrow$  sum of products

Dot product of two sinusoids  $x(t) = \sin f_x t$  and  $y(t) = \sin f_y t$ 

$$x(t) \circ y(t) = \int_{-\infty}^{\infty} (\sin f_x t) (\sin f_y t) dt$$

 $\rightarrow$  again: just a sum of products

 $\rightarrow$  ex.: sin t and sin 2t

More generally:  $x(t) \circ y(t) = \int_{-\infty}^{\infty} e^{if_x t} e^{-if_y t} dt$ 

 $\rightarrow$  since Fourier transform involves complex sinusoids  $\rightarrow$  but same form: sum of products

How to get N mutually orthogonal sinusoids?

- Suppose available frequency lies between f<sub>a</sub> and f<sub>b</sub>

  → bandwidth: W = f<sub>b</sub> f<sub>a</sub>
  → ex.: f<sub>a</sub> = 2.4 GHz, f<sub>b</sub> = 2.5 GHz, W = 100 MHz

  Choose N carrier frequencies as

  f f + (W/N) f + 2(W/N)
  - $\rightarrow f_b, f_b + (W/N), f_b + 2(W/N), \ldots$
  - $\rightarrow$  ex.: N = 100
  - $\rightarrow$  2.4 GHz, 2.401 GHz, 2.402 GHz, ..., 2.499 GHz

Can we squeeze in arbitrarily many carrier frequencies?  $\rightarrow$  in principle, yes; in practice, no

Example: indoor wireless signal propagation

- $\rightarrow$  time duration of single bit: called symbol period  $\tau$
- $\rightarrow$  cannot be too short due to multi-path propagation which causes delay spread
- $\rightarrow$  i.e., time delayed echos may overlap with next bit transmission
- $\rightarrow$  called inter-symbol interference (ISI)
- $\rightarrow$  different from ICI (inter-channel interference)

Given symbol period  $\tau$  to prevent ISI, cannot send bits faster than  $\bar{f} = 1/\tau$  Hz.

 $\rightarrow$  use as orthogonal frequency spacing

Hence number of carrier frequencies is

 $\rightarrow N = W/\bar{f}$ 

Example (wireless): IEEE 802.11g WLANs

 $\rightarrow$  uses OFDM

$$\rightarrow$$
 symbol time  $\tau = 3.2 \ \mu s$ 

 $\rightarrow$  part of IEEE standard

 $\rightarrow \bar{f} = 1/\tau = 312.5 \text{ kHz}$ 

$$\rightarrow W = 20 \text{ MHz}, N = W/\bar{f} = 64$$

Example (wireline): ADSL

- $\rightarrow$  frequency spacing influenced by noise factors
- $\rightarrow$  ADSL:  $\bar{f} = 4.3125 \text{ kHz}$
- $\rightarrow$  part of ITU G.992.1 standard
- $\rightarrow$  UTP (unshielded twisted pair) copper wire
- $\rightarrow$  frequency band: 0–1.104 MHz

$$\rightarrow N = W/\bar{f} = 256$$

Shannon showed that there is a fundamental limitation to reliable data transmission.

- ther wider the bandwidth (Hz) the higher the reliable throughput
- the noisier the channel, the smaller the reliable throughput
  - $\rightarrow$  overhead spent dealing with corrupted bits

Channel Coding Theorem (Shannon): Given bandwidth W, signal power  $P_S$ , noise power  $P_N$ , channel subject to white noise,

$$C = W \log \left(1 + \frac{P_S}{P_N}\right)$$
 bps.

 $\rightarrow P_S/P_N$ : signal-to-noise ratio (SNR)

- Increase bandwidth W (Hz) to proportionally increase reliable throughput
  - $\rightarrow$  e.g., FDM, OFDM
  - $\rightarrow$  best possible way
  - $\rightarrow$  wireless bandwidth: scarce resource
- Power control (e.g., handheld wireless devices)
  - $\rightarrow$  trade-off w.r.t. battery power
  - $\rightarrow$  trade-off w.r.t. multi-user interference: doesn't work if everyone increases power
  - $\rightarrow$  signal-to-interference ratio (SIR)

Signal-to-noise ratio (SNR) is expressed as  $dB = 10 \log_{10}(P_S/P_N).$ 

Ex.: Assuming a decibel level of 30, what is the channel capacity of a telephone line?

First, W = 3000 Hz,  $P_S/P_N = 1000$ . Using Channel Coding Theorem,

 $C = 3000 \log 1001 \approx 30$  Kbps.

 $\longrightarrow$  compare against 28.8 Kbps modems

- $\longrightarrow$  what about 56 Kbps modems?
- $\longrightarrow$  xDSL lines?

- $\rightarrow$  modern communication: mainly for digitizing analog audio (music and voice)
- $\rightarrow$  key issue: digitizing time
- $\rightarrow$  digitizing amplitude: less critical due to log-response of auditory system

Sampling Theorem (Nyquist): Given continuous bandlimited signal s(t) with bandwidth W (Hz), s(t) can be reconstructed from its samples if

$$\nu > 2W$$

where  $\nu$  is the sampling rate.

 $\longrightarrow \nu$ : samples per second

- $\rightarrow$  sensitivity: 20 Hz–20 KHz range (roughly 20 KHz)
- $\rightarrow$  voice: 300 Hz–3.3 KHz (roughly 4 KHz)
- T1 TDM line: 1.544 Mbps
- $\rightarrow$  frame size 193 (24 users, 8 bits-per-user, 1 preamble bit)
- $\rightarrow 8000$  samples per second
- $\rightarrow 193 \times 8000 = 1.544$  Mbps
- CD quality audio: 44100 samples per second
- $\rightarrow$  also denoted Hz (44.1 KHz)
- DVD quality audio: 96 samples per second (and higher)