

# Information Security

## CS 526



### Topic 11: Key Distribution & Agreement, Secure Communication

# Readings for This Lecture

- On Wikipedia
  - [Needham-Schroeder protocol](#)  
(only the symmetric key part)
  - [Public Key Certificates](#)
  - [HTTP Secure](#)



# Outline and Objectives

- Key distribution among multiple parties
- Kerberos
- Distribution of public keys, with public key certificates
- Diffie-Hellman Protocol
- TLS/SSL/HTTPS

# Key Agreement among Multiple Parties

- For a group of  $N$  parties, every pair needs to share a different key
  - What is the number of keys?
- Solutions
  - Symmetric Encryption - Use a central authority, a.k.a. (TTP).
  - Asymmetric Encryption – PKI.

# Needham-Schroeder Shared-Key Protocol

- Parties: A, B, and trusted server T
- Setup: A and T share  $K_{AT}$ , B and T share  $K_{BT}$
- Goal: Mutual entity authentication between A and B; key establishment
- Messages:

$$A \rightarrow T: A, B, N_A \quad (1)$$

$$A \leftarrow T: E[K_{AT}](N_A, B, k, E[K_{BT}](k, A)) \quad (2)$$

$$A \rightarrow B: E[K_{BT}](k, A) \quad (3)$$

$$A \leftarrow B: E[k](N_B) \quad (4)$$

$$A \rightarrow B: E[k](N_B-1) \quad (5)$$

What bad things can happen if there is no  $N_A$ ?

Another subtle flaw in Step 3.

# Kerberos

- Implements the idea of Needham-Schroeder protocol
- Kerberos is a **network authentication protocol**
- Provides authentication and secure communication
- Relies entirely on **symmetric cryptography**
- Developed at MIT:  
<http://web.mit.edu/kerberos/www>
- Used in many systems, e.g., Windows 2000 and later as default authentication protocol



# Kerberos Overview

- One issue of Needham-Schroeder – Needs  $[K_{AT}]$  for every communication.
- Kerberos solution:
  - Separates TTP into an AS and a TGS.
- The client authenticates to AS using a long-term *shared secret* and receives a TGT [SSO].
- Use this TGT to get additional tickets from TGS without resorting to using the shared secret.

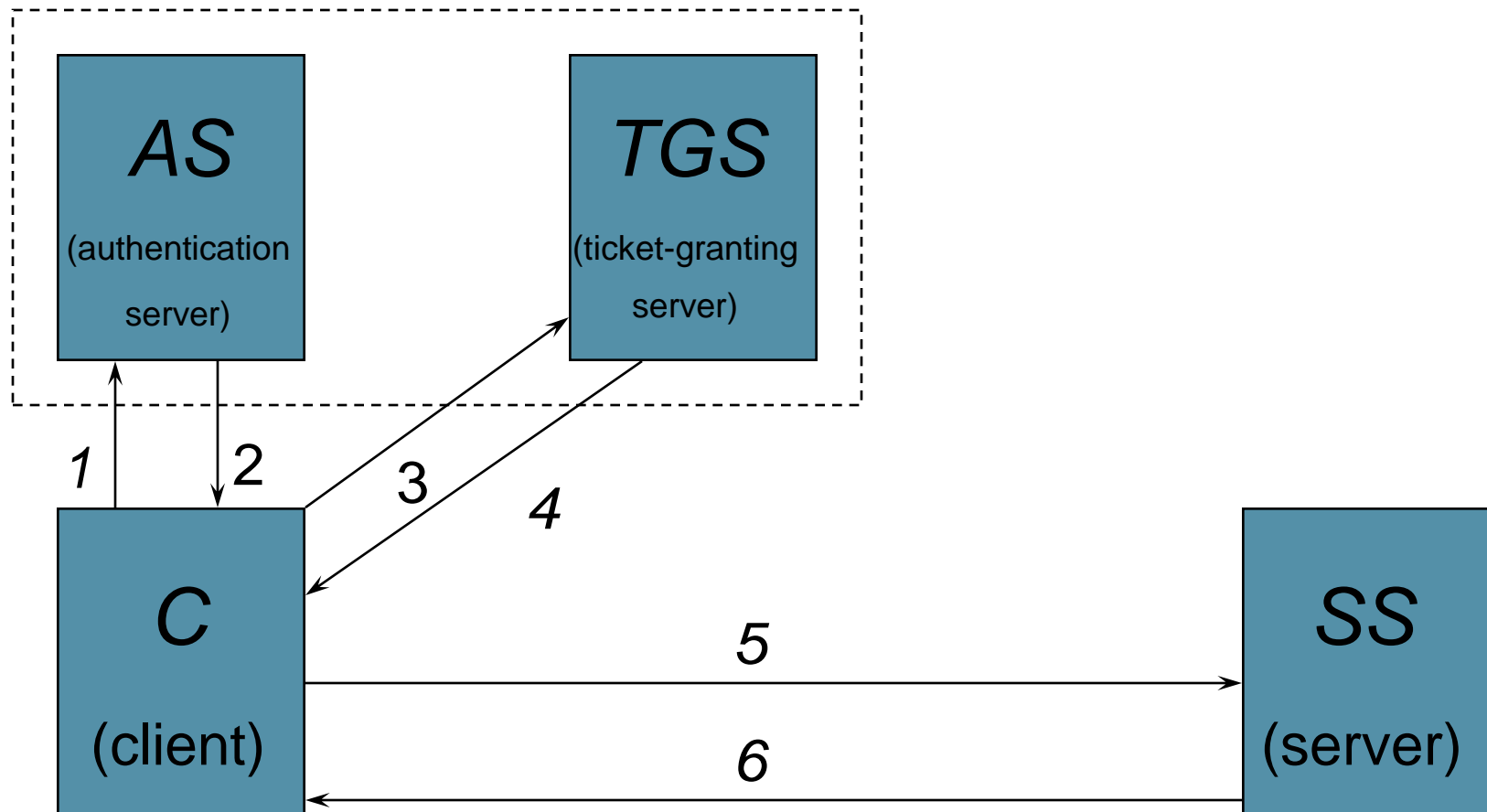
AS = Authentication Server

TGS = Ticket Granting Server

SS = Service Server

TGT = Ticket Granting Ticket

# Kerberos Protocol - 1





# Kerberos Protocol – 2 (Simplified)

1.  $C \rightarrow AS: TGS \parallel N_C$

2.  $AS \rightarrow C: \{K_{C,TGS} \parallel C\}_{K_{AS,TGS}} \parallel \{K_{C,TGS} \parallel N_C \parallel TGS\}_{K_{AS,C}}$

(Note that the **first** part of message 2 is the **ticket granting ticket (TGT)** for the TGS)

3.  $C \rightarrow TGS: SS \parallel N'_C \parallel \{K_{C,TGS} \parallel C\}_{K_{AS,TGS}} \parallel \{C \parallel T_1\}_{K_{C,TGS}}$

4.  $TGS \rightarrow C: \{K_{C,SS} \parallel C\}_{K_{TGS,SS}} \parallel \{K_{C,SS} \parallel N'_C \parallel SS\}_{K_{C,TGS}}$

(Note that the **first** part in message 4 is the **ticket** for the server S).

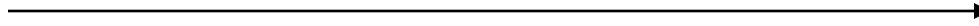
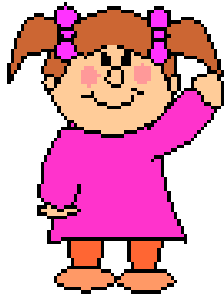
5.  $C \rightarrow SS: \{K_{C,SS} \parallel C\}_{K_{TGS,SS}} \parallel \{C \parallel T_2\}_{K_{C,SS}}$

6.  $SS \rightarrow C: \{T_3\}_{K_{C,SS}}$

# Kerberos Drawback

- Single point of failure:
- Security partially depends on tight clock synchronization.
- Useful primarily inside an organization
  - Does it scale to Internet? What is the main difficulty?

# Public Keys and Trust



- Public Key:  $P_A$
- Secret key:  $S_A$

- Public Key:  $P_B$
- Secret key:  $S_B$

- **How are public keys stored?**
- **How to obtain the public key?**
- **How does Bob know or 'trusts' that  $P_A$  is Alice's public key?**

# Distribution of Public Keys

- **Public announcement:** users distribute public keys to recipients or broadcast to community at large.
- **Publicly available directory:** can obtain greater security by registering keys with a public directory.
- Both approaches have problems, and are vulnerable to forgeries

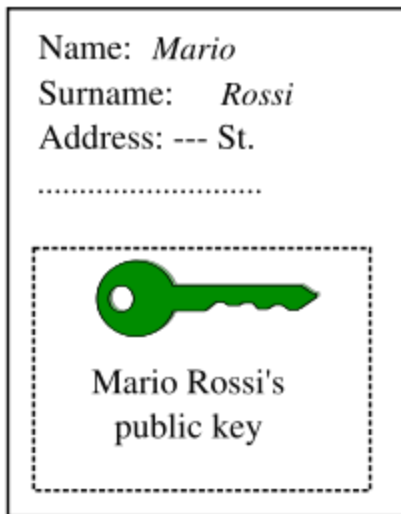


# Public-Key Certificates

- A certificate binds identity (or other information) to public key
- Contents digitally signed by a trusted Public-Key or Certificate Authority (CA)
  - Can be verified by anyone who knows the public-key authority's public-key.
- For Alice to send an encrypted message to Bob, obtains a certificate of Bob's public key

# Public Key Certificates

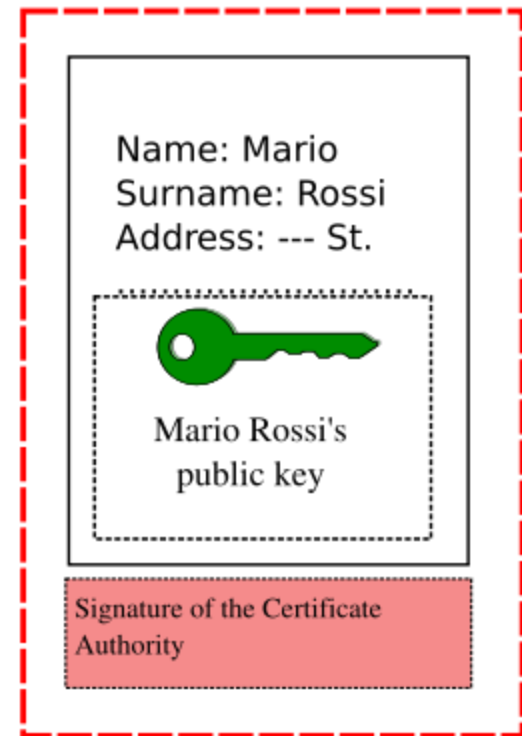
Document containing the public key and identity for Mario Rossi



Certificate Authority's private key



## Mario Rossi's Certificate



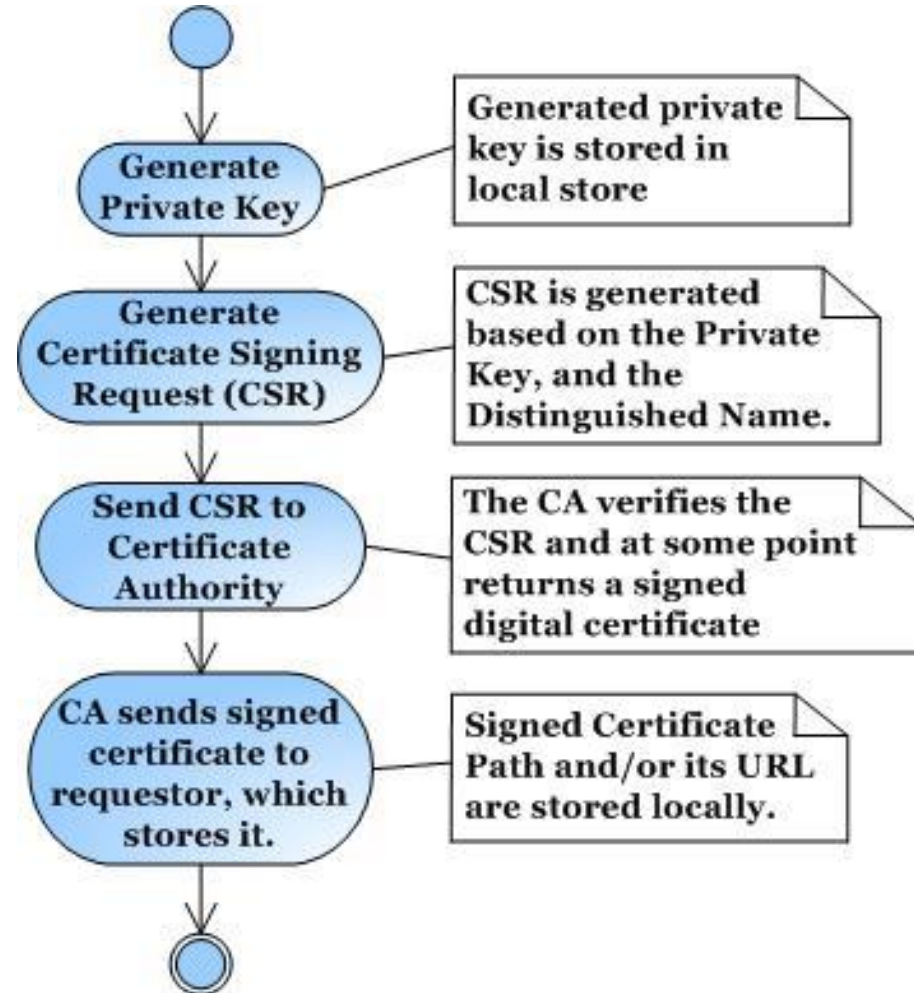
Document signed by the Certificate Authority

# X.509 Certificates

- Part of X.500 directory service standards.
  - Started in 1988
- Defines framework for authentication services:
  - Defines that public keys stored as **certificates** in a public directory.
  - Certificates are **issued and signed** by an entity called **certification authority (CA)**.
- Used by numerous applications: SSL, IPsec, SET
- Example: see certificates accepted by your browser

# How to Obtain a Certificate?

- Define your own CA (use openssl or Java Keytool)
  - Certificates unlikely to be accepted by others
- Obtain certificates from one of the vendors: VeriSign, Thawte, and many others





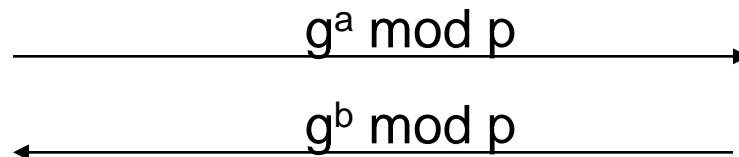
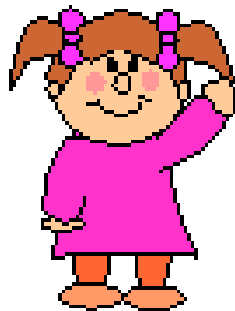
# CAs and Trust

- Certificates are trusted if signature of CA verifies
- Chain of CA's can be formed, head CA is called root CA
- In order to verify the signature, the public key of the root CA should be obtain.
- TRUST is centralized (to root CA's) and hierarchical
- What bad things can happen if the root CA system is compromised?
- How does this compare with the TTP in Needham/Schroeder protocol?

# Key Agreement: Diffie-Hellman Protocol

Key agreement protocol, both A and B contribute to the key

Setup:  $p$  prime and  $g$  generator of  $Z_p^*$ ,  $p$  and  $g$  public.



Pick random, secret (a)

Compute and send  $g^a \bmod p$

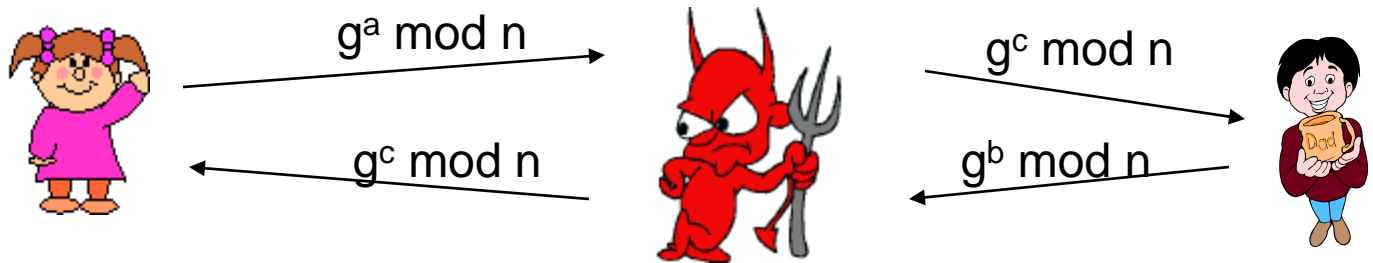
$$K = (g^b \bmod p)^a = g^{ab} \bmod p$$

Pick random, secret (b)

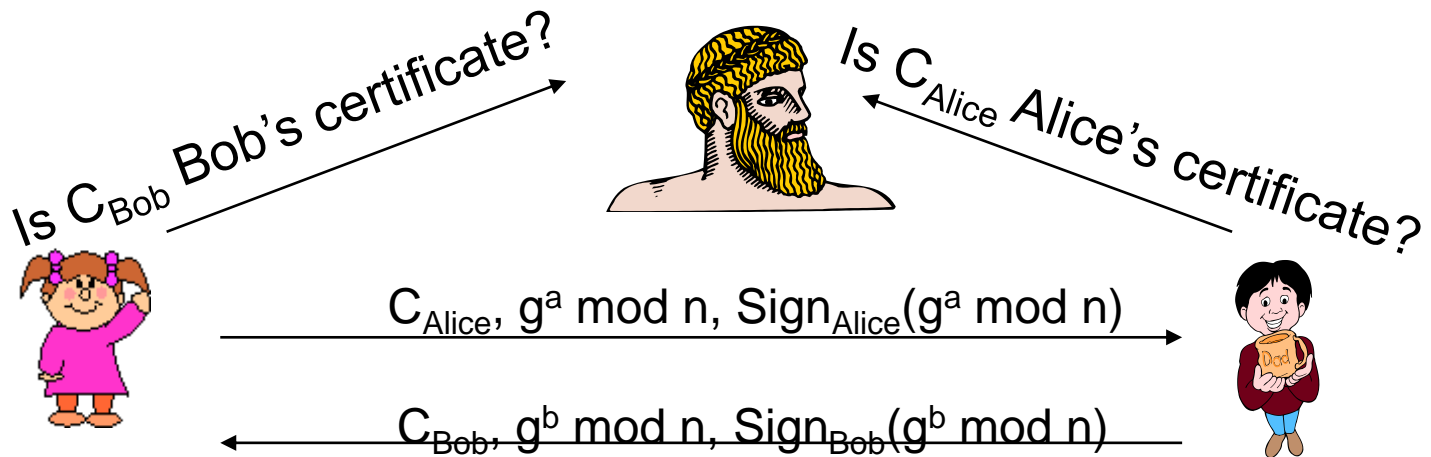
Compute and send  $g^b \bmod p$

$$K = (g^a \bmod p)^b = g^{ab} \bmod p$$

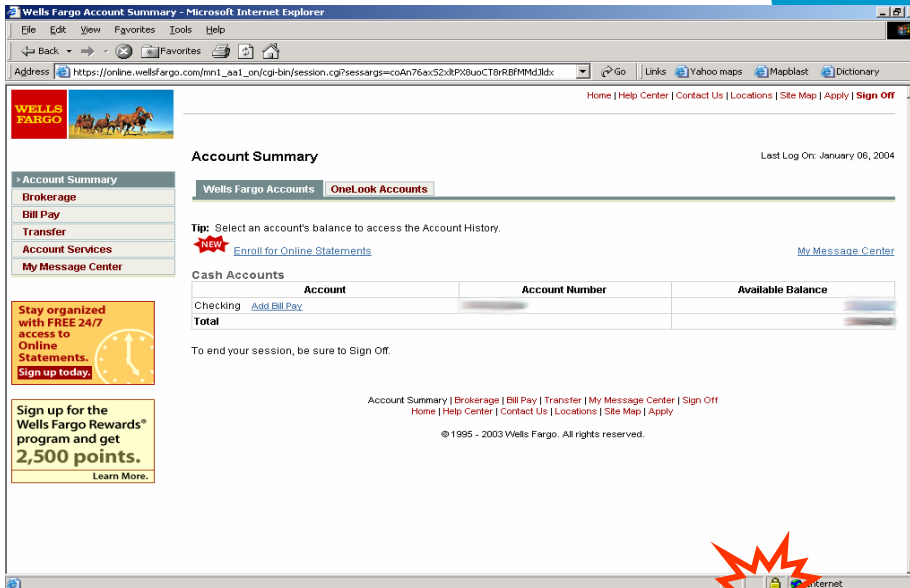
# Authenticated Diffie-Hellman



- Alice computes  $g^{ac} \bmod n$  and Bob computes  $g^{bc} \bmod n$  !!!



# Secure communication



# Transport Layer Security (TLS)

- Predecessors: Secure socket layer (SSL): Versions 1.0, 2.0, 3.0
- TLS 1.0 (SSL 3.1); Jan 1999
- TLS 1.1 (SSL 3.2); Apr 2006
- TLS 1.2 (SSL 3.3); Aug 2008
- Standard for Internet security
  - Originally designed by Netscape
  - Goal: “... provide privacy and reliability between two communicating applications”
- Two main parts
  - Handshake Protocol
    - Establish shared secret key using public-key cryptography
    - Signed certificates for authentication
  - Record Layer
    - Transmit data using negotiated key, encryption function

# Usage of SSL/TLS

- Applied on top of transport layer (typically TCP)
- Used to secure HTTP (HTTPS), SMTP, etc.
- One or both ends can be authenticated using public key and certificates
  - Typically only the server is authenticated
- Client & server negotiate a cipher suite, which includes
  - A key exchange algorithm, e.g., RSA, Diffie-Hellman, SRP, etc.
  - An encryption algorithm, e.g., RC4, Triple DES, AES, etc.
  - A MAC algorithm, e.g., HMAC-MD5, HMC-SHA1, etc.

# Viewing HTTPS web sites

- Browser needs to communicate to the user the fact that HTTPS is used
  - E.g., a golden lock indicator on the bottom or on the address bar
  - Check some common websites
  - When users correctly process this information, can defeat phishing attacks
  - Security problems exist
    - People don't know about the security indicator
    - People forgot to check the indicator
    - Browser vulnerabilities enable incorrect indicator to be shown
    - Use confusing URLs, e.g.,
      - [https:// homebanking.purdueefcu.com@host.evil.com/](https://homebanking.purdueefcu.com@host.evil.com/)
    - Stored certificate authority info may be changed

# Coming Attractions ...

- Software vulnerabilities

