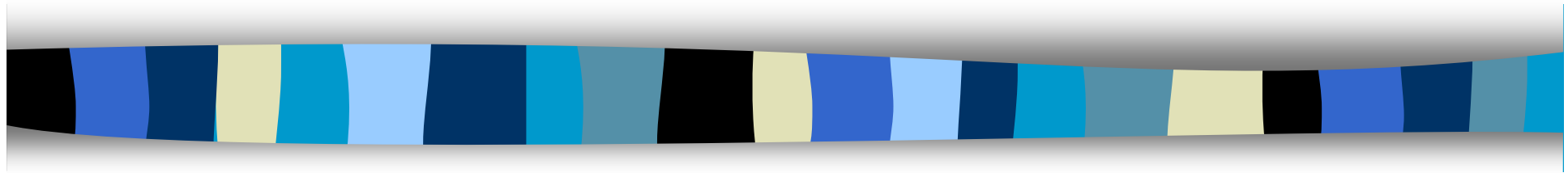


Computer Security

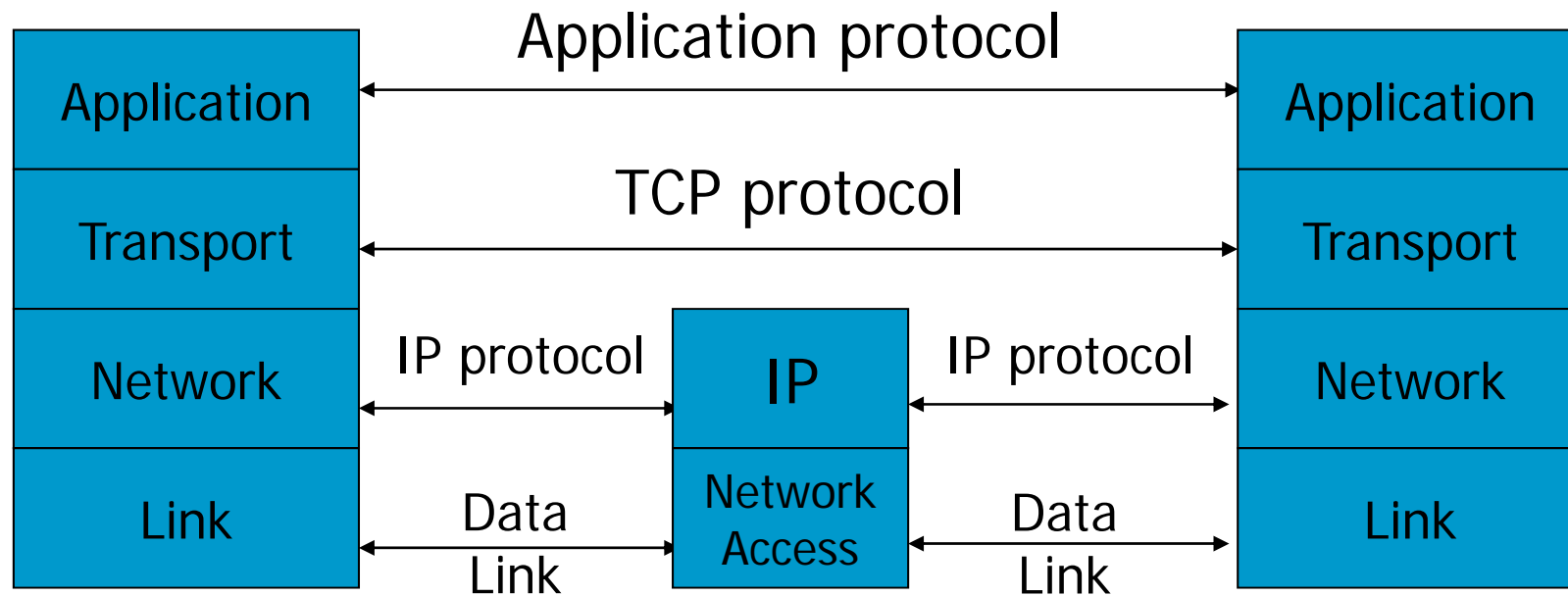
CS 426

Lecture 33



Network Security (1)

Network Protocols Stack



Types of Addresses in Internet

- Media Access Control (MAC) addresses in the network access layer
 - Associated w/ network interface card (NIC)
 - 48 bits or 64 bits
- IP addresses for the network layer
 - 32 bits for IPv4, and 128 bits for IPv6
 - E.g., 128.3.23.3
- IP addresses + ports for the transport layer
 - E.g., 128.3.23.3:80
- Domain names for the application/human layer
 - E.g., www.purdue.edu

Routing and Translation of Addresses

- Translation between IP addresses and MAC addresses
 - Address Resolution Protocol (ARP) for IPv4
 - Neighbor Discovery Protocol (NDP) for IPv6
- Routing with IP addresses
 - TCP, UDP, IP for routing packets, connections
 - Border Gateway Protocol for routing table updates
- Translation between IP addresses and domain names
 - Domain Name System (DNS)

Threats in Networking

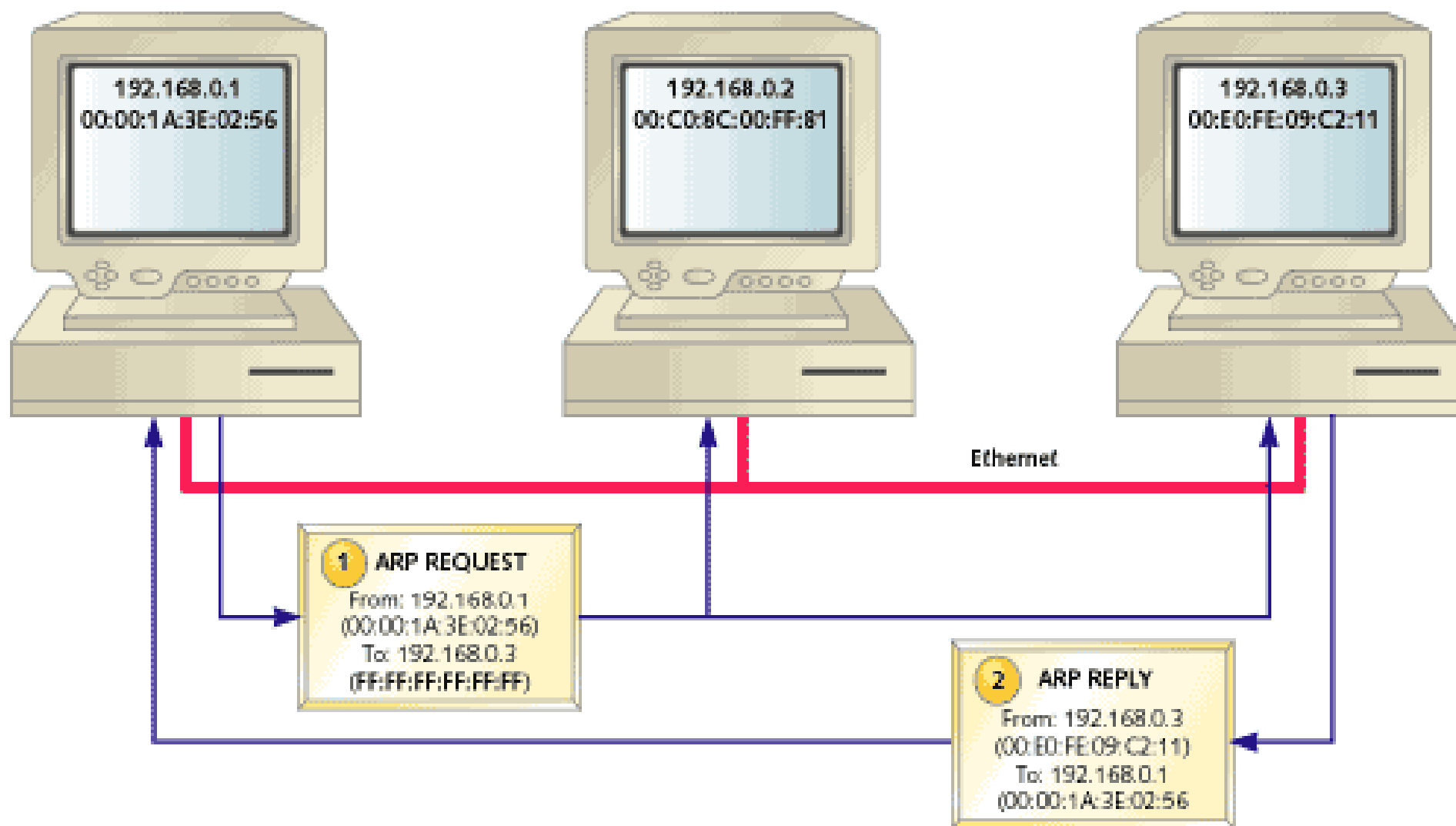
- Confidentiality
 - Packet sniffing
- Integrity
 - Session hijacking
- Availability
 - Denial of service attacks
- Common
 - Address translation poisoning attacks
 - Routing attacks

Concrete Security Problems

- ARP is not authenticated
 - APR spoofing (or ARP poisoning)
- Network packets pass by untrusted hosts
 - Packet sniffing
- TCP state can be easy to guess
 - TCP spoofing attack
- Open access
 - Vulnerable to DoS attacks
- DNS is not authenticated
 - DNS poisoning attacks

Address Resolution Protocol (ARP)

- Primarily used to translate IP addresses to Ethernet MAC addresses
 - The device driver for Ethernet NIC needs to do this to send a packet
- Also used for IP over other LAN technologies, e.g., FDDI, or IEEE 802.11
- Each host maintains a table of IP to MAC addresses
- Message types:
 - ARP request
 - ARP reply
 - ARP announcement



<http://www.netrino.com/Embedded-Systems/How-To/ARP-RARP>

ARP Spoofing (ARP Poisoning)

- Send fake or 'spoofed', ARP messages to an Ethernet LAN.
 - To have other machines associate IP addresses with the attacker's MAC
- Defenses
 - static ARP table
 - DHCP snooping (use access control to ensure that hosts only use the IP addresses assigned to them, and that only authorized DHCP servers are accessible).
 - detection: Arpwatch (sending email when updates occur),
- Legitimate use
 - redirect a user to a registration page before allow usage of the network

IP Internet Protocol

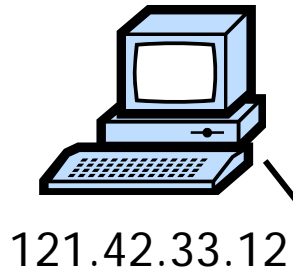
- Connectionless
 - Unreliable
 - Best effort
- Transfer datagram
 - Header
 - Data

Version	Header Length
Type of Service	
Total Length	
Identification	
Flags	Fragment Offset
Time to Live	
Protocol	
Header Checksum	
Source Address of Originating Host	
Destination Address of Target Host	
Options	
Padding	
IP Data	

IP Routing



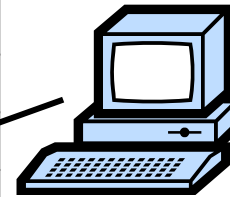
Meg



121.42.33.12

Packet	
Source	121.42.33.12
Destination	132.14.11.51
Sequence	5

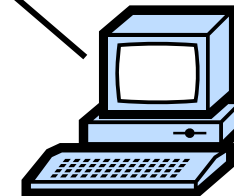
Office gateway



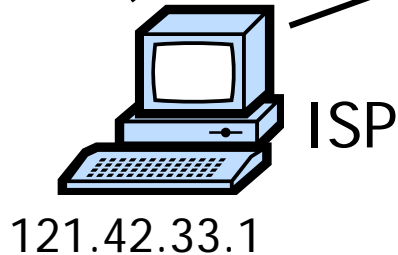
132.14.11.1



Tom



132.14.11.51

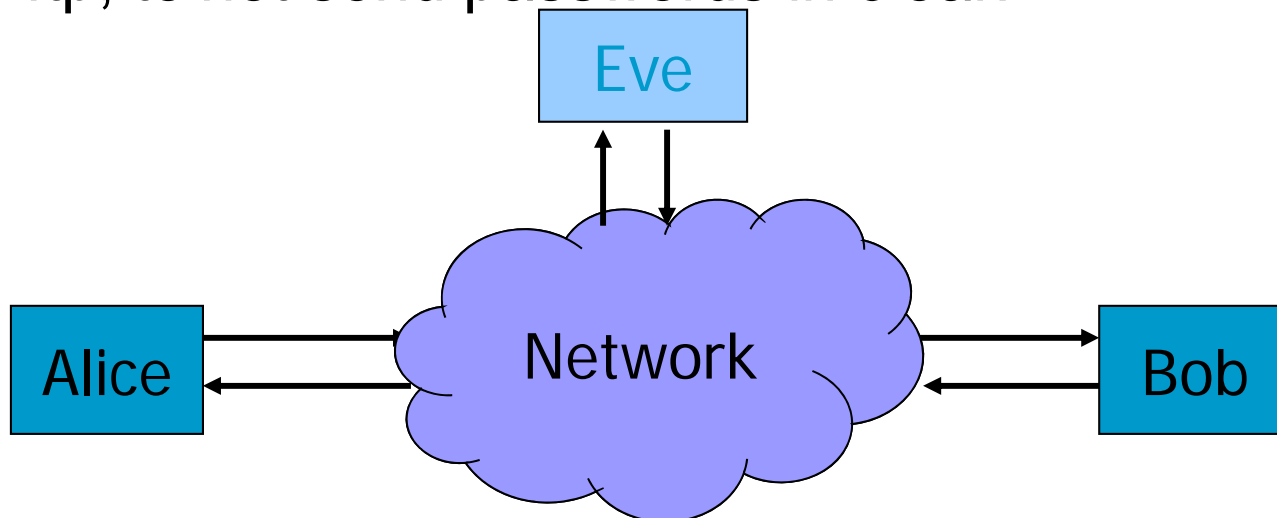


121.42.33.1

- Internet routing uses numeric IP address
- Typical route uses several hops

Packet Sniffing

- Promiscuous Network Interface Card reads all packets
 - Read all unencrypted data (e.g., “ngrep”)
 - ftp, telnet send passwords in clear!



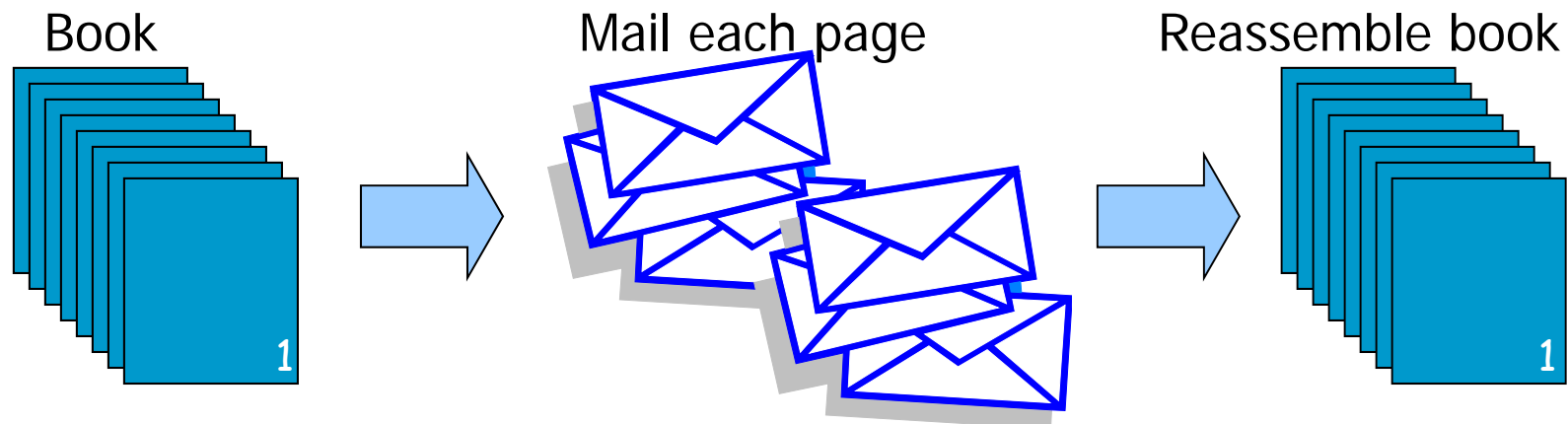
Prevention: Encryption (IPSEC, TLS)

User Datagram Protocol

- IP provides routing
 - IP address gets datagram to a specific machine
- UDP separates traffic by port (16-bit number)
 - Destination port number gets UDP datagram to particular application process, e.g., 128.3.23.3:53
 - Source port number provides return address
- Minimal guarantees
 - No acknowledgment
 - No flow control
 - No message continuation

Transmission Control Protocol

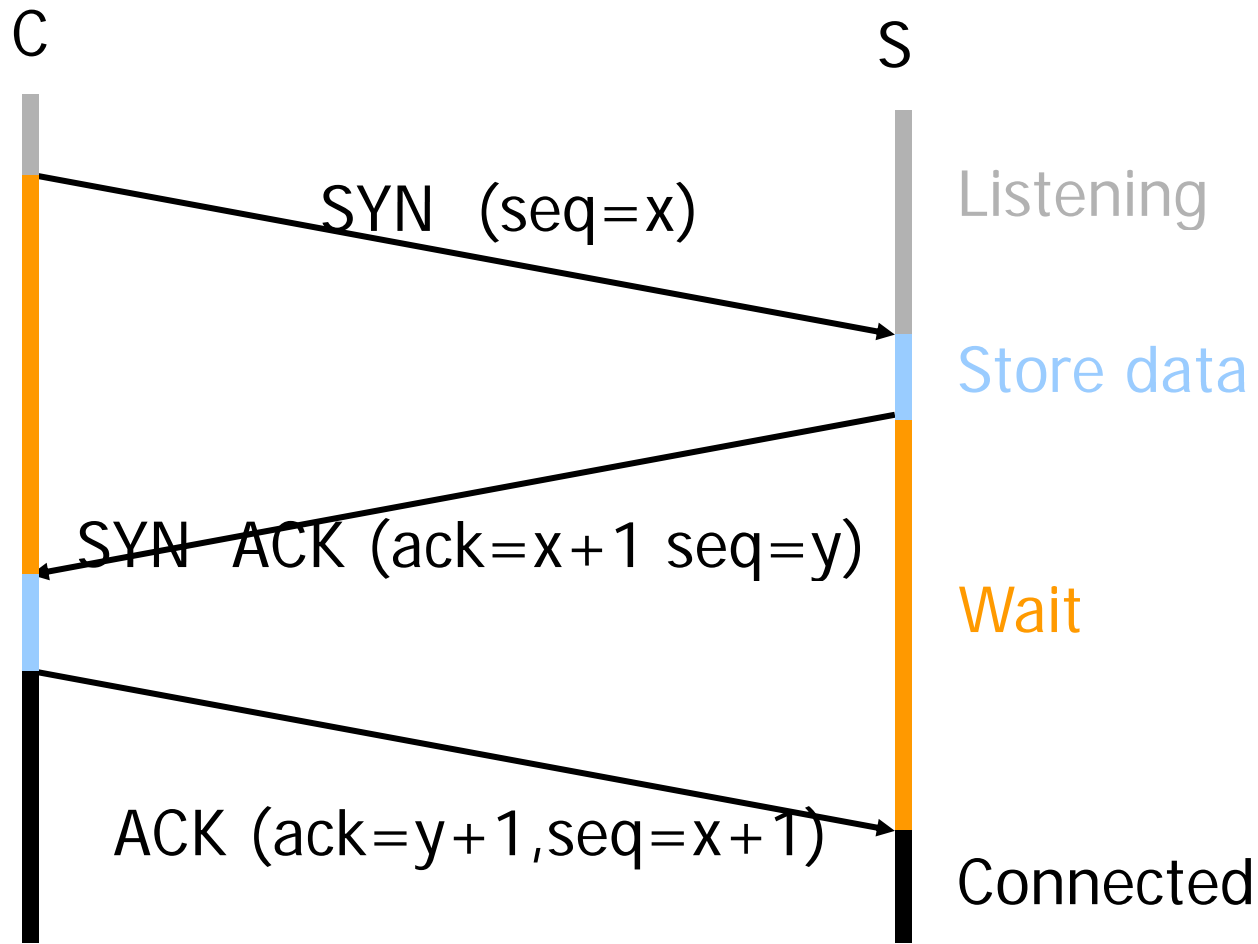
- Connection-oriented, preserves order
 - Sender
 - Break data into packets
 - Attach sequence numbers
 - Receiver
 - Acknowledge receipt; lost packets are resent
 - Reassemble packets in correct order



TCP Sequence Numbers

- Sequence number (32 bits) – has a dual role:
 - If the SYN flag is set, then this is the initial sequence number. The sequence number of the actual first data byte is this sequence number plus 1.
 - If the SYN flag is clear, then this is the accumulated sequence number of the first data byte of this packet for the current session.
- Acknowledgment number (32 bits) –
 - If the ACK flag is set then this the next sequence number that the receiver is expecting.
 - This acknowledges receipt of all prior bytes (if any).

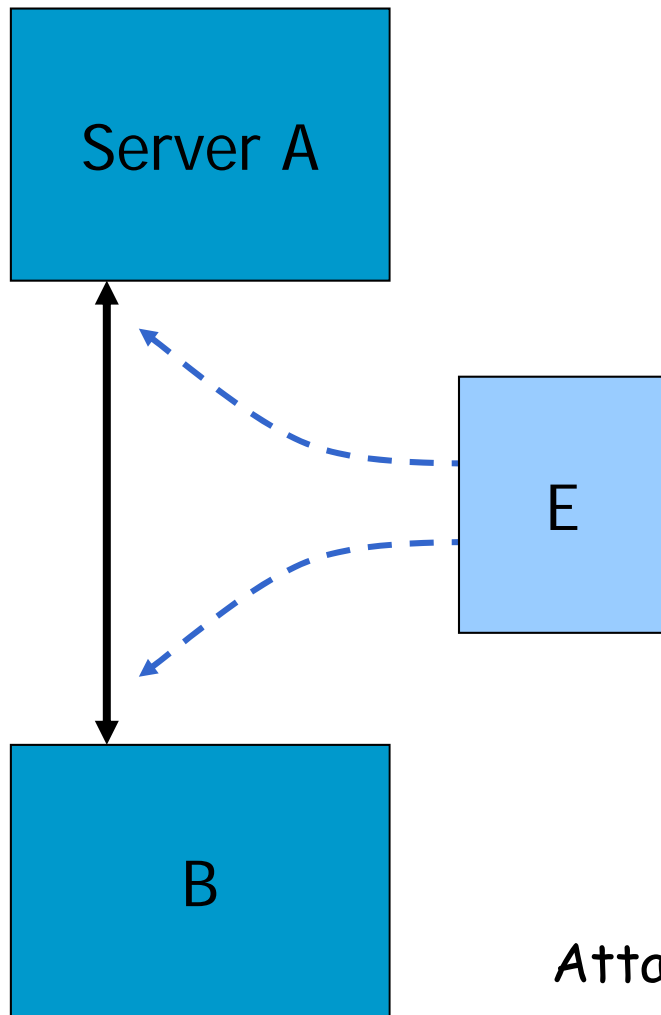
TCP Handshake



TCP sequence prediction attack

- Predict the sequence number used to identify the packets in a TCP connection, and then counterfeit packets.
- Adversary: do not have full control over the network, but can inject packets with fake source IP addresses
 - E.g., control a computer on the local network
- TCP sequence numbers are used for authenticating packets
- Initial seq# needs high degree of unpredictability
 - If attacker knows initial seq # and amount of traffic sent, can estimate likely current values
 - Some implementations are vulnerable

Blind TCP Session Hijacking



- A, B trusted connection
 - Send packets with predictable seq numbers
- E impersonates B to A
 - Opens connection to A to get initial seq number
 - DoS B's queue
 - Sends packets to A that resemble B's transmission
 - E cannot receive, but may execute commands on A

Attack can be blocked if E is outside firewall.

Risks from Session Hijacking

- Inject data into an unencrypted server-to-server traffic, such as an e-mail exchange, DNS zone transfers, etc.
- Inject data into an unencrypted client-to-server traffic, such as ftp file downloads, http responses.
- IP addresses often used for preliminary checks on firewalls or at the service level.
- Hide origin of malicious attacks.
- Carry out MITM attacks on weak cryptographic protocols.
 - often result in warnings to users that get ignored
- Denial of service attacks, such as resetting the connection.

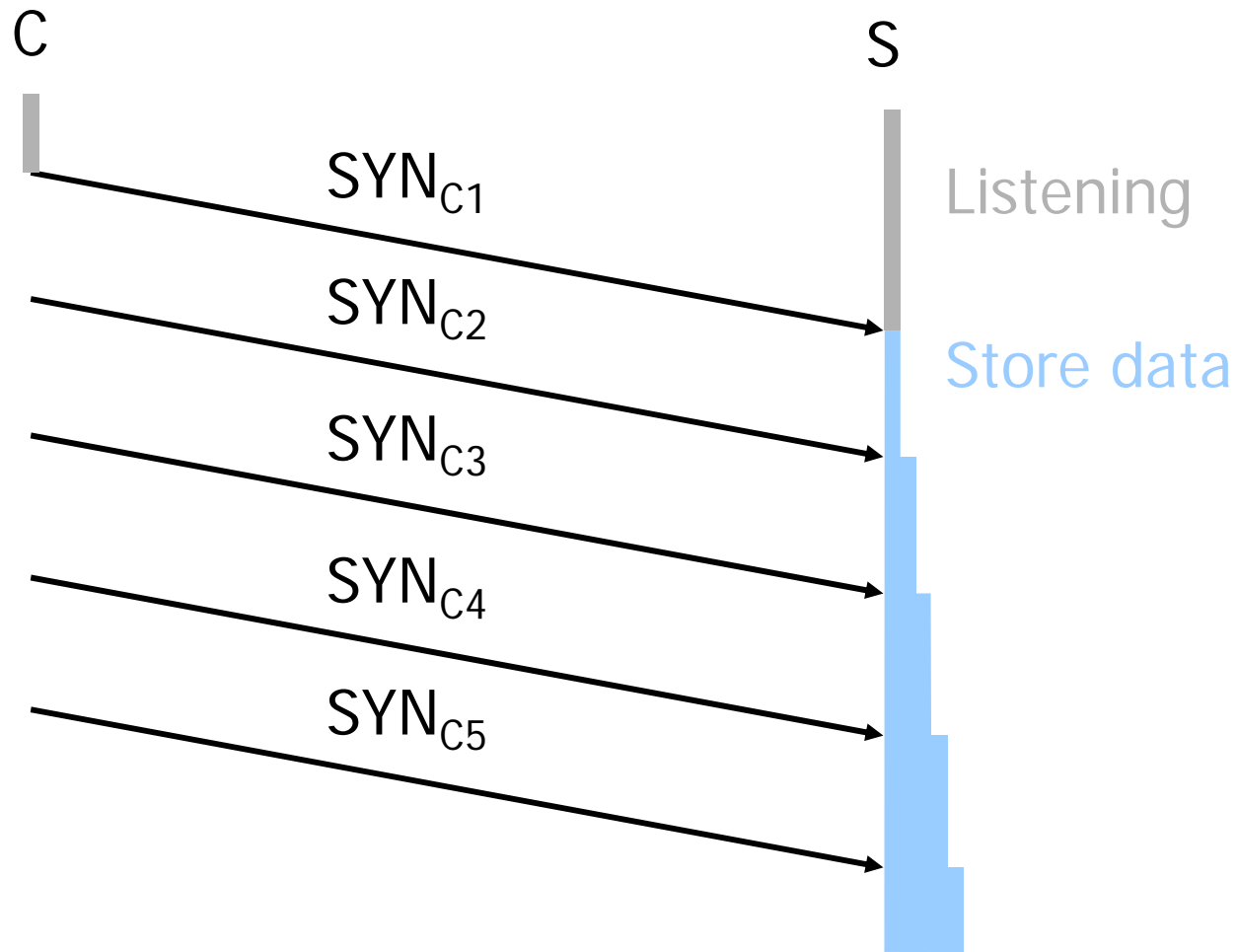
DoS vulnerability caused by session hijacking

- Suppose attacker can guess seq. number for an existing connection:
 - Attacker can send Reset packet to close connection. Results in DoS.
 - Naively, success prob. is $1/2^{32}$ (32-bit seq. #'s).
 - Most systems allow for a large window of acceptable seq. #'s
 - Much higher success probability.
- Attack is most effective against long lived connections, e.g. BGP.

Categories of Denial-of-service Attacks

	Stopping services	Exhausting resources
Locally	<ul style="list-style-type: none">• Process killing• Process crashing• System reconfiguration	<ul style="list-style-type: none">• Spawning processes to fill the process table• Filling up the whole file system• Saturate comm bandwidth
Remotely	<ul style="list-style-type: none">• Malformed packets to crash buggy services	<ul style="list-style-type: none">• Packet floods (Smurf, SYN flood, DDoS, etc)

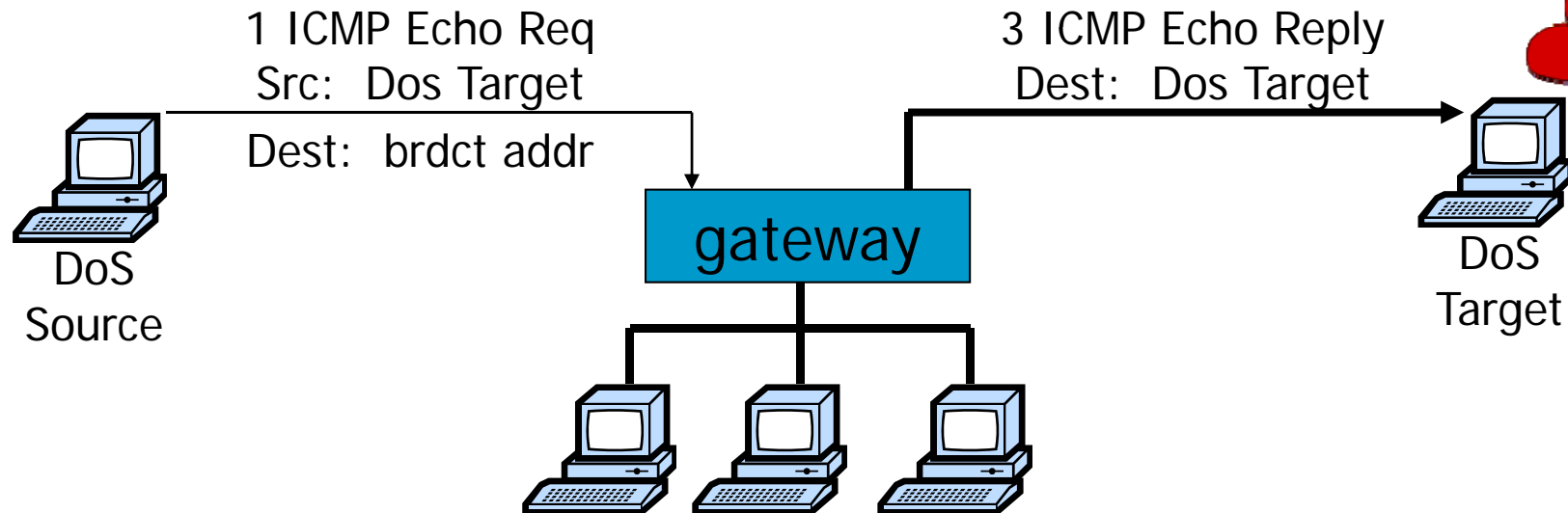
SYN Flooding



SYN Flooding

- Attacker sends many connection requests
 - Spoofed source addresses
- Victim allocates resources for each request
 - Connection requests exist until timeout
 - Old implementations have a small and fixed bound on half-open connections
- Resources exhausted \Rightarrow requests rejected
- No more effective than other channel capacity-based attack today

Smurf DoS Attack



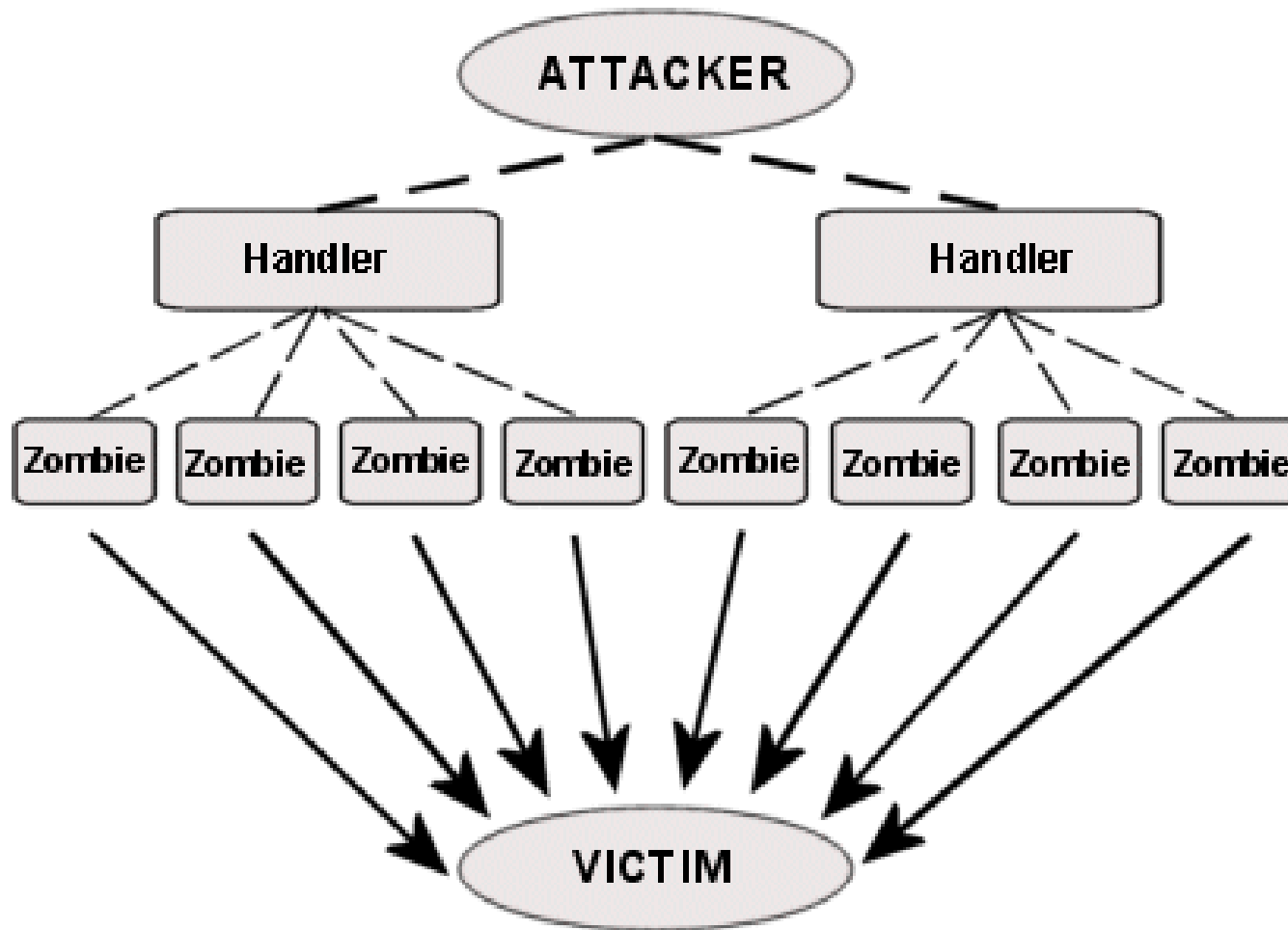
- Send ping request to broadcast addr (ICMP Echo Req)
 - Lots of responses:
 - Every host on target network generates a ping reply (ICMP Echo Reply) to victim
 - Ping reply stream can overload victim
- Prevention: reject external packets to broadcast address

Internet Control Message Protocol

- Provides feedback about network operation
 - Error reporting
 - Reachability testing
 - Congestion Control
- Example message types
 - Destination unreachable
 - Time-to-live exceeded
 - Parameter problem
 - Redirect to better gateway
 - Echo/echo reply - reachability test
 - Timestamp request/reply - measure transit delay

Distributed DoS (DDoS)

Architecture of a DDoS Attack



Hiding DDoS Attacks

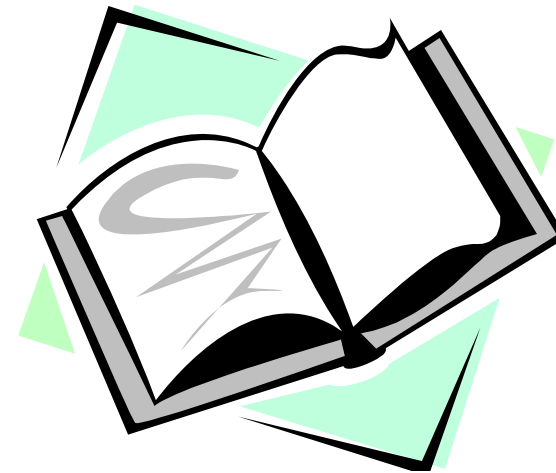
- Reflection
 - Find big sites with lots of resources, send packets with spoofed source address, response to victim
 - PING => PING response
 - SYN => SYN-ACK
- Pulsing zombie floods
 - each zombie active briefly, then goes dormant;
 - zombies taking turns attacking
 - making tracing difficult

Cryptographic network protection

- Solutions above the transport layer
 - Examples: SSL and SSH
 - Protect against session hijacking and injected data
 - Do not protect against denial-of-service attacks caused by spoofed packets
- Solutions at network layer
 - Use cryptographically random ISNs [RFC 1948]
 - More generally: IPsec
 - Can protect against
 - session hijacking and injection of data
 - denial-of-service attacks using session resets

Readings for This Lecture

- Optional Reading
 - [Steve Bellovin: A Look Back at “Security Problems in the TCP/IP Protocol Suite”](#)



Coming Attractions ...

- DNS Security

