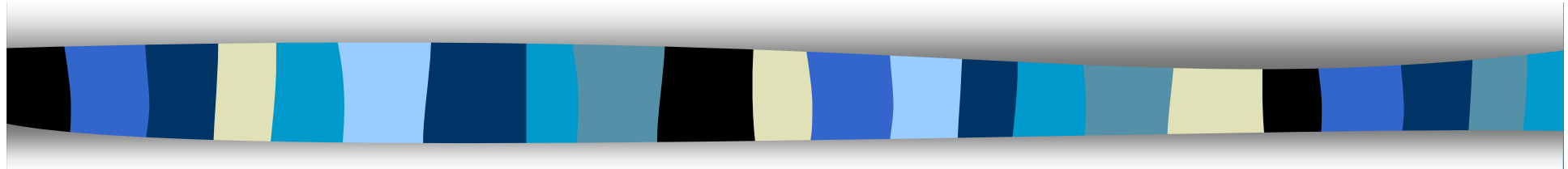


Computer Security

CS 426

Lecture 8



User Authentication

User Authentication

- Using a method to validate users who attempt to access a computer system or resources, to ensure they are authorized
- Types of user authentication
 - Something you know
 - User accounts with passwords
 - Something you have
 - Smart cards or other security tokens
 - Something you are
 - Biometrics

Scenarios Requiring Authentication

- Scenarios
 - Logging into a local computer
 - Logging into a computer remotely
 - Access web sites
- Potential vulnerabilities to consider when client authenticating server
 - channel between the client and the server
 - server compromise
 - client compromise
 - social engineering
 - weak passwords
- **What are some threats?**

Threats to Passwords

- Online guessing attempts
- Offline dictionary attacks
- Login spoofing
- Shoulder surfing
- Social engineering
 - e.g., pretexting: creating and using an invented scenario (the pretext) to persuade a target to release information or perform an action and is usually done over the telephone

Storing Passwords (UNIX Case Study)

- Old UNIX
 - The file `/etc/passwd` stores $H(\text{password})$ together with each user's login name, user id, home directory, login shell, etc.
 - file must be world readable
 - Brute force attacks possible even if H is one-way
 - how to brute-force when trying to obtain password of any account on a system with many accounts?
- New UNIX
 - $H(\text{password})$ stored in `/etc/shadow`, readable only by root

Password Salts

- Store $[r, H(\text{password}, r)]$ rather than $H(\text{password})$
 - r is randomly chosen for each password
 - r is public
 - similar to Initial Vector in CBC & CTR modes
- Benefits
 - dictionary attacks much more difficult
 - cost of attacking a single account remains the same
 - if two users happen to choose the same password, it doesn't immediately show

Mechanisms to Defend Against Dictionary and Guessing Attacks

- Protect stored passwords (use both cryptography & access control)
- Disable accounts with multiple failed attempts

Mechanisms to Avoid Weak Passwords

- Allow long passphrases
- Randomly generate passwords
- Check the quality of user-selected passwords
 - use a number of rules
 - run dictionary attack tools
- Give user suggestions/guidelines in choosing passwords
 - e.g., think of a sentence and select letters from it, “It’s 12 noon and I am hungry” => “I’S12&IAH”
 - Using both letter, numbers, and special characters
- Mandate password expiration
- Things to remember: Usability issues

Mechanisms to Defend Against Login Spoofing: Trusted Path

- Attacks:
 - write a program showing a login window on screen and record the passwords
 - put su in current directory
- Defense: Trusted Path
 - Mechanism that provides confidence that the user is communicating with what the real server (typically Trusted Computing Base in OSes)
 - attackers can't intercept or modify whatever information is being communicated.
 - defends attacks such as fake login programs
 - Example: Ctrl+Alt+Del for log in on Windows

Defending Against Other Threats

- Use ideas from recent research:
 - graphical passwords,
 - combine with typing
- Go beyond passwords
 - security tokens
 - biometrics
 - 2-factor authentication
 - US Banks are required to use 2-factor authentication by end of 2006 for online banking

Using Passwords Over Insecure Channel

- One-time passwords
 - Each password is used only once
 - Defend against passive adversaries who eavesdrop and later attempt to impersonate
- Challenge response
 - Send a response related to both the password and a challenge
- Zero knowledge proof of knowledge

How to do One-Time Password

- Time-synchronized OTP
- Shared lists of one-time passwords
- Using a hash chain (Lamport)
 - $h(s), h(h(s)), h(h(h(s))), \dots, h^{1000}(s)$
 - use these values as passwords in reverse order



Lamport's One-Time Password: Details

- One-time setup:
 - A selects a value w , a hash function $H()$, and an integer t , computes $w_0 = H^t(w)$ and sends w_0 to B
 - B stores w_0
- Protocol: to identify to B for the i^{th} time, $1 \leq i \leq t$
 - A sends to B: $A, i, w_i = H^{t-i}(w)$
 - B checks $i = i_A, H(w_i) = w_{i-1}$
 - if both holds, $i_A = i_A + 1$

Challenge-Response Protocols

- Goal: one entity authenticates to other entity proving the knowledge of a secret, ‘challenge’
- Approach: Use time-variant parameters to prevent replay, interleaving attacks, provide uniqueness and timeliness
 - e.g., nonce (used only once), timestamps

Challenge-response based on symmetric-key crypto

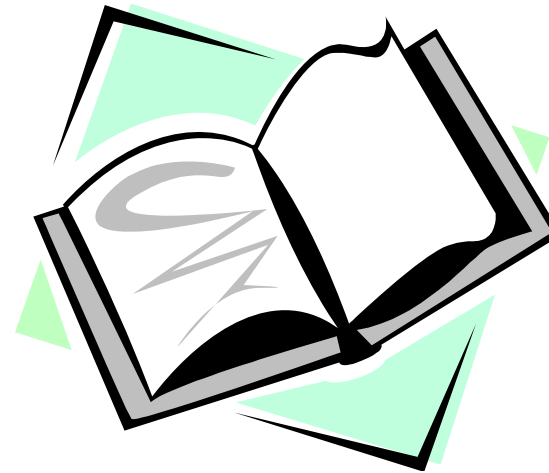
- Unilateral authentication, timestamp-based
 - A to B: $\text{MAC}_K(t_A, B)$
- Unilateral authentication, nonce-based
 - B to A: r_B
 - A to B: $\text{MAC}_K(r_B, B)$
- Mutual authentication, nonce-based
 - B to A: r_B
 - A to B: $r_A, \text{MAC}_K(r_A, r_B, B)$
 - B to A: $\text{MAC}_K(r_B, r_A)$

Issues to Consider in Password Systems

- Which types of attacks to defend against?
 - targeted attack on one account
 - attempt to penetrate any account on a system
 - attempt to penetrate any account on any system
 - service denial attack
- Whether to protect users against each other?
- Can users be trained? Will they follow the suggestions?
- Will the passwords be used in other systems?
- Whether the passwords will be used in a controlled environment?

Readings for This Lecture

- Wikipedia
 - [Salt_\(cryptography\)](#)
 - [Password cracking](#)
 - [Trusted path](#)
 - [One time password](#)



Coming Attractions ...

- Unix Access Control

