CS 580: Algorithm Design and Analysis

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Spring 2019

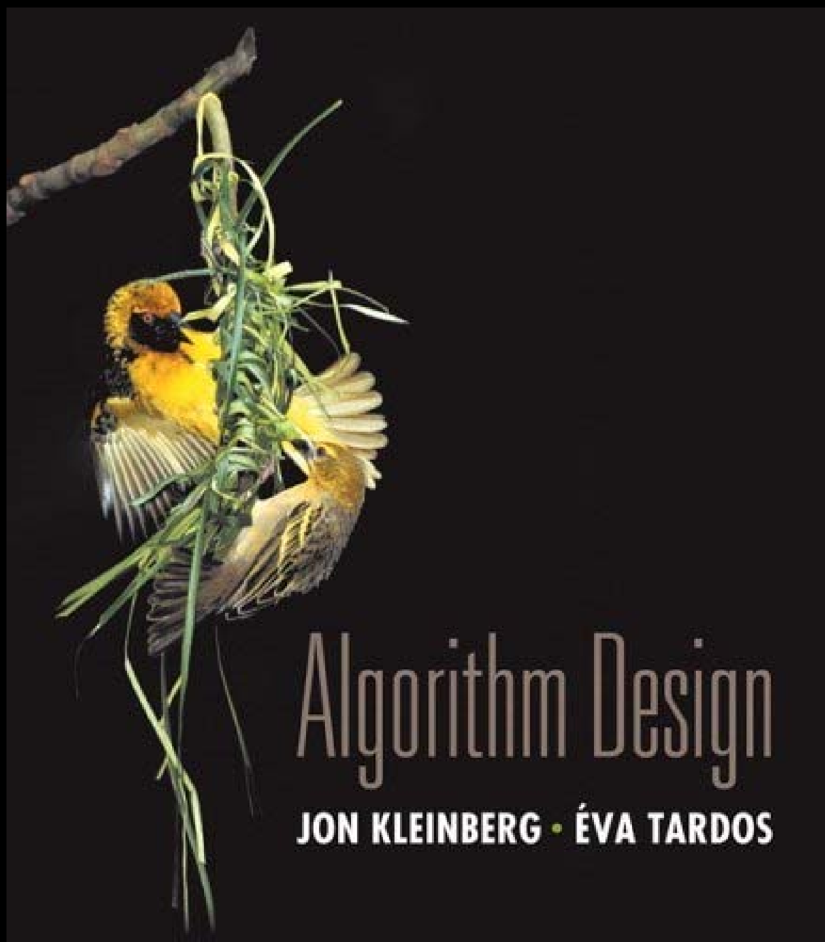
Midterm 2. April 3 @ 8PM (EE 170) [Tomorrow Night!]
Practice Midterm Released (+ Solutions)
3x5 Index Card (Double Sided)
Class Canceled: Thursday (April 4)

Midterm 2

- When?
  - April 3rd from 8PM to 10PM (2 hours)
- Where?
  - EE 170
- What can I bring?
  - 3x5 inch index card with your notes (double sided)
  - No electronics (phones, computers, calculators etc…)

Midterm 2

- When?
  - April 3rd from 8PM to 10PM (2 hours)
- Where?
  - EE 170
- What material should I study?
  - The midterm will cover recent topics more heavily
    - Network Flow
      - Max-Flow Min-Cut, Augmenting Paths, etc…
      - Ford Fulkerson, Dinic's Algorithm etc…
      - Applications of Network Flow (e.g., Maximum Bipartite Matching)
    - Linear Programming
    - NP-Completeness
      - Polynomial time reductions, P, NP, NP-Hard, NP-Completess, coNP
    - PSPACE (minimal coverage; some basic questions)



Chapter 9

PSPACE: A Class of Problems Beyond NP

Slides by Kevin Wayne.

Geography Game

Geography. Alice names capital city c of country she is in. Bob names a capital city c' that starts with the letter on which c ends. Alice and Bob repeat this game until one player is unable to continue. Does Alice have a forced win?

Ex. Budapest → Tokyo → Ottawa → Ankara → Amsterdam → Moscow → Washington → Nairobi → …

Geography on graphs. Given a directed graph $G = (V, E)$ and a start node s, two players alternate turns by following, if possible, an edge out of the current node to an unvisited node. Can first player guarantee to make the last legal move?

Remark. Some problems (especially involving 2-player games and AI) defy classification according to P, EXPTIME, NP, and NP-complete.

9.1 PSPACE

PSPACE

P. Decision problems solvable in polynomial time.

PSPACE. Decision problems solvable in polynomial space.

Observation. P $\subseteq$ PSPACE.

poly-time algorithm can consume only polynomial space

PSPACE

Binary counter. Count from 0 to $2^n - 1$ in binary.
Algorithm. Use n bit odometer.

Claim. 3-SAT is in PSPACE.
Pf.

- Enumerate all 2^n possible truth assignments using counter.
- Check each assignment to see if it satisfies all clauses. ▪

Theorem. NP $\subseteq$ PSPACE.
Pf. Consider arbitrary problem Y in NP.

- Since Y $\leq_P$ 3-SAT, there exists algorithm that solves Y in poly-time plus polynomial number of calls to 3-SAT black box.
- Can implement black box in poly-space. ▪

9.3 Quantified Satisfiability

Quantified Satisfiability

QSAT. Let $\Phi(x_1, \ldots, x_n)$ be a Boolean CNF formula. Is the following propositional formula true?

$$\exists x_1 \; \forall x_2 \; \exists x_3 \; \forall x_4 \ldots \; \forall x_{n-1} \; \exists x_n \; \Phi(x_1, \ldots, x_n)$$

↑ assume n is odd

Intuition. Amy picks truth value for x_1, then Bob for x_2, then Amy for x_3, and so on. Can Amy satisfy Φ no matter what Bob does?

Ex. $(x_1 \vee x_2) \wedge (x_2 \vee \overline{x_3}) \wedge (\overline{x_1} \vee \overline{x_2} \vee x_3)$

Yes. Amy sets x_1 true; Bob sets x_2; Amy sets x_3 to be same as x_2.

Ex. $(x_1 \vee x_2) \wedge (\overline{x_2} \vee \overline{x_3}) \wedge (\overline{x_1} \vee \overline{x_2} \vee x_3)$

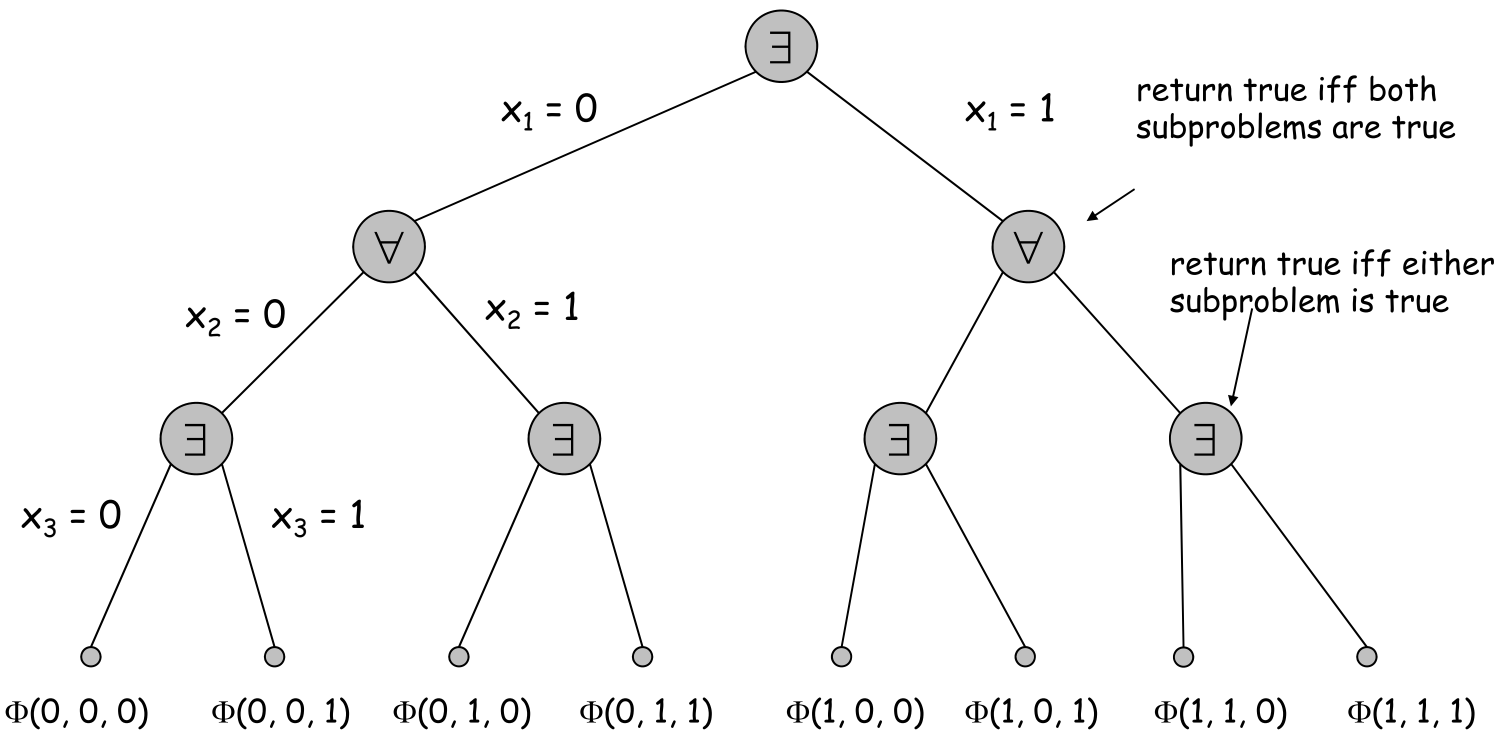
No. If Amy sets x_1 false; Bob sets x_2 false; Amy loses;
if Amy sets x_1 true; Bob sets x_2 true; Amy loses.

QSAT is in PSPACE

Theorem. QSAT $\in$ PSPACE.

Pf. Recursively try all possibilities.

- Only need one bit of information from each subproblem.
- Amount of space is proportional to depth of function call stack.

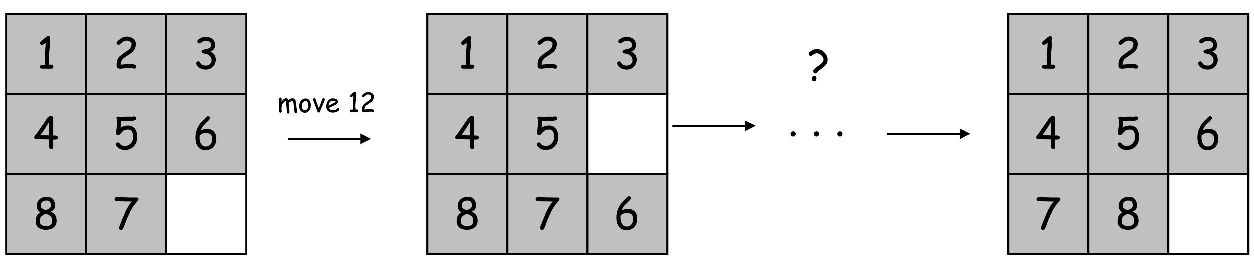


9.4 Planning Problem

15-Puzzle

8-puzzle, 15-puzzle. [Sam Loyd 1870s]

- Board: 3-by-3 grid of tiles labeled 1-8.
- Legal move: slide neighboring tile into blank (white) square.
- Find sequence of legal moves to transform initial configuration into goal configuration.



initial configuration

goal configuration

Planning Problem

Conditions. Set $C = \{ C_1, \ldots, C_n \}$.
Initial configuration. Subset $c_0 \subseteq C$ of conditions initially satisfied.
Goal configuration. Subset $c^* \subseteq C$ of conditions we seek to satisfy.
Operators. Set $O = \{ O_1, \ldots, O_k \}$.

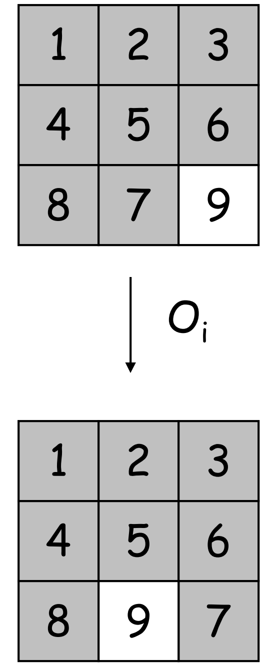
- To invoke operator O_i, must satisfy certain prereq conditions.
- After invoking O_i certain conditions become true, and certain conditions become false.

PLANNING. Is it possible to apply sequence of operators to get from initial configuration to goal configuration?

Examples.

- 15-puzzle.
- Rubik's cube.
- Logistical operations to move people, equipment, and materials.

Planning Problem: 8-Puzzle



Planning example. Can we solve the 8-puzzle?

Conditions. C_{ij}, $1 \le i, j \le 9$. ← C_{ij} means tile i is in square j

Initial state. $c_0 = \{C_{11}, C_{22}, \ldots, C_{66}, C_{78}, C_{87}, C_{99}\}$.

Goal state. $c^* = \{C_{11}, C_{22}, \ldots, C_{66}, C_{77}, C_{88}, C_{99}\}$.

Operators.

- Precondition to apply $O_i = \{C_{11}, C_{22}, \ldots, C_{66}, C_{78}, C_{87}, C_{99}\}$.
- After invoking O_i, conditions C_{79} and C_{97} become true.
- After invoking O_i, conditions C_{78} and C_{99} become false.

Solution. No solution to 8-puzzle or 15-puzzle!

Diversion: Why is 8-Puzzle Unsolvable?

8-puzzle invariant. Any legal move preserves the parity of the number of pairs of pieces in reverse order (inversions).

3	1	2
4	5	6
8	7	

3 inversions
1-3, 2-3, 7-8

$\longrightarrow$

3	1	2
4	5	6
8		7

3 inversions
1-3, 2-3, 7-8

$\longrightarrow$

3	1	2
4		6
8	5	7

5 inversions
1-3, 2-3, 7-8, 5-8, 5-6

1	2	3
4	5	6
7	8	

0 inversions

$\not\longrightarrow$

1	2	3
4	5	6
8	7	

1 inversion: 7-8

Planning Problem: Binary Counter

Planning example. Can we increment an n-bit counter from the all-zeroes state to the all-ones state?

Conditions. $C_1, \ldots, C_n$. ← C_i corresponds to bit i = 1

Initial state. $c_0 = \phi$. ← all 0s

Goal state. $c^* = \{C_1, \ldots, C_n\}$. ← all 1s

Operators. $O_1, \ldots, O_n$.

- To invoke operator O_i, must satisfy $C_1, \ldots, C_{i-1}$. ← i-1 least significant bits are 1
- After invoking O_i, condition C_i becomes true. ← set bit i to 1
- After invoking O_i, conditions $C_1, \ldots, C_{i-1}$ become false.

set i-1 least significant bits to 0

Solution. $\{\,\} \Rightarrow \{C_1\} \Rightarrow \{C_2\} \Rightarrow \{C_1, C_2\} \Rightarrow \{C_3\} \Rightarrow \{C_3, C_1\} \Rightarrow \ldots$

Observation. Any solution requires at least $2^n - 1$ steps.

Planning Problem: In Exponential Space

Configuration graph G.

- Include node for each of 2^n possible configurations.
- Include an edge from configuration c' to configuration c'' if one of the operators can convert from c' to c''.

PLANNING. Is there a path from c_0 to c^* in configuration graph?

Claim. PLANNING is in EXPTIME.
Pf. Run BFS to find path from c_0 to c^* in configuration graph. ▪

Note. Configuration graph can have 2^n nodes, and shortest path can be of length = $2^n - 1$.

↑
binary counter

Planning Problem: In Polynomial Space

Theorem. PLANNING is in PSPACE.

Pf.

- Suppose there is a path from c_1 to c_2 of length L.
- Path from c_1 to midpoint and from midpoint to c_2 are each $\leq L/2$.
- Enumerate all possible midpoints.
- Apply recursively. Depth of recursion = $\log_2 L$. ▪

```
boolean hasPath(c1, c2, L) {
   if (L ≤ 1) return correct answer

                enumerate using binary counter
   foreach configuration c' {
      boolean x = hasPath(c1, c', L/2)
      boolean y = hasPath(c', c2, L/2)
      if (x and y) return true
   }
   return false
}
```

9.5 PSPACE-Complete

PSPACE-Complete

PSPACE. Decision problems solvable in polynomial space.

PSPACE-Complete. Problem Y is PSPACE-complete if (i) Y is in PSPACE and (ii) for every problem X in PSPACE, $X \leq_P Y$.

Theorem. [Stockmeyer-Meyer 1973] QSAT is PSPACE-complete.

Theorem. PSPACE $\subseteq$ EXPTIME.
Pf. Previous algorithm solves QSAT in exponential time, and QSAT is PSPACE-complete. ▪

Summary. P $\subseteq$ NP $\subseteq$ PSPACE $\subseteq$ EXPTIME.

it is known that P $\neq$ EXPTIME, but unknown which inclusion is strict; conjectured that all are

PSPACE-Complete Problems

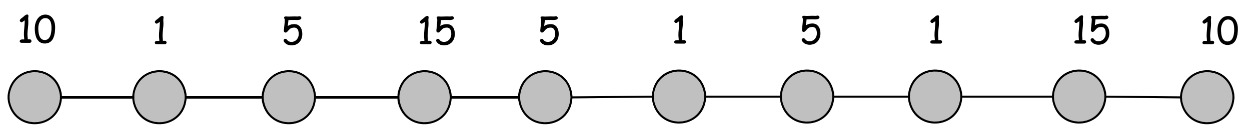
More PSPACE-complete problems.

- Competitive facility location.
- Natural generalizations of games.
  - Othello, Hex, Geography, Rush-Hour, Instant Insanity
  - Shanghai, go-moku, Sokoban
- Given a memory restricted Turing machine, does it terminate in at most k steps?
- Do two regular expressions describe different languages?
- Is it possible to move and rotate complicated object with attachments through an irregularly shaped corridor?
- Is a deadlock state possible within a system of communicating processors?

Competitive Facility Location

Input. Graph with positive edge weights, and target B.
Game. Two competing players alternate in selecting nodes. Not allowed to select a node if any of its neighbors has been selected.

Competitive facility location. Can second player guarantee at least B units of profit? (Player one might play vindictively to minimize B's profit)



Yes if B = 20; no if B = 25.

Competitive Facility Location

Claim. COMPETITIVE-FACILITY is PSPACE-complete.

Pf.

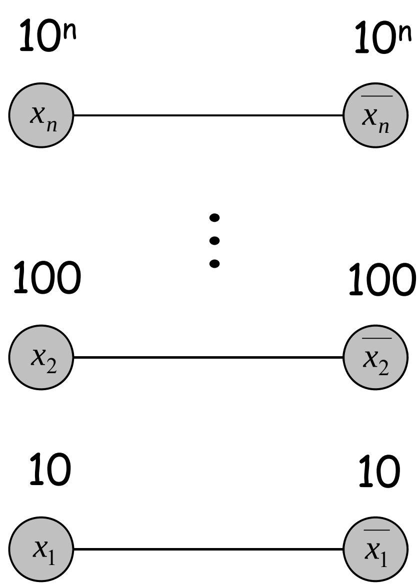
- To solve in poly-space, use recursion like QSAT, but at each step there are up to n choices instead of 2.
- To show that it's complete, we show that QSAT polynomial reduces to it. Given an instance of QSAT, we construct an instance of COMPETITIVE-FACILITY such that player 2 can force a win iff QSAT formula is false.

Competitive Facility Location

assume n is odd

Construction. Given instance $\Phi(x_1, \ldots, x_n) = C_1 \wedge C_1 \wedge \ldots C_k$ of QSAT.

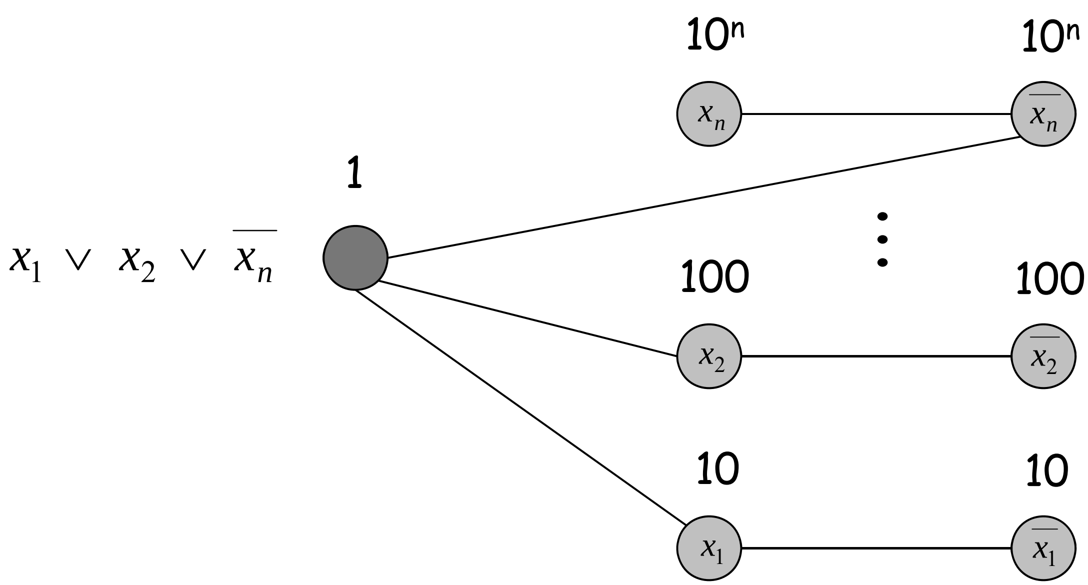
- Include a node for each literal and its negation and connect them.
  - at most one of x_i and its negation can be chosen
- Choose $c \geq k+2$, and put weight c^i on literal x^i and its negation; set $B = c^{n-1} + c^{n-3} + \ldots + c^4 + c^2 + 1$.
  - ensures variables are selected in order $x_n, x_{n-1}, \ldots, x_1$.
- As is, player 2 will lose by 1 unit: $c^{n-1} + c^{n-3} + \ldots + c^4 + c^2$.

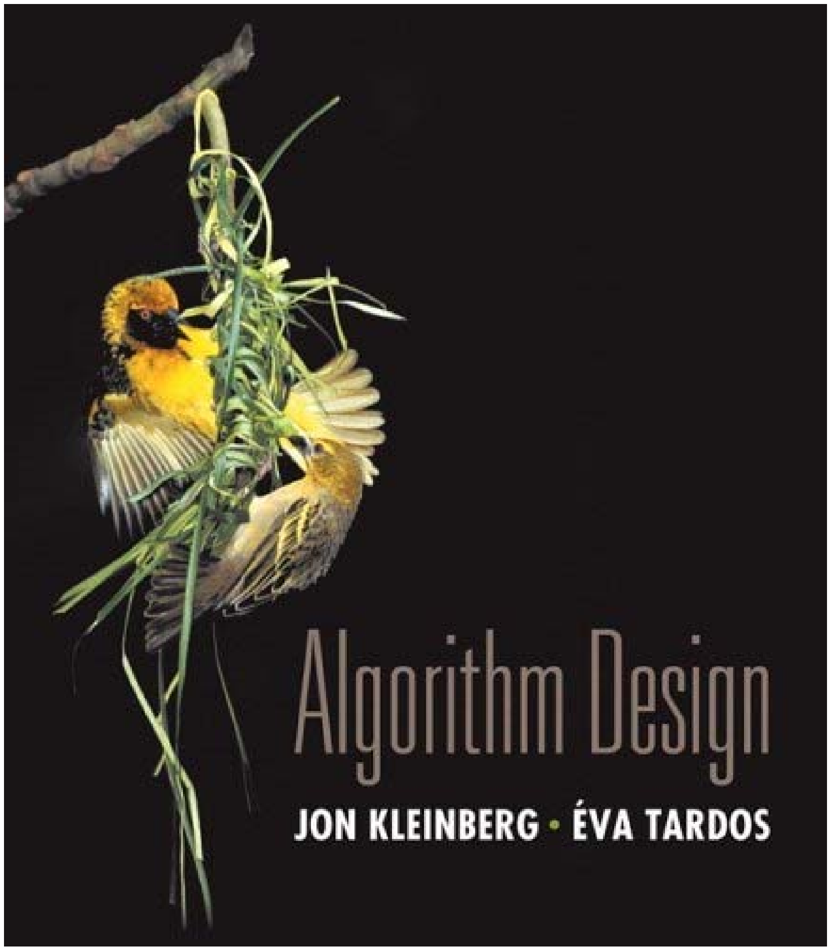


Competitive Facility Location

Construction. Given instance $\Phi(x_1, \ldots, x_n) = C_1 \wedge C_1 \wedge \ldots C_k$ of QSAT.

- Give player 2 one last move on which she can try to win.
- For each clause C_j, add node with value 1 and an edge to each of its literals.
- Player 2 can make last move iff truth assignment defined alternately by the players failed to satisfy some clause. ▪





Approximation Algorithms

Approximation Algorithms

Q. Suppose I need to solve an NP-hard problem. What should I do?
A. Theory says you're unlikely to find a poly-time algorithm.

Must sacrifice one of three desired features.

- Solve problem to optimality.
- Solve problem in poly-time.
- Solve arbitrary instances of the problem.

ρ-approximation algorithm.

- Guaranteed to run in poly-time.
- Guaranteed to solve arbitrary instance of the problem
- Guaranteed to find solution within ratio ρ of true optimum.

Challenge. Need to prove a solution's value is close to optimum, without even knowing what optimum value is!

11.1 Load Balancing

Load Balancing

Input. m identical machines; n jobs, job j has processing time t_j.

- Job j must run contiguously on one machine.
- A machine can process at most one job at a time.

Def. Let $J(i)$ be the subset of jobs assigned to machine i. The load of machine i is $L_i = \Sigma_{j \in J(i)} t_j$.

Def. The makespan is the maximum load on any machine $L = \max_i L_i$.

Load balancing. Assign each job to a machine to minimize makespan.

Load Balancing: List Scheduling

List-scheduling algorithm.

- Consider n jobs in some fixed order.
- Assign job j to machine whose load is smallest so far.

play

```
List-Scheduling(m, n, t_1,t_2,…,t_n) {
   for i = 1 to m {
      L_i ← 0          ← load on machine i
      J(i) ← φ         ← jobs assigned to machine i
   }

   for j = 1 to n {
      i = argmin_k L_k         ← machine i has smallest load
      J(i) ← J(i) ∪ {j}        ← assign job j to machine i
      L_i ← L_i + t_j          ← update load of machine i
   }
   return J(1), …, J(m)
}
```

Implementation. $O(n \log m)$ using a priority queue.

Load Balancing: List Scheduling Analysis

Theorem. [Graham, 1966] Greedy algorithm is a 2-approximation.

- First worst-case analysis of an approximation algorithm.
- Need to compare resulting solution with optimal makespan L*.

Lemma 1. The optimal makespan $L^* \geq \max_j t_j$.
Pf. Some machine must process the most time-consuming job. ▪

Lemma 2. The optimal makespan $L^* \geq \frac{1}{m} \sum_j t_j$
Pf.

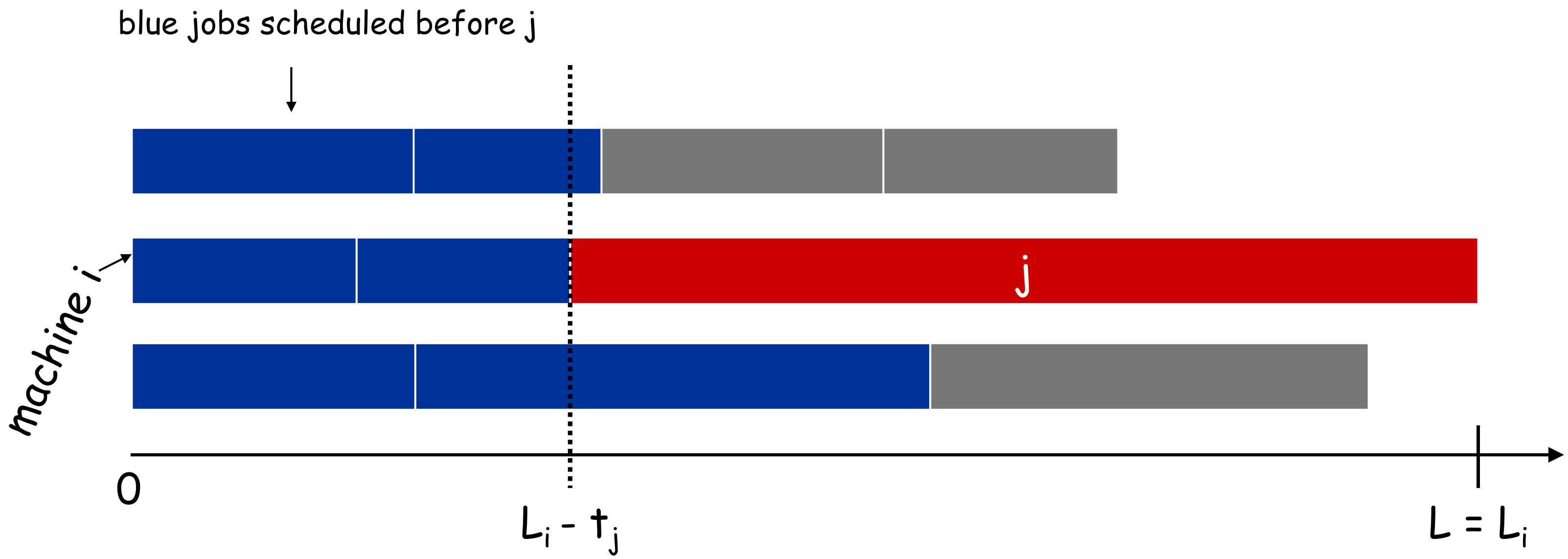
- The total processing time is $\Sigma_j t_j$.
- One of m machines must do at least a 1/m fraction of total work. ▪

Load Balancing: List Scheduling Analysis

Theorem. Greedy algorithm is a 2-approximation.

Pf. Consider load L_i of bottleneck machine i.

- Let j be last job scheduled on machine i.
- When job j assigned to machine i, i had smallest load. Its load before assignment is $L_i - t_j \Rightarrow L_i - t_j \le L_k$ for all $1 \le k \le m$.



Load Balancing: List Scheduling Analysis

Theorem. Greedy algorithm is a 2-approximation.

Pf. Consider load L_i of bottleneck machine i.

- Let j be last job scheduled on machine i.
- When job j assigned to machine i, i had smallest load. Its load before assignment is $L_i - t_j \;\Rightarrow\; L_i - t_j \le L_k$ for all $1 \le k \le m$.
- Sum inequalities over all k and divide by m:

$$L_i - t_j \le \frac{1}{m}\sum_{k=1}^{m} L_k$$

$$= \frac{1}{m}\sum_{k=1}^{n} t_k \quad \le L^*$$

Lemma 1

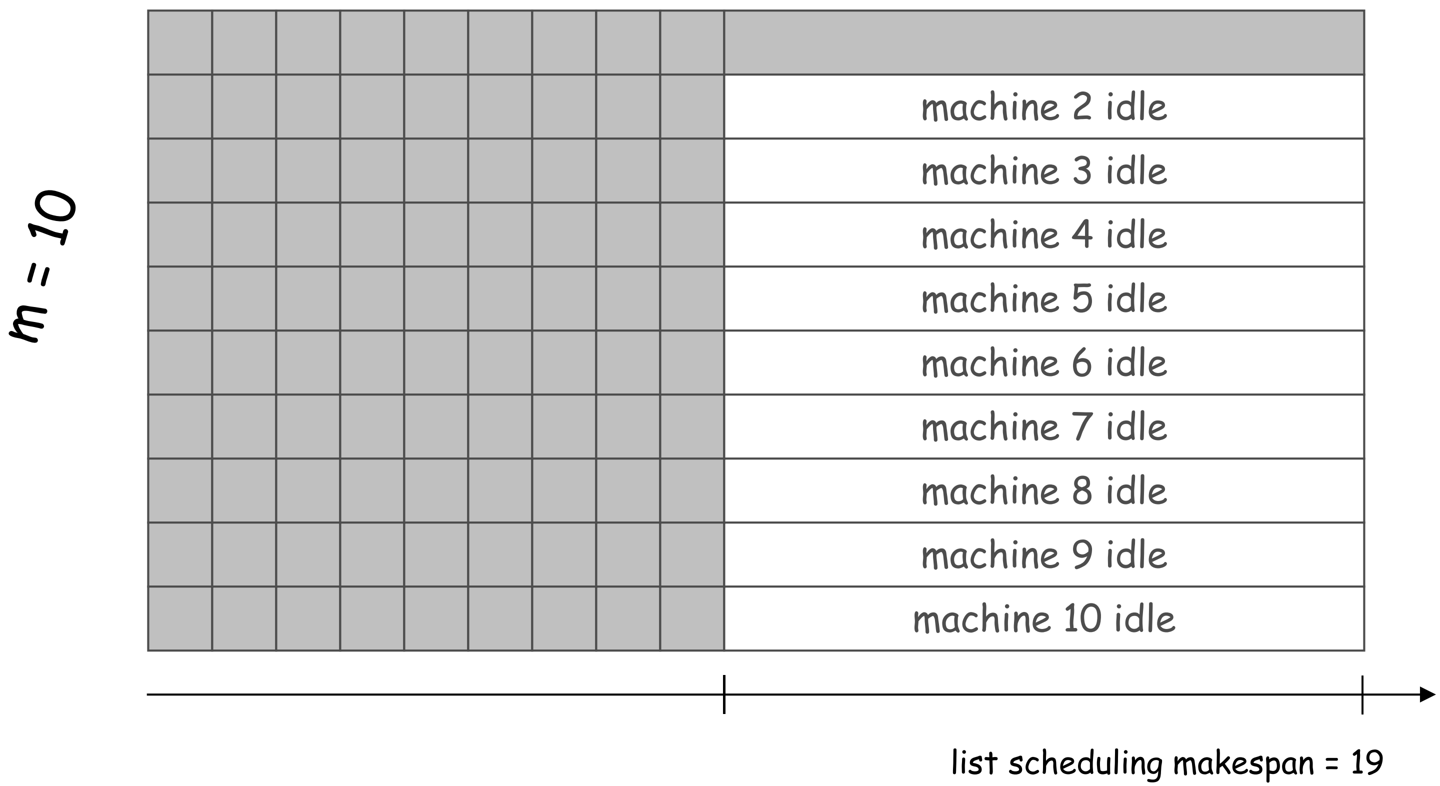
Now $L_i = \underbrace{(L_i - t_j)}_{\le L^*} + \underbrace{t_j}_{\le L^*} \le 2L^*$ ▪

Lemma 2

Load Balancing: List Scheduling Analysis

Q. Is our analysis tight?
A. Essentially yes.

Ex: m machines, m(m-1) jobs length 1 jobs, one job of length m



Load Balancing: List Scheduling Analysis

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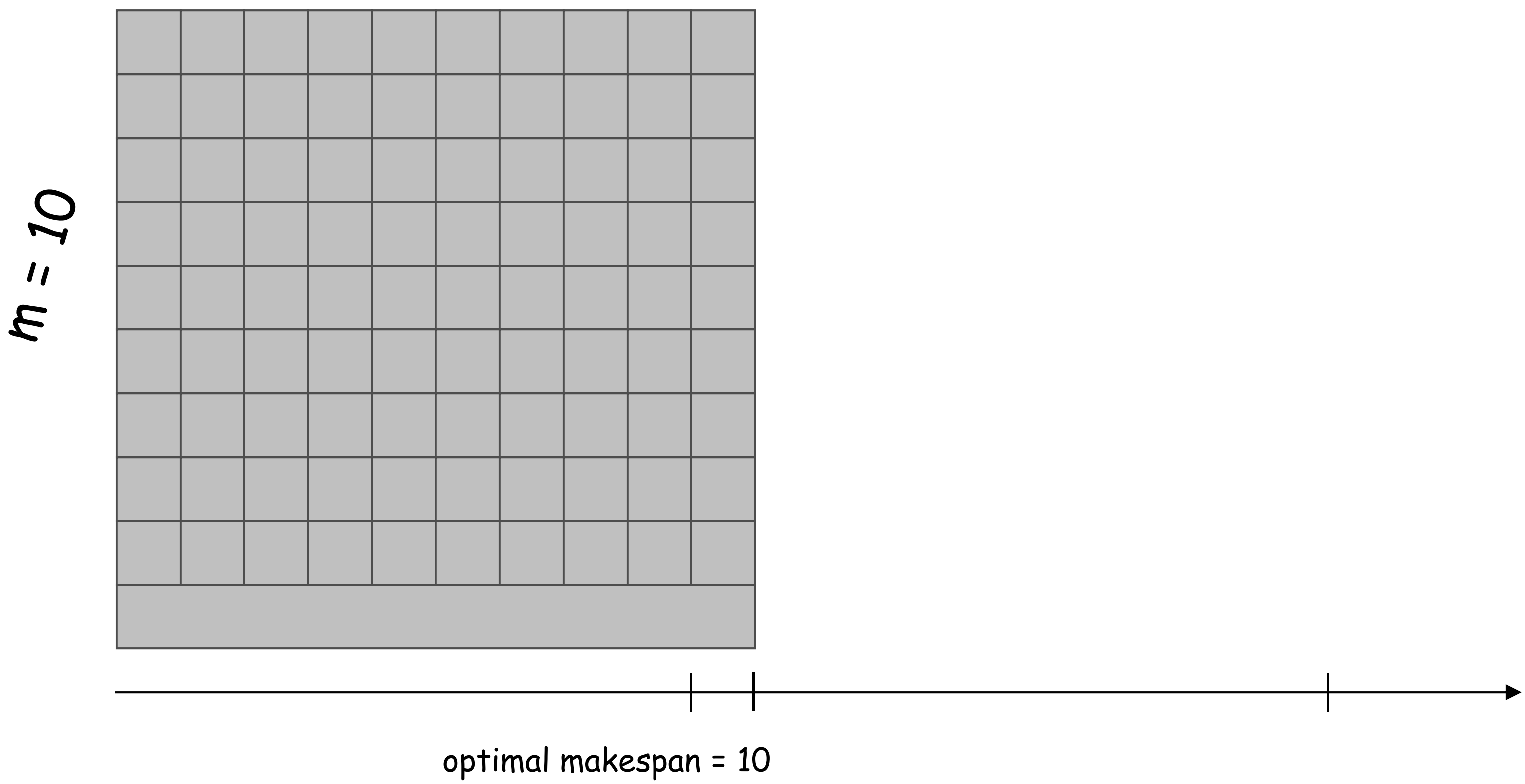

Load Balancing: LPT Rule

Longest processing time (LPT). Sort n jobs in descending order of processing time, and then run list scheduling algorithm.

```
LPT-List-Scheduling(m, n, t1,t2,…,tn) {
   Sort jobs so that t1 ≥ t2 ≥  … ≥ tn

   for i = 1 to m {
      Li ← 0           ← load on machine i
      J(i) ← φ         ← jobs assigned to machine i
   }

   for j = 1 to n {
      i = argmink Lk          ← machine i has smallest load
      J(i) ← J(i) ∪ {j}       ← assign job j to machine i
      Li ← Li + tj            ← update load of machine i
   }
   return J(1), …, J(m)
}
```

Load Balancing: LPT Rule

Observation. If at most m jobs, then list-scheduling is optimal.
Pf. Each job put on its own machine. ▪

Lemma 3. If there are more than m jobs, $L^* \geq 2\, t_{m+1}$.
Pf.
- Consider first m+1 jobs $t_1, \ldots, t_{m+1}$.
- Since the t_i's are in descending order, each takes at least t_{m+1} time.
- There are m+1 jobs and m machines, so by pigeonhole principle, at least one machine gets two jobs. ▪

Theorem. LPT rule is a 3/2 approximation algorithm.
Pf. Same basic approach as for list scheduling.

$$L_i \;=\; \underbrace{(L_i - t_j)}_{\leq\, L^*} \;+\; \underbrace{t_j}_{\leq\, \frac{1}{2} L^*} \;\leq\; \tfrac{3}{2} L^* . \;\blacksquare$$

Lemma 3
(by observation, can assume number of jobs > m)

Load Balancing: LPT Rule

Q. Is our 3/2 analysis tight?
A. No.

Theorem. [Graham, 1969] LPT rule is a 4/3-approximation.
Pf. More sophisticated analysis of same algorithm.

Q. Is Graham's 4/3 analysis tight?
A. Essentially yes.

Ex: m machines, $n = 2m+1$ jobs, 2 jobs of length $m+1$, $m+2$, …, $2m-1$ and one job of length m.