

CS 381 – FALL 2019

Week 5.3, Monday, Sept 16

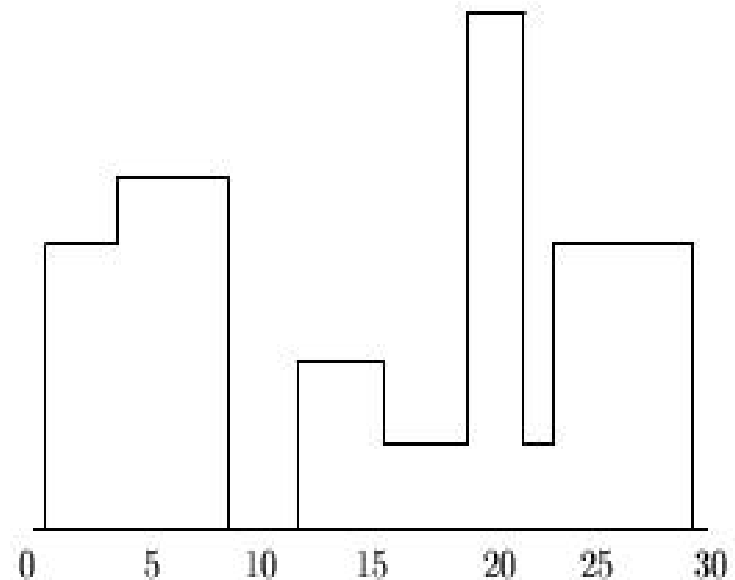
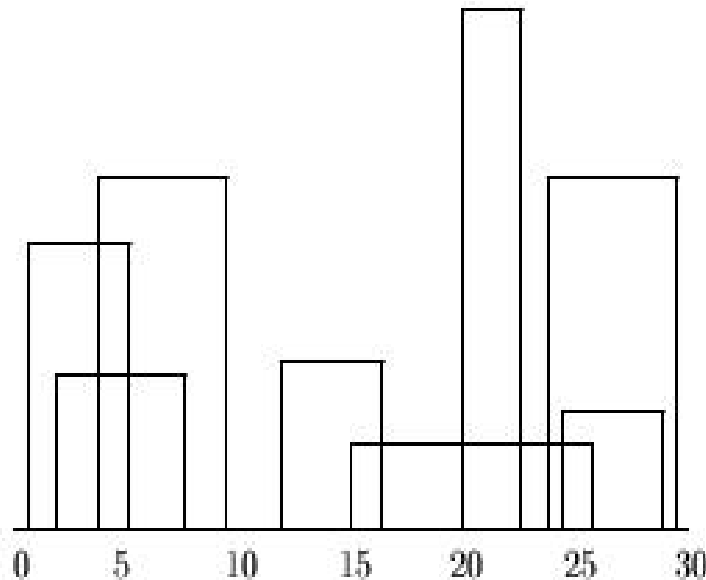
Guest Lecture (Prof. Hambrusch)
Homework 2 Due Tonight!
September 16th, 2019 @ 11:59PM on Gradescope

Homework 2 Reminders

- ▣ Must include collaborator/resource statement
 - No credit for solutions that don't include CR statement
- ▣ Must type solutions
 - Only allowed to scan hand drawn graphs/diagrams
 - No credit if the entire solution is a scan of your handwritten homework
 - Expectation to use math notation \sqrt{n} vs $n^{(1/2)}$
- ▣ Remember to select the appropriate pages for each problem on Gradescope

Skyline problem

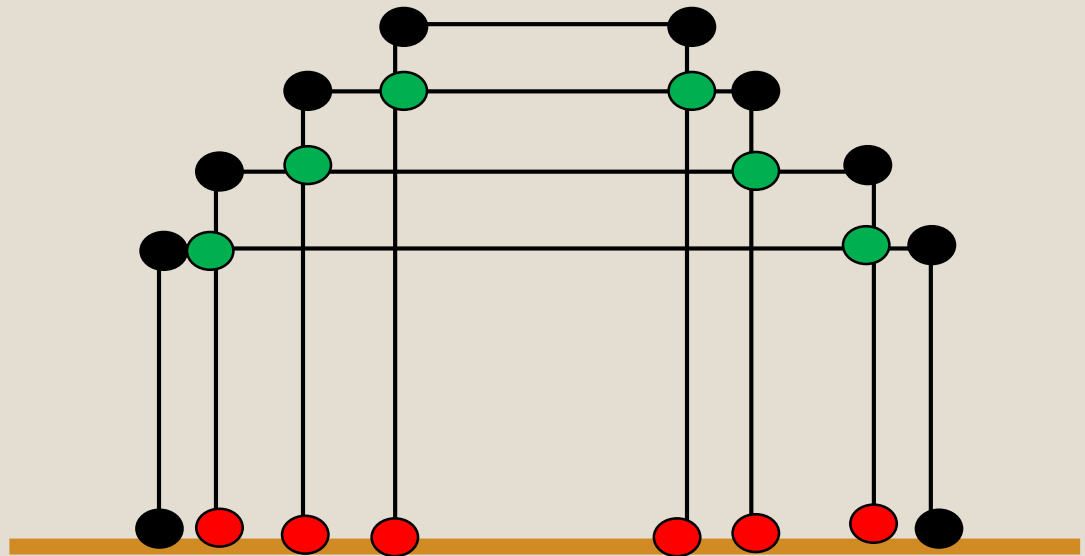
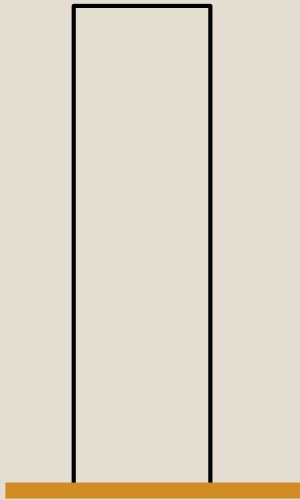
- Given are n “buildings” positioned on a common horizontal line. Building i is described by the triple (l_i, r_i, h_i) , $1 \leq i \leq n$.
- The skyline is the union of the outline formed by the n buildings. Represented by coordinates forming the corners.



Towards a first solution ...

How many corners can a skyline have?

- Minimum of 4 and maximum of $4n$ ($\#new \leq \#hidden$)



Formalizing the Problem...

Building i is described by the triple (l_i, r_i, h_i) , $1 \leq i \leq n$.

The skyline is represented by its corner coordinates

How many corners can a skyline have?

- Minimum of 4 and maximum of $4n$

How should one maintain the skyline?

- Store corners by non-decreasing x -coordinates

Formalizing the Problem...

Input: List of n triples (l_i, r_i, h_i) , $1 \leq i \leq n$ encoding the location of building i .

- Encodes:
 - Left corner $(l_i, 0)$ of building i
 - Right corner $(r_i, 0)$ of building i
 - Height h_i of building i

Output: List of (at most $4n$) corner points (x_i, y_i) sorted by x -coordinate i.e., $x_1 < x_2 < \dots$

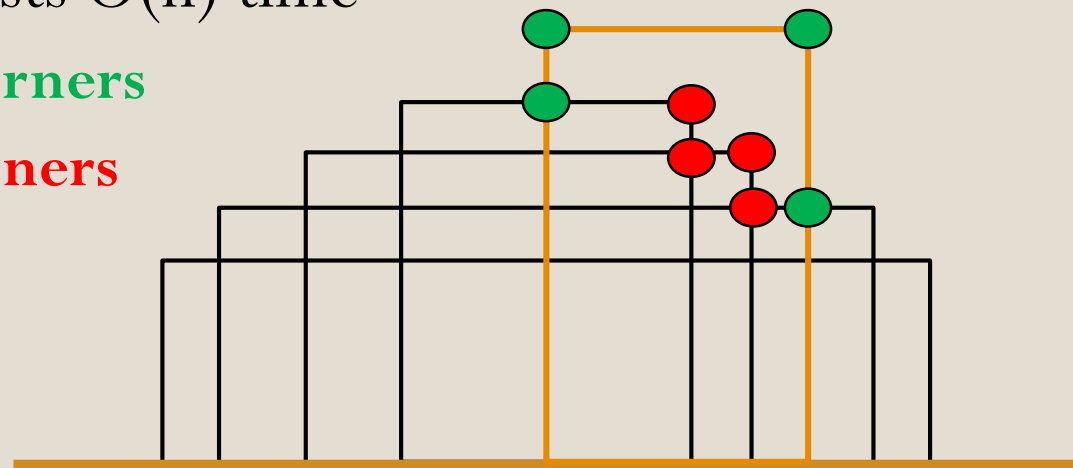
A first solution ...

The skyline is represented by its corner coordinates.

Store corners by non-decreasing x-coordinates

Assume we build the skyline by adding one building at a time

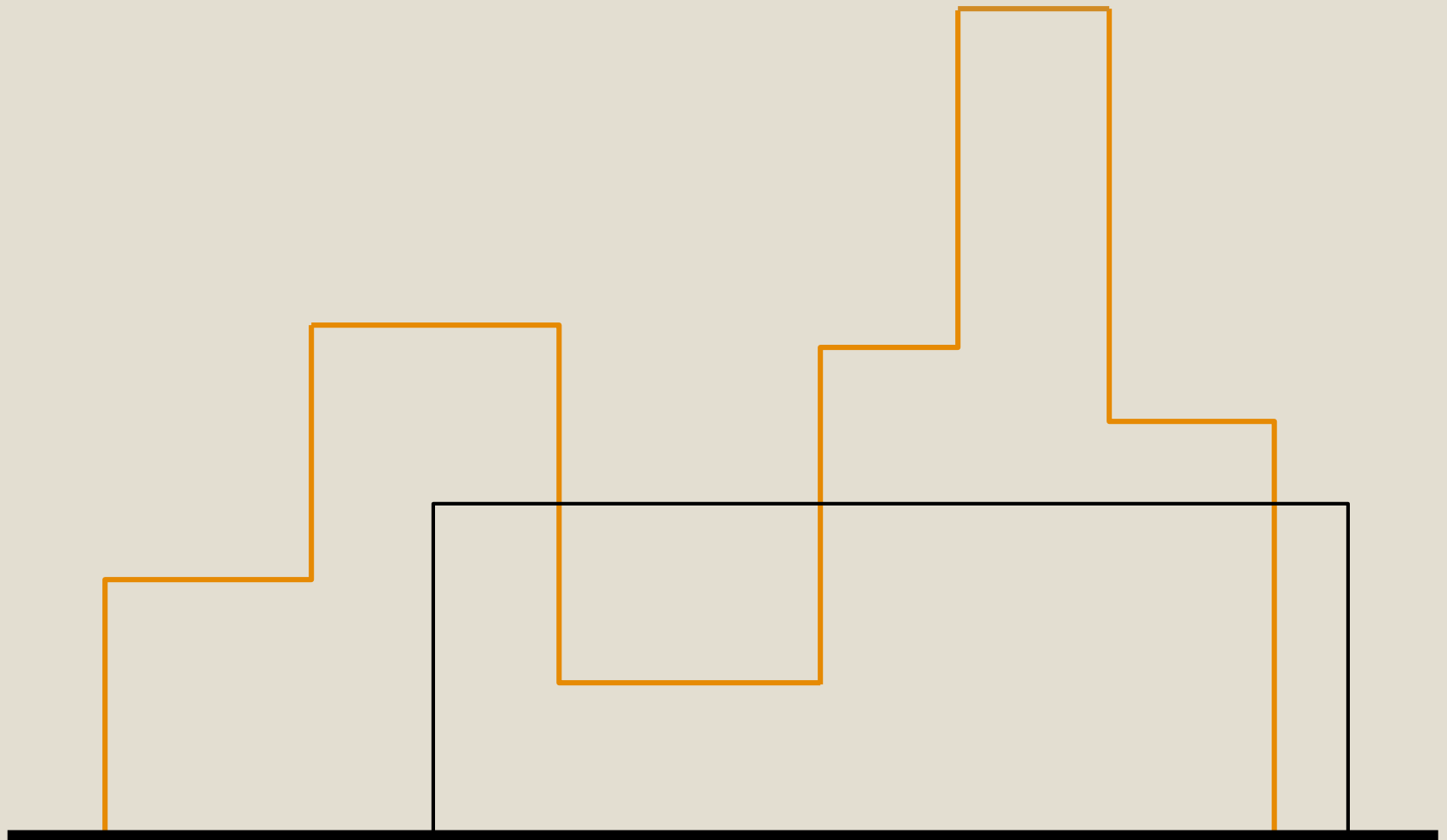
- Adding one building costs $O(n)$ time
 - Add (up to) 4 new corners
 - Remove (up to) 4 corners
- Yields $O(n^2)$ algorithm



Can we do better than $O(n^2)$?

Can we work with the sorted skyline more effectively?

- If we use binary search, we need to store the skyline in an array
- If we use an array, we need to move elements in the array after inserting into the current skyline
- Worst case seems to still be n^2
- Consider other data structures?
 - Balanced search trees?
 - Structures for searching in 2 dimensions?



A divide and conquer solution

1. Split the n buildings into two sets $B1$ and $B2$ of $n/2$ building each (no need to sort the buildings).
2. Find the skyline $S1$ for set $B1$.
3. Find the skyline $S2$ for set $B2$.
4. Merge the two skylines $S1$ and $S2$.
 - Traverse $S1$ and $S2$ and pick up the piece belonging to the new skyline from one of them.

$$T(n) = 2T(n/2) + cn \quad \text{and} \quad T(1) = 1$$

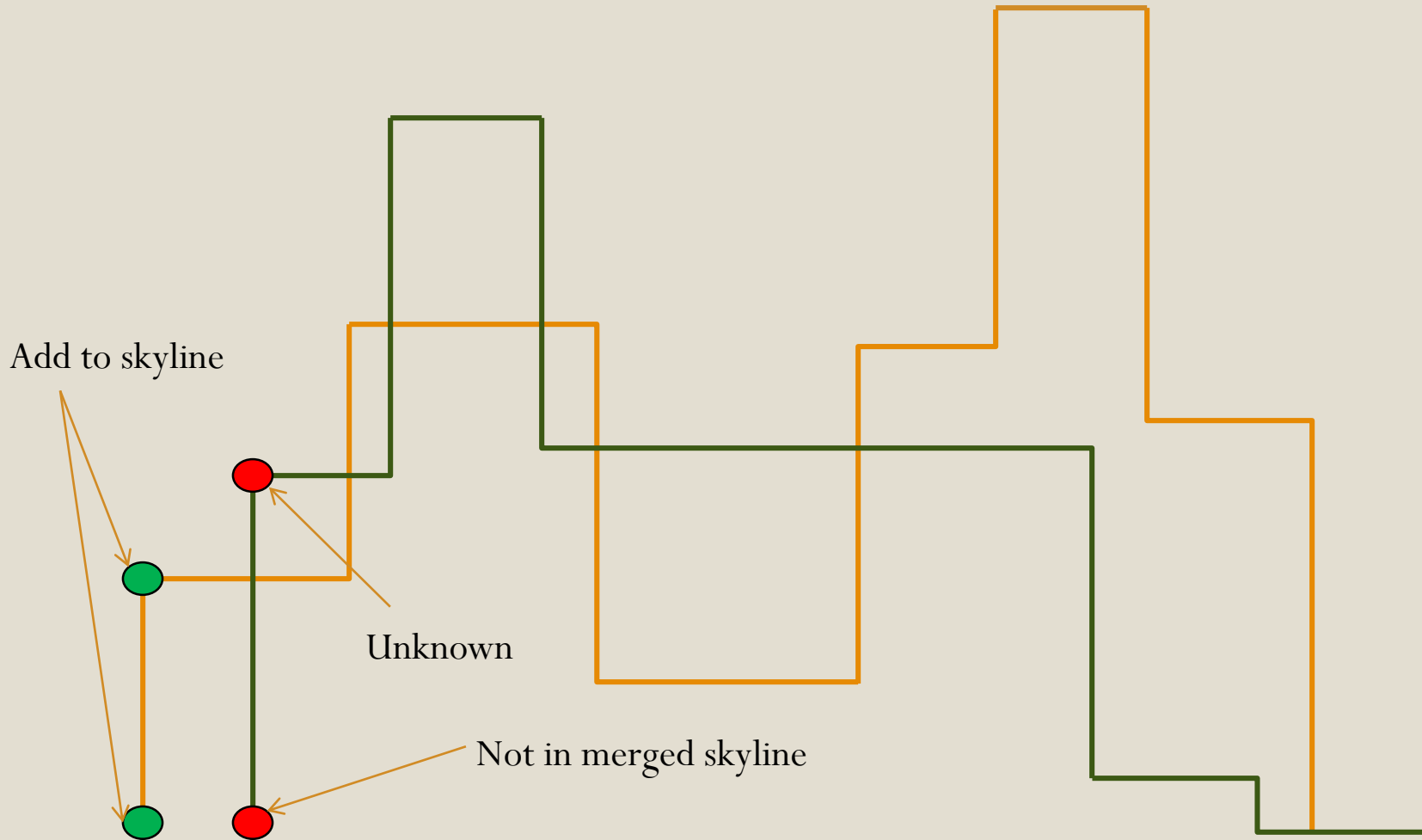
$$T(n) = O(n \log n)$$

Merge Skylines in Linear Time

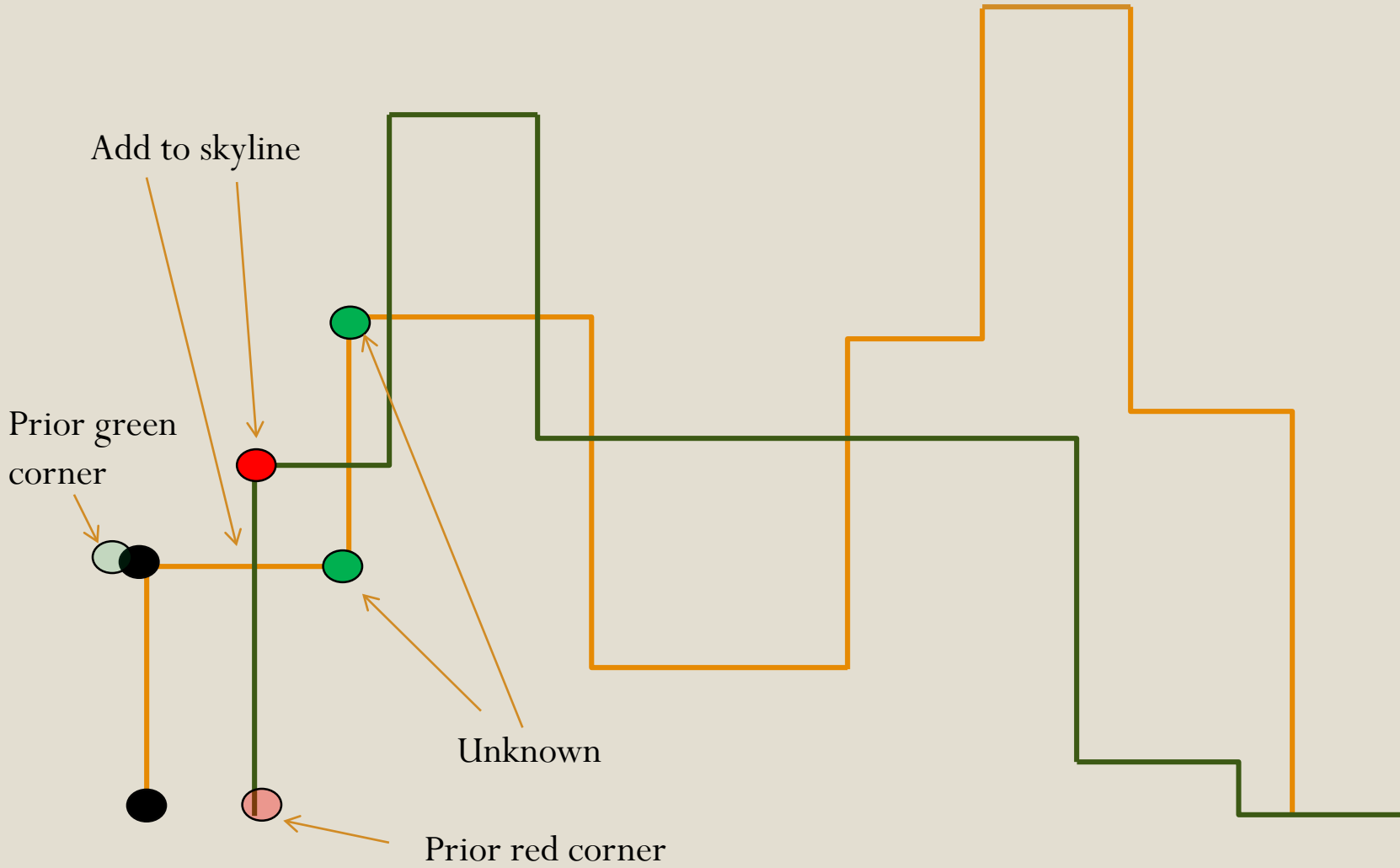
Skyline A: $(x_{1,A}, y_{1,A}), (x_{2,A}, y_{2,A}) \dots$

Skyline B: $(x_{1,B}, y_{1,B}), (x_{2,A}, y_{2,A}) \dots$

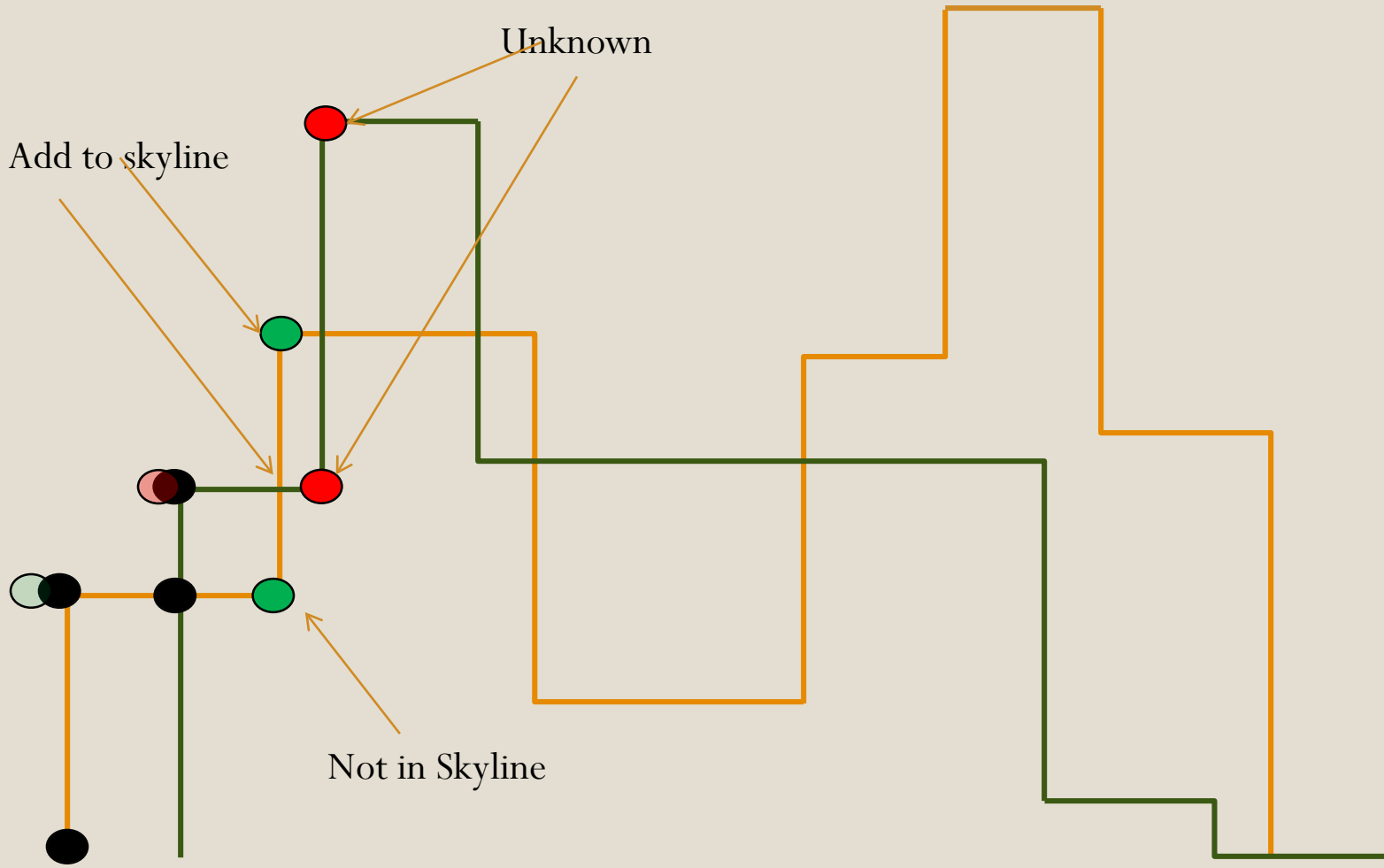
Traverse skyline A and B (left to right) and pick up the piece belonging to the new skyline from one of them.



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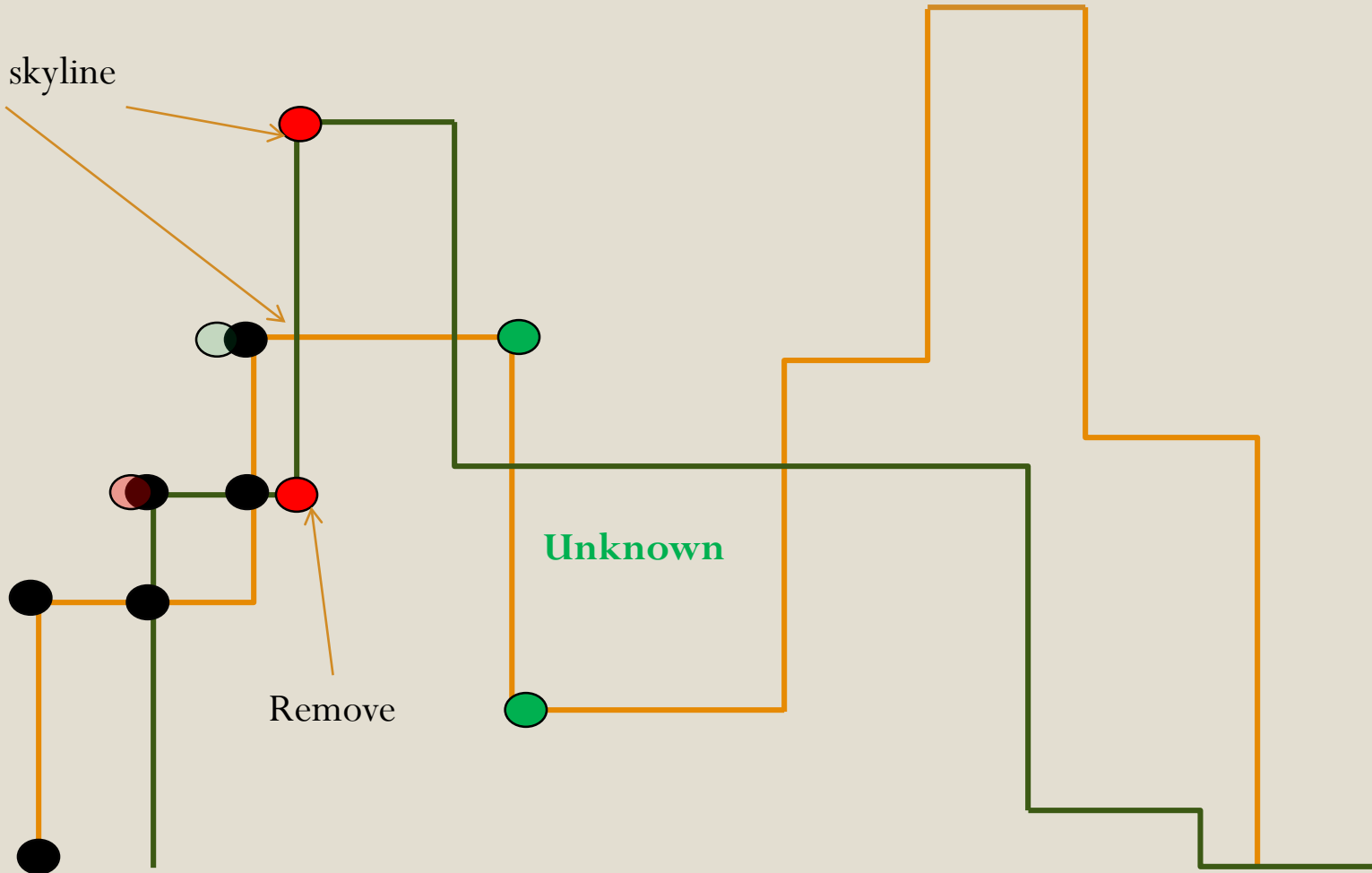


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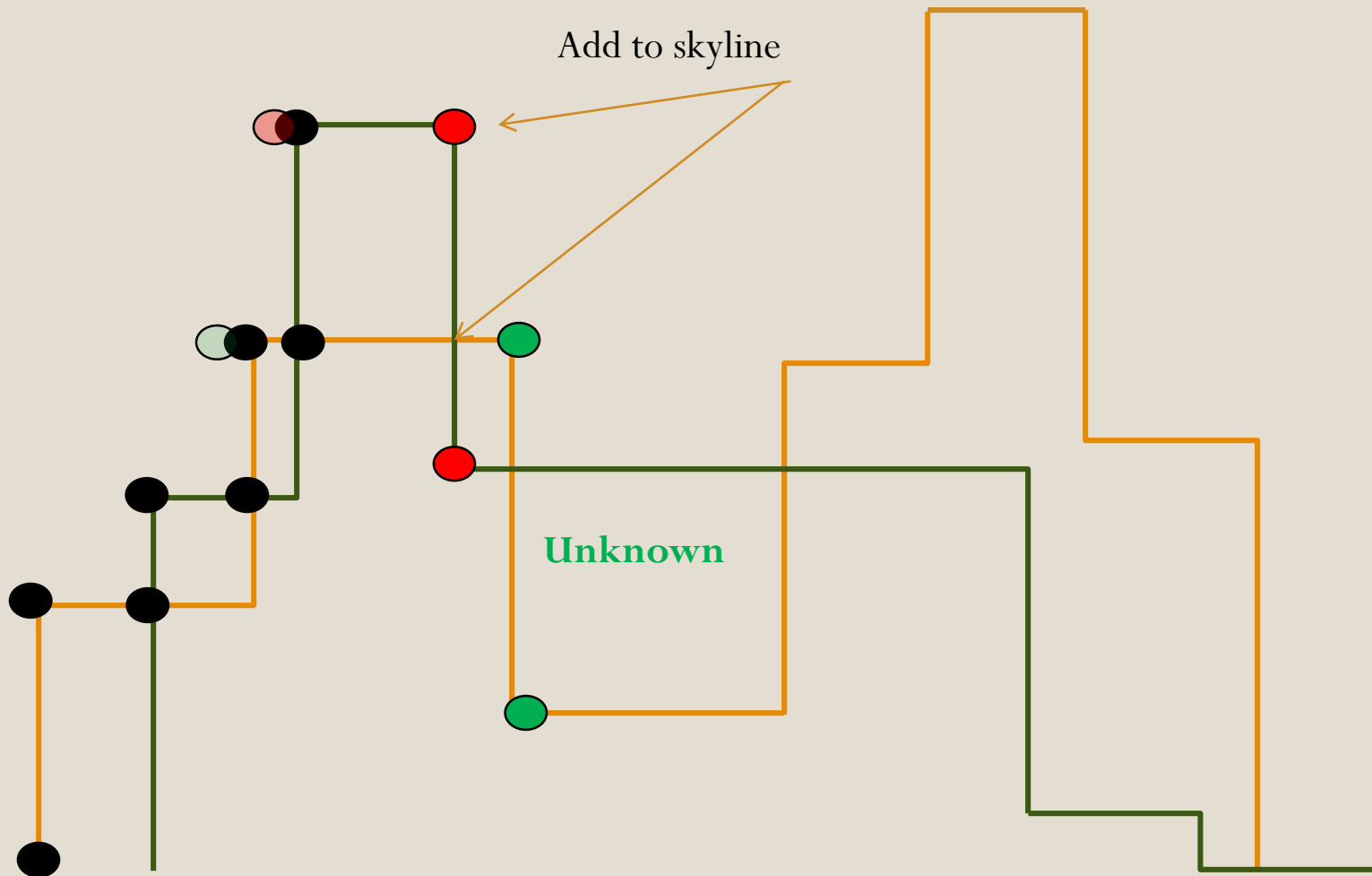
Add to skyline



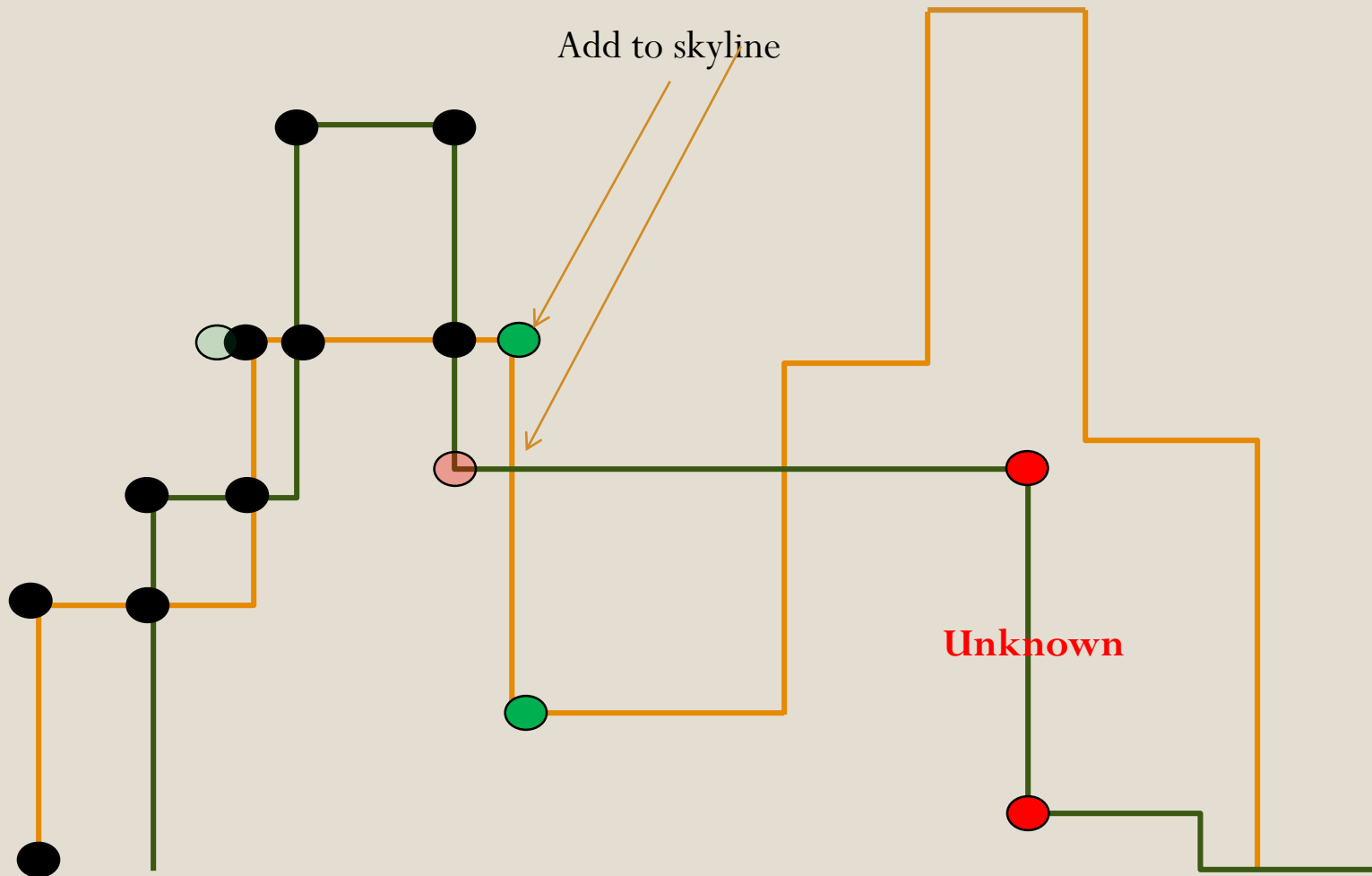
Remove

Unknown

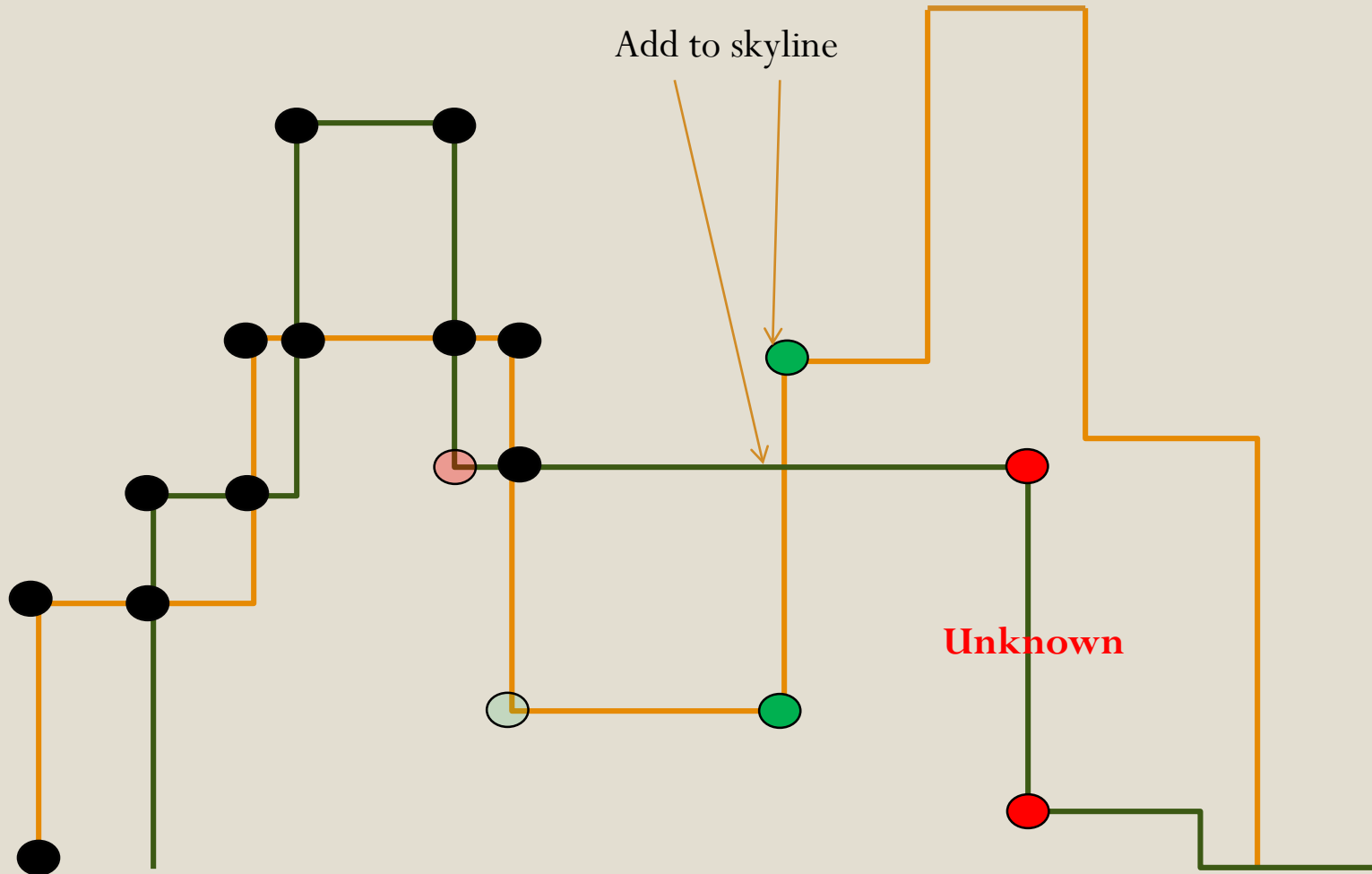
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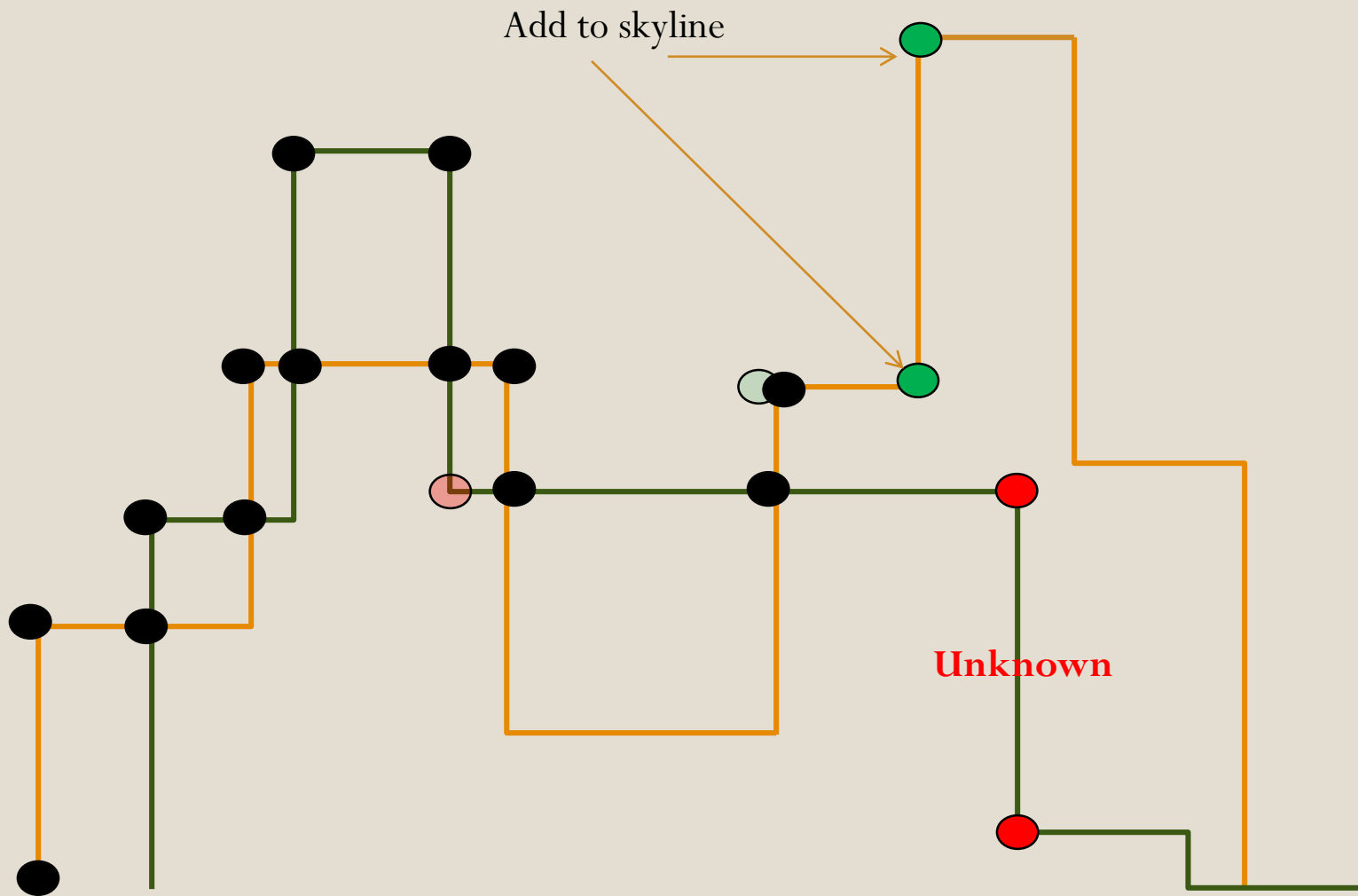
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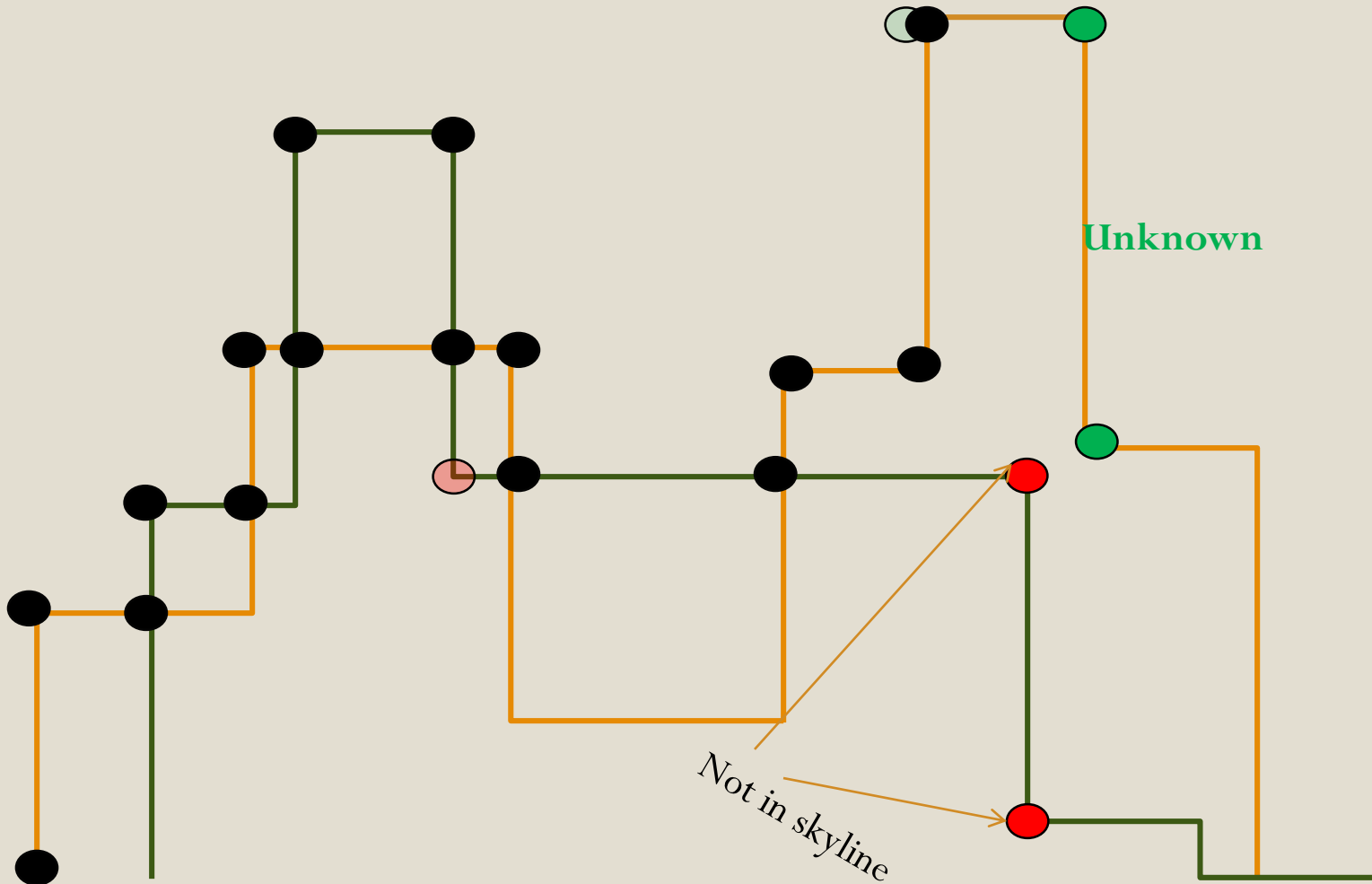
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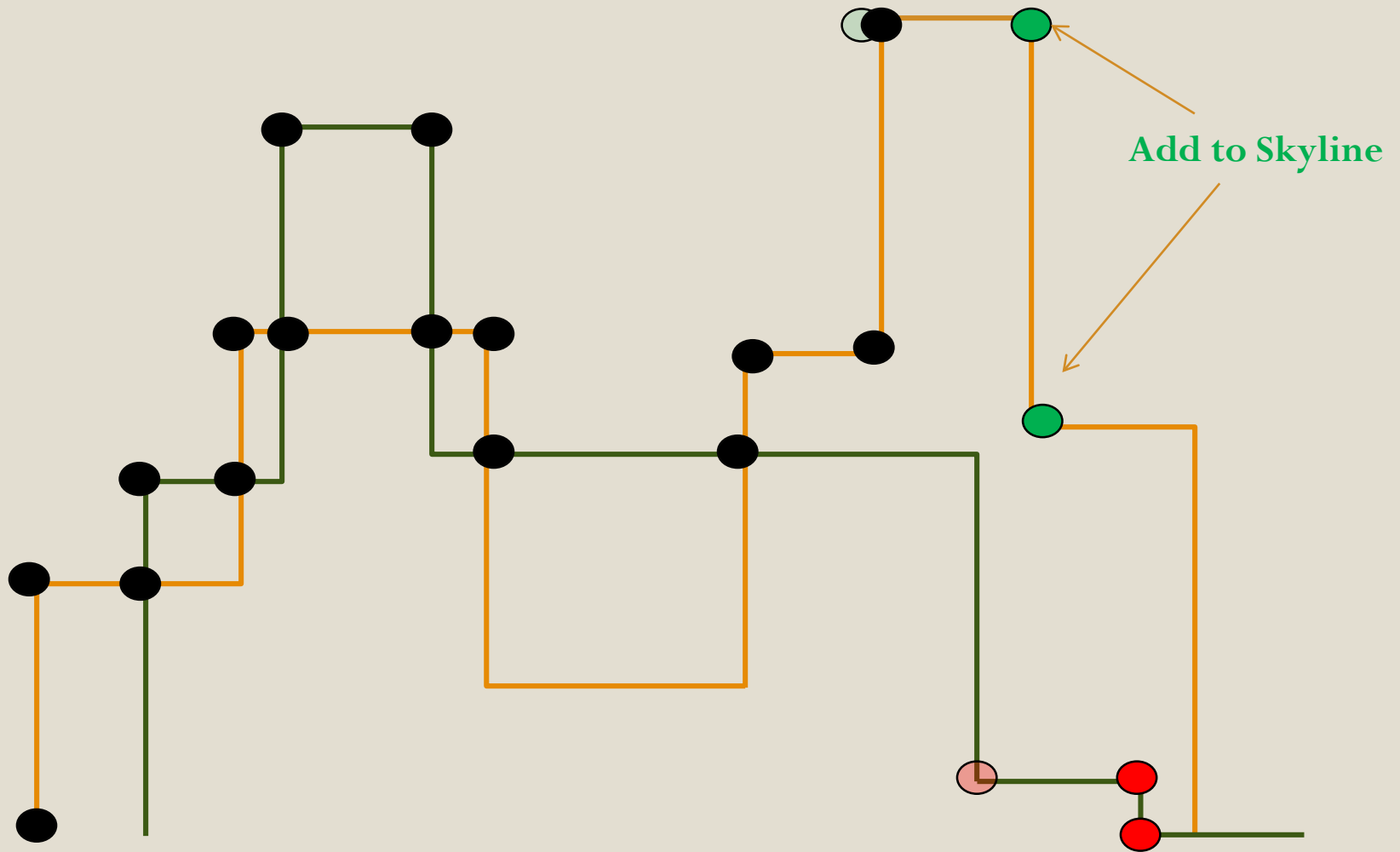
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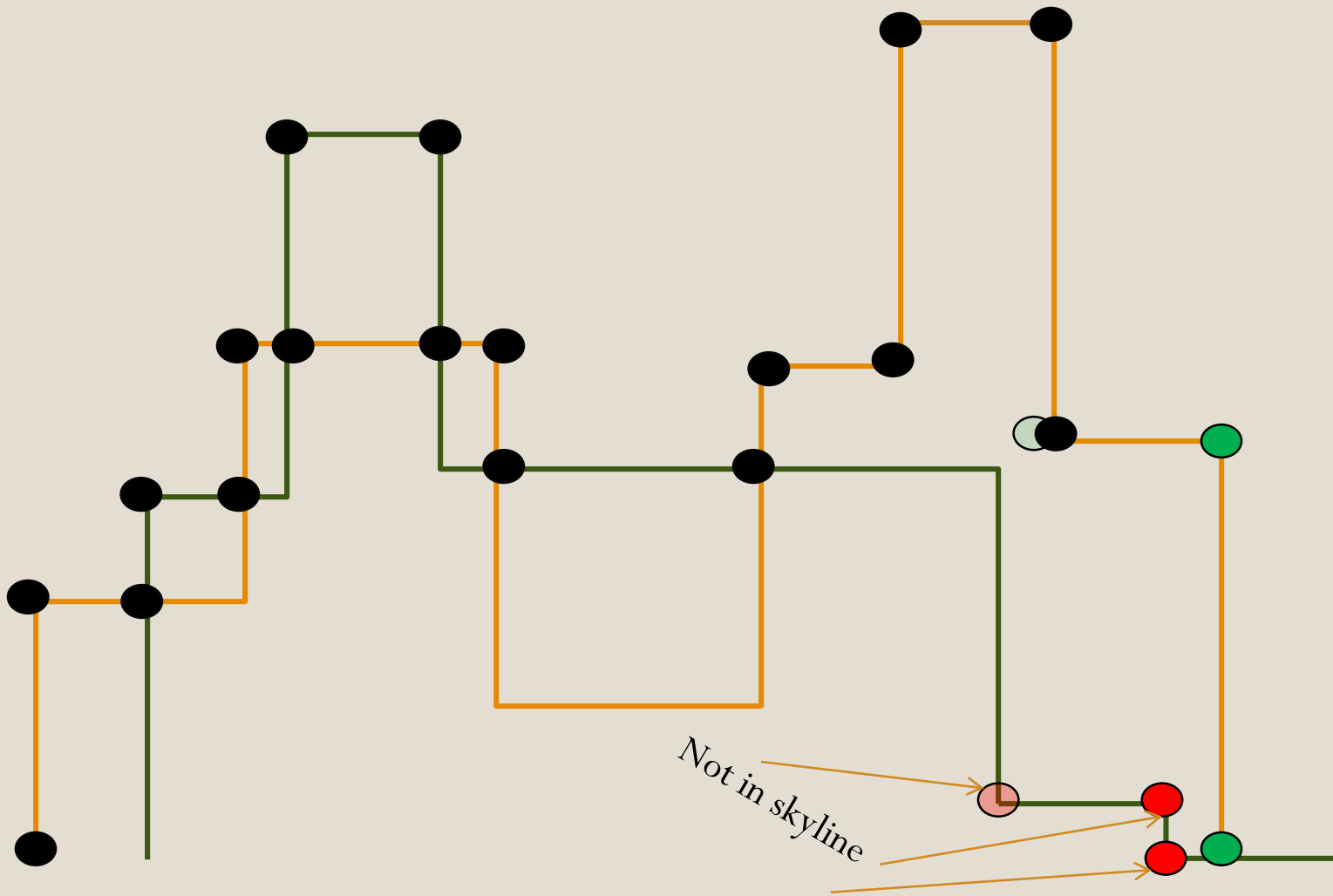
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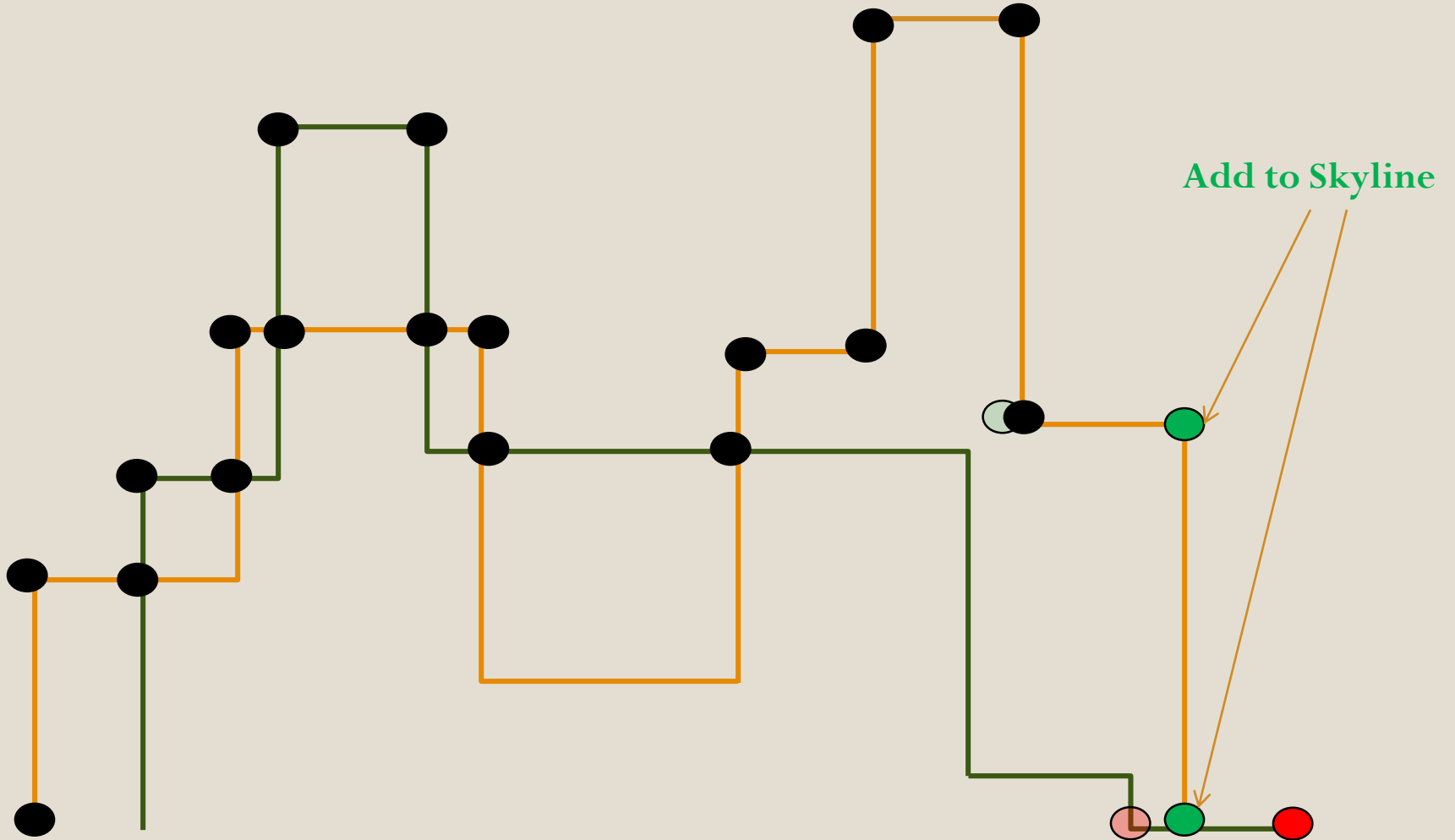
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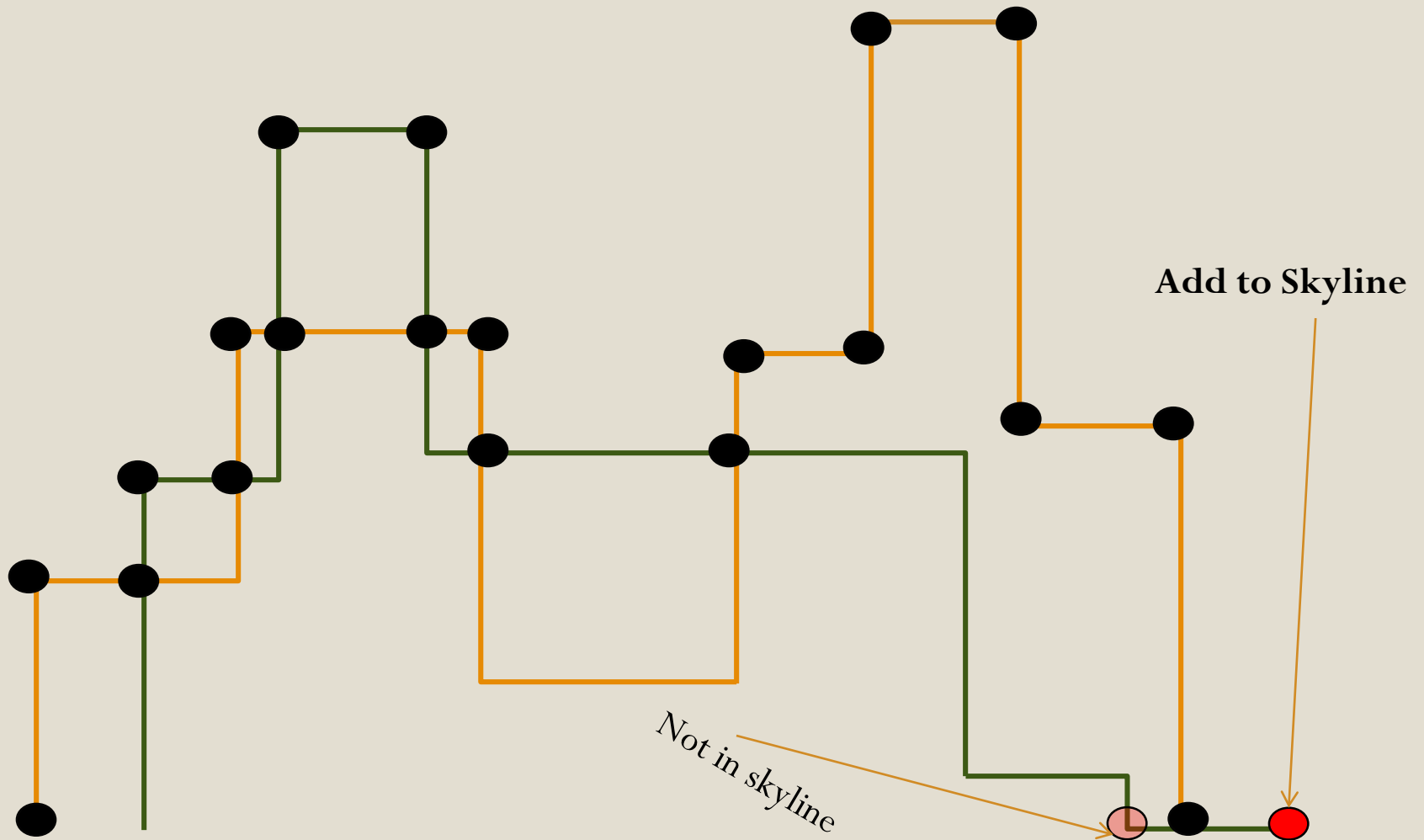
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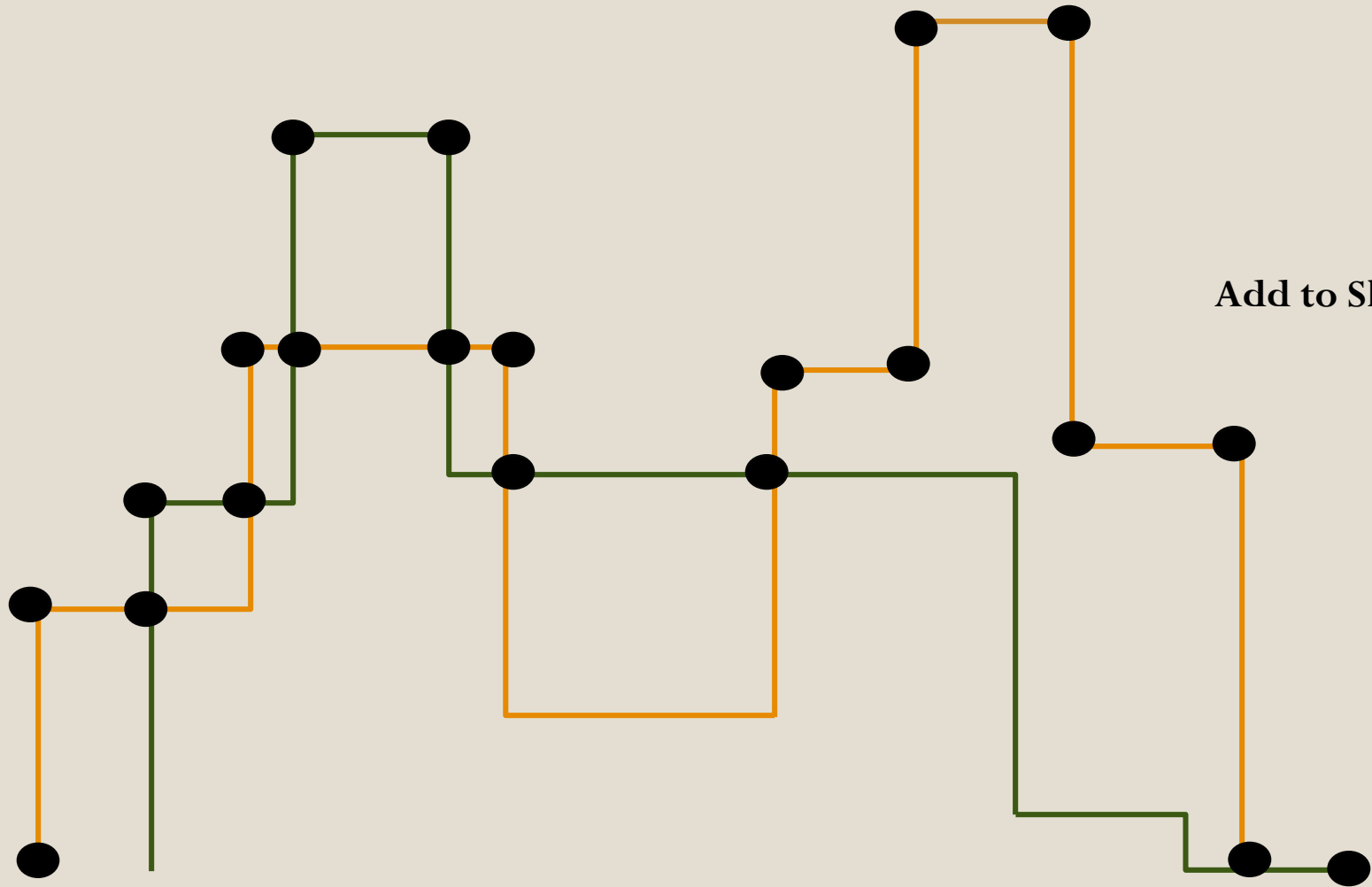
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Merge Skylines in Linear Time

- Iterate through skyline A and B in order
- Keep track of current/previous corners from skyline A and B
- Compute new corners (intersections), and decide which corners to include/exclude
- Requires $O(1)$ work in each iteration