## Homework 2 (150 points)

- 1. (10 + 10 + 20 points) An Interesting Concentration. Let X be the random variable over the sample space  $\{1, 2, ...\}$  such that  $\mathbb{P}[\mathbb{X} = i] = 2^{-i}$ .
  - (a) Compute  $\mu = \mathbb{E}[\mathbb{X}]$ .
  - (b) Define  $\mathbb{Y} = \mathbb{X} \mu$ . For  $0 \leq h \leq \ln 2$ , compute  $\mathbb{E}\left[\exp(h\mathbb{Y})\right]$ .
  - (c) Define  $\mathbb{S}_n = \mathbb{Y}^{(1)} + \cdots + \mathbb{Y}^{(n)}$ . Find the concentration bound for  $\mathbb{P}[\mathbb{S}_n \ge t]$  using the technique of Chernoff bound.

- 2. (10 + 10 + 20 points) Concentration of Sum of Poisson Distribution. Let X be the random variable over the sample space  $\{0, 1, ...\}$  such that  $\mathbb{P}[\mathbb{X} = i] = \exp(-\mu)\frac{\mu^i}{i!}$ .
  - (a) Prove that  $\mathbb{E}[\mathbb{X}] = \mu$ .
  - (b) Define  $\mathbb{Y} = \mathbb{X} \mu$ . For positive *h*, compute  $\mathbb{E}\left[\exp(h\mathbb{Y})\right]$ .
  - (c) Define  $\mathbb{S}_n = \mathbb{Y}^{(1)} + \cdots + \mathbb{Y}^{(n)}$ . Find the concentration bound for  $\mathbb{P}[\mathbb{S}_n \ge t]$  using the technique of Chernoff bound. (You might find it useful to use a variable m such that  $m = n\mu$  in the final bound.)

- 3. (10 + 10 points) Coin Tossing. Let X be the uniform distribution over the sample space  $\{0, 1\}$ .
  - (a) Let  $S_n = \mathbb{X}^{(1)} + \cdots + \mathbb{X}^{(n)}$ . Given a fixed values of m, how will you choose n such that  $\mathbb{P}[\mathbb{S}_n \ge m] \le (1-\varepsilon)$ ?
  - (b) Use the above result to prove the concentration bound in Problem 1 part c.

4. (10 points) Concentration of Matrix rank. Let  $\mathbb{M}$  be a distribution over  $n \times n$  matrices, where each element is selected uniformly and independently at random from the set  $\Omega$ . State and prove a concentration bound for the rank of  $\mathbb{M}$  around its median or mean.

5. (40 points) Prefix-sum of Coins are Close to their respective Mean. Let  $\mathbb{X}$  be a distribution over  $\{0,1\}$  such that  $\mathbb{P}[\mathbb{X}=1] = p$  and  $\mathbb{P}[\mathbb{X}=0] = (1-p)$ . We consider the sum  $S_n = \mathbb{X}^{(1)} + \cdots + \mathbb{X}^{(n)}$ .

Chernoff-Hoeffding's bound states the following. It says that the probability of the sum  $S_n$  exceeding the expectation by t is very small. For example, we can say that

$$\mathbb{P}\left[\mathbb{S}_n \ge pn + t\right] \le \exp\left(-2t^2/n\right)$$

Intuitively, suppose we reject any outcome of the coins such that  $\mathbb{S}_n \ge pn + t$ . Then, this bound says that the probability of rejecting is at most  $\exp\left(-2t^2/n\right)$ .

We want to claim that " $\mathbb{S}_n$  never exceeded the expectation in any prefix." Let me elaborate. Suppose we reject any coin such that  $\mathbb{S}_i \ge p \cdot i + t$  for any  $i \in \{1, \ldots, n\}$ . Formally, we reject if there exists  $i \in \{1, \ldots, n\}$  such that  $\mathbb{S}_i \ge p \cdot i + t$ . Note that this rejection rule is *more stringent* than the previous rejection criterion. Our goal is to prove that this rejection probability is small. In particular, prove that

$$\mathbb{P}\left[\exists i \in \{1, \dots, n\} \text{ s.t. } \mathbb{S}_i \ge p \cdot i + t\right] \le \exp\left(-2t^2/n\right)$$

Isn't this amazing? This bound is identical to the Chernoff-Hoeffding bound!

6. (Extra Credit) New bounds for Hoeffding's Lemma. Surprise me with a new statement/proof of Hoeffding's Lemma!