

Homework 1

1. ((4+1) + (4+1) points) Use Jensen's Inequality to prove the following inequalities.

- Let A , B and C be the three angles of a triangle. Prove that $\sin A + \sin B + \sin C \leq \frac{3\sqrt{3}}{2}$. Find the exact criterion for equality.
- Let \mathbb{X} be a random variable over a finite sample space Ω . We define the entropy of the random variable \mathbb{X} as $H(\mathbb{X}) := -\sum_{x \in \Omega} \mathbb{P}[\mathbb{X} = x] \lg \mathbb{P}[\mathbb{X} = x]$. Prove that $H(\mathbb{X}) \leq \lg|\Omega|$. Find the exact criterion for equality.

2. (10 points) Let M be a $n \times m$ matrix such that each of its entries is 0 or 1. Let $B = \{(i, j) : i \in [n], j \in [m], M(i, j) = 1\}$. Suppose $|B| \geq \varepsilon \cdot (nm)$, i.e., the set B covers at least an ε fraction of the entries of the matrix. Let R be the set of rows i such that $\sum_{j=1}^m M(i, j) \geq (\varepsilon/2) \cdot m$. Intuitively, R is the set of rows in the matrix where at least $\varepsilon/2$ fraction of the entries at in the set B . Prove that $|R| \geq (\varepsilon/2)n$.

3. (10 points) Let $\pi(x)$ represent the number of prime numbers that are less than x . For example, $\pi(16) = |\{2, 3, 5, 7, 11, 13\}| = 6$. The celebrated "Prime Number Theorem" shows that $\pi(x)$ roughly behaves like $x/\lg x$, for large enough x . In this question assume that $\pi(x) = x/\lg x$.

Let $r(x)$ represent the number of bits that are needed to represent the number x in binary. For example, $r(9) = 4$.

Let $\Omega_n = \{p : p \text{ is a prime, and } r(p) \leq n\}$. Let \mathbb{X}_n be uniform distribution over Ω_n . What is $\mathbb{E}[r(\mathbb{X}_n)]$?

4. (10 points) Let $f(x)$ be a convex upward function for $x \geq 1$. Prove that

$$\sum_{i=1}^n f(i) \leq \frac{f(1) + f(n)}{2} + \int_1^n f(x) dx$$

5. (3 + 3 + 2 + 5 + 2 points) The "Stirling Approximation" states that $x!$ is roughly $\sqrt{2\pi x} \left(\frac{x}{e}\right)^x$, for large enough x . So, this approximation cannot be applied to small values of x .

Let $p \in (0, 1/2)$ be a constant. Let $B(n, p)$ be the number of n -bit binary strings that have at most pn 1s. Note that we have:

$$B(n, p) = \sum_{0 \leq i \leq pn} \binom{n}{i}$$

We will upper bound this number.

(a) Prove that, for $1 \leq i \leq pn$, we have:

$$\binom{n}{i-1} \leq \frac{p}{1-p} \binom{n}{i}$$

(b) Prove that, for $0 \leq k \leq pn$, we have:

$$\binom{n}{pn-k} \leq \left(\frac{p}{1-p}\right)^k \binom{n}{pn}$$

(c) Prove that

$$B(n, p) \leq \left(\frac{1-p}{1-2p}\right) \binom{n}{pn}$$

(d) Let $h: [0, 1] \rightarrow [0, 1]$ be the binary entropy function defined as follows:

$$h(x) = -x \lg x - (1-x) \lg(1-x)$$

Use Stirling Approximation to show that

$$\binom{n}{pn} \text{ is roughly } \frac{2^{h(p)n}}{\sqrt{2p(1-p)\pi n}}$$

(e) Prove that, for large enough n , $B(n, p) \leq 2^{h(p)n}$.

6. (Extra Credit) Suppose we are given a black-box B . If we ping this black-box, it outputs a uniformly random real number in the range $[0, 1]$.

Let S_n denote the set of all permutations of size n (A size n permutation is a bijection from $[n] \rightarrow [n]$). We are interested in constructing an algorithm that uses this black-box to output a permutation that is drawn uniformly at random from S_n . Consider the following conjectured approach:

function UNIFORM_PERMUTATION_GENERATOR(n)

for $i = 1$ to n **do**

 Ping B to obtain to obtain a sample x_i

end for

 Let $(x_{i_1}, \dots, x_{i_n})$ be the sorting of the sequence (x_1, \dots, x_n) in an increasing order. (If there are collisions, i.e. there are $i < j$ such that $x_i = x_j$, then the sorted sequence has x_i before x_j)

 Let π be the permutation such that, for $k \in [n]$, we have $\pi(k) = i_k$.

 Output π

end function

Let \mathbb{X}_n be the random variable corresponding to the output of the algorithm described above. Prove or disprove that, for any $\pi \in S_n$, we have

$$\mathbb{P}[\mathbb{X}_n = \pi] = \frac{1}{n!}$$