

Lecture 22: Linearity Testing

Note: Boolean Functions

- In general, we consider boolean functions to be functions that map $\{0, 1\}^n \rightarrow \{0, 1\}$
- For ease of analysis we will interpret boolean functions to be $\{0, 1\}^n \rightarrow \{+1, -1\}$
- In particular, the boolean output 0 will correspond to the output +1, and the boolean output 1 will correspond to the output -1
- For example, $S \cdot x$ is a linear function that takes x as input and outputs 0/1 as output. In our context, the function χ_S corresponds to this linear function

Blum-Luby-Rubinfeld Linearity Testing

- Given oracle access to a function f ascertain whether the function is close to linear or not

BLR f :

- Sample $a, b \sim \mathbb{U}_{\{0,1\}^n}$
- Evaluate $x = f(a)$, $y = f(b)$, and $z = f(a + b)$
- Return $(x \cdot y == z)$

We want to prove: “The probability that this algorithm outputs true is high” if and only if “ f is close to linear”

Agreement of two Boolean Functions

- Let $f, g: \{0, 1\}^n \rightarrow \{+1, -1\}$ be two boolean functions
- The probability that f and g agree is

$$\mathbb{P} \left[f(\mathbb{U}_{\{0,1\}^n}) = g(\mathbb{U}_{\{0,1\}^n}) \right]$$

- Note that

$$\mathbb{P} \left[f(\mathbb{U}_{\{0,1\}^n}) = g(\mathbb{U}_{\{0,1\}^n}) \right] = 1 - \varepsilon \text{ if and only if } \langle f, g \rangle = 1 - 2\varepsilon$$

- And recall that $\langle f, g \rangle = \mathbb{E} [f(\mathbb{U})g(\mathbb{U})]$

BLR-Testing Theorem Formulation

The claim

“The probability that this algorithm outputs true is high” if and only if “ f is close to linear”

Is equivalent to proving

“ $\mathbb{E} [f(\mathbb{A})f(\mathbb{B})f(\mathbb{A} + \mathbb{B})]$ is high” if and only if “ f is close to linear”

Is equivalent to proving

“ $\mathbb{E} [f(\mathbb{A})f(\mathbb{B})f(\mathbb{A} + \mathbb{B})]$ is high” if and only if “There exists $S \in \{0, 1\}^n$ such that $\langle f, \chi_S \rangle$ is high”

Is equivalent to proving

“ $\mathbb{E} [f(\mathbb{A})f(\mathbb{B})f(\mathbb{A} + \mathbb{B})]$ is high” if and only if “ $\max_S \hat{f}(S)$ is high”

$$\mathbb{E} [f(\mathbb{A})f(\mathbb{B})f(\mathbb{A} + \mathbb{B})] = \frac{1}{N^2} \sum_{a,b} f(a)f(b)f(a+b)$$

Recall that $f = \sum_S \hat{f}(S)\chi_S$. So, we have

$$\begin{aligned} & \frac{1}{N^2} \sum_{a,b} f(a)f(b)f(a+b) \\ &= \frac{1}{N^2} \sum_{a,b} \left(\sum_R \hat{f}(R)\chi_R(a) \right) \left(\sum_S \hat{f}(S)\chi_S(b) \right) \left(\sum_T \hat{f}(T)\chi_T(a+b) \right) \\ &= \frac{1}{N^2} \sum_{a,b,R,S,T} \hat{f}(R)\hat{f}(S)\hat{f}(T)\chi_R(a)\chi_S(b)\chi_T(a+b) \end{aligned}$$

$$\begin{aligned}
& \frac{1}{N^2} \sum_{a,b,R,S,T} \hat{f}(R)\hat{f}(S)\hat{f}(T)\chi_R(a)\chi_S(b)\chi_T(a)\chi_T(b) \\
&= \frac{1}{N^2} \sum_{a,R,S} \hat{f}(R)\hat{f}(S)\chi_R(a) \sum_T \hat{f}(T)\chi_T(a) \sum_b \chi_S(b)\chi_T(b) \\
&= \frac{1}{N} \sum_{a,R,S} \hat{f}(R)\hat{f}(S)\chi_R(a) \sum_T \hat{f}(T)\chi_T(a) \langle \chi_S, \chi_T \rangle \\
&= \frac{1}{N} \sum_{a,R,S} \hat{f}(R)\hat{f}(S)\chi_R(a) \hat{f}(S)\chi_S(a) \\
&= \frac{1}{N} \sum_S \hat{f}(S)^2 \sum_R \hat{f}(R) \sum_a \chi_R(a)\chi_S(a) \\
&= \sum_S \hat{f}(S)^2 \sum_R \hat{f}(R) \langle \chi_R, \chi_S \rangle = \sum_S \hat{f}(S)^3
\end{aligned}$$

“ $\mathbb{E} [f(\mathbb{A})f(\mathbb{B})f(\mathbb{A} + \mathbb{B})]$ is high” if and only if “ $\max \hat{f}(S)$ is high”

Is equivalent to proving

“ $\sum_S \hat{f}(S)^3$ is high” if and only if “ $\max \hat{f}(S)$ is high”

This result is false unless we add the constraint that f is boolean
(Think: Why is this false?)

- Since f is boolean, we have $\langle f, f \rangle = 1$
- By Parseval's identity we get

$$1 = \sum_S \hat{f}(S)^2$$

BLR-Theorem Statement (Final)

Under the constraint

$$\sum_S \hat{f}(S)^2 = 1$$

We have

“ $\sum_S \hat{f}(S)^3 \approx 1$ ” if and only if “ $\max \hat{f}(S) \approx 1$ ”

Formalize and Prove this inequality

Given the promise that f is close to a linear function, can we recover S such that f is close to χ_S ?