Lecture 21: Extractors (Leftover Hash Lemma)

Randomness Extraction

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2-Universal Hash Function Family

- Let $\mathcal{F}_{n,m}$ be the set of all function $f: \{0,1\}^n \to \{0,1\}^m$
- *H* is a distribution over the sample space $\mathcal{F}_{n,m}$

Definition (2-Universal Hash Function Family)

For every distinct $x_1, x_2 \in \{0, 1\}^n$, we have:

$$\mathbb{P}_{h\sim H}[h(x_1)=h(x_2)]\leqslant \frac{1}{2^m}$$

- We want that the sampling $h \sim H$ can be efficiently performed by a randomized algorithm that takes a sample from U_d
- Intuitively, two separate inputs collide under h at the same probability that they collide under a random function from $\mathcal{F}_{n,m}$

Theorem (LHL)

Let H be a 2-universal Hash Function Family. For any X that is an (n, k)-source, the following is true:

$$2\text{SD}\left(\left(H,H(X)\right),\left(H,\mathbb{U}_{\{0,1\}^m}\right)\right) \leq \sqrt{\frac{M-1}{K}}$$

- That is, H is a good extractor for (n, k)-sources
- So, we need to construct the family *H* that can be sampled using only *d*-bits of randomness, and we want *d* to be as small as possible
- Note about the proof: We will see a more general Fourier-based proof, because there is another result, namely "Lopsided-LHL," that (as far as I know) cannot be proven using elementary combinatorial techniques

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- We will use $M = 2^m$ and $K = 2^k$
- We will use U_m to represent the distribution $\mathbb{U}_{\{0,1\}^m}$

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Leftover Hash Lemma

• We bound the SD as follows:

$$2SD ((H, H(X)), (H, U_m))$$

$$= \mathbb{E}_{h \sim H} \left[2SD (h(X), U_m) \right]$$

$$= \mathbb{E}_{h \sim H} \left[\sum_{y \in \{0,1\}^m} |h(X)(y) - U_m(y)| \right]$$

$$\leqslant \mathbb{E}_{h \sim H} \left[M^{1/2} \left(\sum_{y \in \{0,1\}^m} (h(X)(y) - U_m(y))^2 \right)^{1/2} \right], \quad \text{Cauchy-Schwartz}$$

$$= M\mathbb{E}_{h \sim H} \left[\sqrt{\|h(X) - U_m\|_2^2} \right]$$

$$\leqslant M \sqrt{\mathbb{E}_{h \sim H} \left[\|h(X) - U_m\|_2^2 \right]}, \quad \text{Jensen's}$$

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Leftover Hash Lemma

• Let us upper bound $\|h(X) - U_m\|_2^2$

$$\begin{split} & \|h(X) - U_m\|_2^2 \\ &= \sum_{S \in \{0,1\}^m} (h(\widehat{X}) - U_m)(S)^2, \qquad \text{Parseval's} \\ &= \sum_{S \in \{0,1\}^m: \ S \neq \emptyset} \widehat{h(X)}(S)^2 \\ &= \sum_{S \in \{0,1\}^m} \widehat{h(X)}(S)^2 - \widehat{h(X)}(S = \emptyset)^2 \\ &= \|h(X)\|_2^2 - 1/M^2 \end{split}$$

• So, we have the bound:

$$2\mathrm{SD}\left((H,H(X)),(H,U_m)\right) \leqslant M\sqrt{\mathbb{E}_{h\sim H}\left[\left\|h(X)\right\|_2^2 - M^{-2}\right]}$$

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• So, it suffices to upper bound
$$\mathbb{E}_{h\sim H}\left[\left\|h(X)\right\|_{2}^{2}\right]$$

$$= \mathbb{E}_{h \sim H} \left[\|h(X)\|_{2}^{2} \right]$$

= $\mathbb{E}_{h \sim H} \mathbb{E}_{y \sim U_{m}} \left[h(X)(y)^{2} \right]$
= $\mathbb{E}_{h \sim H} \mathbb{E}_{y \sim U_{m}} \left[\mathbb{P} \left[h(X^{(1)}) = y \land h(X^{(2)}) = y \right] \right]$
= $\mathbb{E}_{h \sim H} \mathbb{E}_{y \sim U_{m}} \left[\mathbb{P} \left[X^{(1)} = X^{(2)} \right] \mathbb{P} \left[h(X^{(1)}) = h(X^{(2)}) = y | X^{(1)} = X^{(2)} \right] \right]$
+ $\mathbb{E}_{h \sim H} \mathbb{E}_{y \sim U_{m}} \left[\mathbb{P} \left[X^{(1)} \neq X^{(2)} \right] \mathbb{P} \left[h(X^{(1)}) = h(X^{(2)}) = y | X^{(1)} \neq X^{(2)} \right] \right]$

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• The first term:

$$\mathbb{P}\left[X^{(1)} = X^{(2)}\right] \mathbb{E}_{h \sim H} \frac{1}{M} \sum_{y \in \{0,1\}^m} \mathbb{P}\left[h(X^{(1)}) = h(X^{(2)}) = y|X^{(1)} = X^{(2)}\right]$$

= $\mathbb{P}\left[X^{(1)} = X^{(2)}\right] \mathbb{E}_{h \sim H} \frac{1}{M} \mathbb{P}\left[h(X^{(1)}) = h(X^{(2)})|X^{(1)} = X^{(2)}\right]$
= $\mathbb{P}\left[X^{(1)} = X^{(2)}\right] \mathbb{E}_{h \sim H} \frac{1}{M} \cdot 1$
= $\frac{1}{M} \cdot \mathbb{P}\left[X^{(1)} = X^{(2)}\right]$

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$$\begin{split} &\frac{1}{M} \cdot \mathbb{P}\left[X^{(1)} \neq X^{(2)}\right] \mathbb{E}_{h \sim H} \mathbb{P}\left[h(X^{(1)}) = h(X^{(2)}) | X^{(1)} \neq X^{(2)}\right] \\ &\leq \frac{1}{M^2} \mathbb{P}\left[X^{(1)} \neq X^{(2)}\right] \\ &= \frac{1}{M^2} (1 - \mathbb{P}\left[X^{(1)} = X^{(2)}\right]) \end{split}$$

• So, we have:

$$E_{h\sim H} \left[\left\| h(X) \right\|_{2}^{2} \right] - \frac{1}{M^{2}}$$

$$\leq \mathbb{P} \left[X^{(1)} = X^{(2)} \right] \left(\frac{1}{M} - \frac{1}{M^{2}} \right)$$

$$\leq \frac{1}{K} \left(\frac{1}{M} - \frac{1}{M^{2}} \right)$$

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• So, overall we have:

$$2\mathrm{SD}\left((H,H(X)),(H,U_m)\right) \leqslant \sqrt{rac{M}{K}-rac{1}{K}}$$

• Hence the result

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