Lecture 11: Generalized Lovász Local Lemma

Lovász Local Lemma

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- We design an experiment with independent random variables $\mathbb{X}_1,\ldots,\mathbb{X}_m$
- We define bad events $\mathbb{A}_1, \ldots, \mathbb{A}_n$ where the bad event \mathbb{A}_i depends on the variables $(\mathbb{X}_{k_1}, \ldots, \mathbb{X}_{k_{n_i}})$
- We define $Vbl_i = \{k_1, \dots, k_{n_i}\}$, the set of all variables that the bad event A_i depends on
- The bad event \mathbb{A}_i can depend on the bad event \mathbb{A}_j if $Vbl_i \cap Vbl_j \neq \emptyset$
- Suppose each bad event \mathbb{A}_i depends on at most d other bad events
- Suppose we show that, for each bad event A_i, the probability of its occurrence P [A_i] ≤ p
- If $ep(d+1) \leqslant 1$, then

$$\mathbb{P}\left[\neg \mathbb{A}_{1}, \dots \neg \mathbb{A}_{n}\right] \geqslant \left(1 - \frac{1}{d+1}\right)^{n} > 0$$

Lovász Local Lemma

- We design an experiment with independent random variables $\mathbb{X}_1, \ldots, \mathbb{X}_m$
- We define bad events $\mathbb{A}_1, \ldots, \mathbb{A}_n$
- Let D_i be the set of indices of bad events that A_i depends on
- Suppose we exhibit the existence of numbers (x_1, \ldots, x_n) such that the following holds. For each *i*, we have:

$$\mathbb{P}\left[\mathbb{A}_{i}
ight]\leqslant x_{i}\prod_{j\in D_{i}}(1-x_{j})$$

• Then the following holds.

$$\mathbb{P}\left[\neg \mathbb{A}_1, \ldots, \neg \mathbb{A}_n\right] \geqslant \prod_{i=1}^n (1-x_i) > 0$$

Lovász Local Lemma

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- Prove Lovász Local Lemma using Generalized Lovász Local Lemma
- The numbers (x_1, \ldots, x_n) are <u>n</u>ot probabilities that add up to 1. This is an incorrect intuition
- Prove the following corollary of the generalized Lovász Local Lemma

Corollary

If, for all i, we have
$$\sum_{j\in D_i}\mathbb{P}\left[\mathbb{A}_j
ight] < 1/4$$
, then

$$\mathbb{P}\left[\neg \mathbb{A}_{1}, \ldots, \neg \mathbb{A}_{n}\right] \geqslant \prod_{i=1}^{n} (1 - 2\mathbb{P}\left[\mathbb{A}_{i}\right]) > 0$$

• Prove the results in the previous lecture using this corollary albeit with slightly worse parameters

Definition (Frugal Coloring)

A $\beta\text{-}\mathsf{Frugal}$ Coloring of a graph satisfies the following two conditions

It is a valid coloring, and

2 In the neighborhood N(v) of any vertex v, there are at most β vertices with the same color.

For example, a 1-frugal coloring of G is a coloring of G^2

We will show the following result

Theorem

For $\beta \in \mathbb{N}$, and a graph G with maximum degree $\Delta \ge \beta^{\beta}$ there exists a β -frugal coloring using $16\Delta^{1+1/\beta}$ colors.

Note that a graph with maximum degree Δ can be 1-frugally colored with Δ^2+1 colors. We will prove the general result using Lovász Local Lemma

Randomly color the vertices of the graph using C colors. We will consider two types of bad events.

- A_e, where e ∈ E(G): If the two vertices at the endpoints of the edge e receive the same color. These will be called type-1 bad events.
- $\mathbb{B}_{\{u_1,\ldots,u_{\beta+1}\}}$, where $u_1,\ldots,u_{\beta+1} \in V(G)$: Suppose there exists a vertex v such that $u_1,\ldots,u_{\beta+1}$ are distinct vertices in N(G) with identical color. These will be called type-2 bad events.

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- Note that one type-1 bad event \mathbb{A}_e can depends on at most 2Δ other type-1 bad events $\mathbb{A}_{e'}$
- We are now interested in computing how many type-2 bad events can \mathbb{A}_{e} depend on. Consider a type-2 bad event $\mathbb{B}_{u_1,\ldots,u_{\beta+1}}$ such that that are in $u_1,\ldots,u_{\beta+1}\in N(\nu)$. Suppose e = (a, b). Note that a has at most Δ neighbors. So, there are at most Δ possible ways of choosing v. Note that we have $\begin{pmatrix} \Delta \\ \beta \end{pmatrix} \text{ ways of choosing the remaining vertices} \\ \{u_1, \ldots, u_{\beta+1}\} \setminus \{a\}. \text{ Similar case for } b \text{ as well. So, there are} \\ \text{at most } 2\Delta \begin{pmatrix} \Delta \\ \beta \end{pmatrix} \text{ type-2 events that } \mathbb{A}_e \text{ can depends on.}$

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Frugal Coloring

- Similarly, a type-2 event $\mathbb{B}_{u_1,...,u_{\beta+1}}$ can depends on $(\beta + 1)\Delta$ other type-1 bad events and $(\beta + 1)\Delta \begin{pmatrix} \Delta \\ \beta \end{pmatrix}$ other type-2 bad events
- Note that

$$\mathbb{P}\left[\mathbb{A}_{e}\right] \leqslant \frac{1}{C}$$
$$\mathbb{P}\left[\mathbb{B}_{u_{1},...,u_{\beta+1}}\right] \leqslant \frac{1}{C^{\beta}}$$

 $\bullet\,$ So, to prove that a $\beta\mbox{-frugal}$ coloring exists, it suffices to prove that

$$(eta+1)\Delta\cdotrac{1}{C}+(eta+1)\Deltaiggl(\Delta \ eta iggr)\cdotrac{1}{C^eta} <rac{1}{4}$$

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• We can use the upper bound
$$\binom{n}{k} \leq \left(\frac{en}{k}\right)^k$$
 to upper-bound the expression

• This is left as an exercise

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Now we are interested in computing the solution that is guaranteed by Lovász Local Lemma.

function Seq_LLL($X = \{X_1, ..., X_m\}, A = \{A_1, ..., A_n\}$) $X \leftarrow$ Random Evaluation while $\exists i \text{ s.t. } A_i$ is satisfied do Pick arbitrary A_i that is satisfied Re-sample all X_j such that $j \in Vbl_i$ end while Output Xend function

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Theorem

Suppose there exists (x_1, \ldots, x_n) such that, for all $i \in [n]$, we have $\mathbb{P}[\mathbb{A}_i] \leq x_i \prod_{j \in D_i} (1 - x_j)$. Then the expected number of times sequential Moser-Tardos samples the event A_i is at most $x_i/(1 - x_i)$ and, hence, the expected number of execution of the inner loop is at most $\sum_{i \in [n]} x_i/(1 - x_i)$.

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function Parallel LLL(X, A)
    X \leftarrow Random Evaluation
    while \exists i \text{ s.t. } \mathbb{A}_i is satisfied do
         Let S be a maximal independent set in the dependency
graph restricted to satisfied A_is
         \mathbb{X}_{\bigcup_{k \in S} \mathsf{Vbl}_k} \leftarrow \mathsf{Random Evaluation}
    end while
    Output 𝗶
end function
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Theorem

Suppose there exists an $\varepsilon > 0$ and (x_1, \ldots, x_n) such that $\mathbb{P}[\mathbb{A}_i] \leq (1 - \varepsilon) x_i \prod_{j \in D_i} (1 - x_j)$. The expected number of inner loops before all bad events are avoided is at most $O\left(\frac{1}{\varepsilon} \sum_{i \in [n]} \frac{x_i}{1 - x_i}\right)$.

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