

## Lecture 06: Repeated Squaring

# Recall

- Let  $(G, \circ)$  be a group with generator  $g$
- We define  $g^0 = e$ , where  $e \in G$  is the identity element of  $G$
- We define  $g^i = \overbrace{g \circ g \circ \cdots \circ g}^{i\text{-times}}$
- For example, the group  $(\mathbb{Z}_7^*, \times)$  is generated by 3 but not 2

# Motivation of Efficient Algorithm to Compute Exponentiation

- Suppose  $p$  is a prime number that is represented using 1000-bits
- Note that the number  $p$  is in the range  $[2^{999}, 2^{1000})$ . We shall summarize this by stating that  $p$  is roughly (in the order of)  $2^{1000}$ .
- Suppose we are interested to work on the field  $(\mathbb{Z}_p^*, \times)$  with generator  $g$
- Given input  $i \in \{0, 1, \dots, p - 1\}$ , we are interested in computing  $g^i \in \mathbb{Z}_p^*$

# First Attempt

Exp ( $i$ ):

- 1 prod =  $e$
- 2 For index in the range  $\{1, \dots, i\}$ :
  - 1 prod = prod  $\circ$   $g$
- 3 Return prod

- Note that this algorithm runs the inner loop  $i$  times. The number  $i$  can take values  $\{0, 1, \dots, p - 2\}$ . For example, if  $i \geq 2^{500}$  then the algorithm will run the inner loop more than the number of atoms in the universe. Effectively, the algorithm is useless
- The algorithm takes  $O(i)$  run-time. The size of the input  $i$  is  $\log i$ . So, this algorithm is an exponential time algorithm

## Second Attempt I

Exp ( $i$ ):

- 1 If  $i = 0$ : Return  $e$
- 2 If  $i$  is even:
  - 1  $\alpha = \text{Exp}(i/2)$
  - 2 Return  $\alpha \circ \alpha$
- 3 If  $i$  is odd:
  - 1  $\alpha = \text{Exp}((i - 1)/2)$
  - 2 Return  $\alpha \circ \alpha \circ g$

- Note that the argument to Exp becomes smaller by one-bit in recursive call. So, the algorithm performs (at most) 1000 recursive call. This is an efficient algorithm because it runs in time  $O(\log i)$

### A Few Optimizations.

- Testing whether  $i$  is even or not can be performed by computing  $i \& 1$  (here,  $\&$  is the bit-wise and of the binary representation of  $i$  and 1)
- Computing  $(i/2)$  when  $i$  is even, or computing  $(i - 1)/2$  when  $i$  is odd can be achieved by  $i \gg 1$  (that is, right-shift the binary representation of  $i$  by one position)

## Second Attempt III

The code shall look as follows

Exp ( $i$ ):

- ① If  $i = 0$ : Return  $e$
- ②  $j \gg 1$
- ③ If  $(i \& 1) == 0$ :
  - ①  $\alpha = \text{Exp}(j)$
  - ② Return  $\alpha \circ \alpha$
- ④ else:
  - ①  $\alpha = \text{Exp}(j)$
  - ② Return  $\alpha \circ \alpha \circ g$

## Second Attempt IV

- 1 The algorithm makes recursive calls. Can we further optimize and avoid recursive function calls? That is, can we unroll the recursion into a for loop?



# Final Attempt I

In the following code, we assume that we represent the prime  $p$  using  $t$ -bits. For example, we were considering  $t = 1000$  in the ongoing example.

We perform a preprocessing step to compute the following global variables.

## Global Preprocessing.

- 1 For index in the set  $\{0, 1, \dots, t - 1\}$ :
  - 1 If index == 0:  $\alpha_{\text{index}} = g$  and  $c_{\text{index}} = 1$
  - 2 Else:  $\alpha_{\text{index}} = \alpha_{\text{index}-1} \circ \alpha_{\text{index}-1}$  and  $c_{\text{index}} = (c_{\text{index}-1} \lll 1)$

- Note that  $\alpha_{\text{index}} = g^{2^{\text{index}}}$ , for all  $\text{index} \in \{0, 1, \dots, t - 1\}$
- Further, note that  $c_{\text{index}} = 2^{\text{index}}$ , for all  $\text{index} \in \{0, 1, \dots, t - 1\}$

## Final Attempt II

We shall use the preprocessed data to compute the exponentiation

Exp ( $i$ ):

- 1 prod = e
- 2 For index in the set  $\{0, 1, \dots, t - 1\}$ :
  - 1 If  $(i < c_{\text{index}})$  : Break
  - 2 If  $(i \& c_{\text{index}}) \neq 0$ : prod = prod  $\circ$   $\alpha_{\text{index}}$
- 3 Return prod

- Note that the test “the  $(1 + \text{index})$ -th bit in the binary representation of  $i$  is 1” is identical to the test  $(i \& c_{\text{index}}) \neq 0$
- If this test passes, then prod is multiplied by  $\alpha_{\text{index}} = g^{2^{\text{index}}}$
- Prove: This approach correctly calculates  $g^i$
- Note that the runtime is  $O(\log i)$  (that is, the algorithm is efficient)

- Let us consider a problem that shall use all the facts we studied about groups and fields in the last two lectures. There are multiple solutions with varying degree of complexities
- Compute

$$17^{2020} \pmod{23}$$

**Solution 1.**

- We can use repeated squaring directly to compute

$$\begin{aligned} 17^1 & \pmod{23} \\ 17^2 & \pmod{23} \\ 17^4 & \pmod{23} \\ & \vdots \\ 17^{1024} & \pmod{23} \end{aligned}$$

- Write 2020 in binary and compute  $17^{2020} \pmod{23}$  using the values computed above
- Although this is a correct and a tractable way to compute this value, it is computationally intensive and prone to errors (without a calculator)

## Solution 2.

- In homework you will prove that  $x^p = x \pmod p$ , where  $p$  is a prime and  $x$  is any integer
- You can use this fact to simplify the computation of  $17^{2020} \pmod{23}$  as follows

$$\begin{aligned} 17^{2020} \pmod{23} &= 17^{23} \cdot 17^{1997} \pmod{23} \\ &= \left(17^{23}\right)^2 \cdot 17^{1974} \pmod{23} \\ &\quad \vdots \\ &= \left(17^{23}\right)^{87} \cdot 17^{19} \pmod{23} \\ &= (17)^{87} \cdot 17^{19} \pmod{23}, && \text{using } x^p = x \pmod{p} \\ &= 17^{106} \pmod{23} \\ &= \left(17^{23}\right)^4 \cdot 17^{14} \pmod{23} \\ &= (17)^4 \cdot 17^{14} \pmod{23}, && \text{using } x^p = x \pmod{p} \\ &= 17^{18} \pmod{23} \end{aligned}$$

- This final expression can be computed using the repeated squaring technique

## Solution 3.

- In homework you will prove that  $x^{p-1} = 1 \pmod p$ , where  $p$  is a prime and  $x$  is any integer NOT divisible by  $p$  (there are also alternate proofs of this statement by considering the size of the subgroup of  $(\mathbb{Z}_p^*, \times)$  that is generated by  $x$ )
- So, we can compute the expression as follows

$$\begin{aligned} 17^{2020} \pmod{23} &= 17^{22} \cdot 17^{1998} \pmod{23} \\ &= \left(17^{22}\right)^2 \cdot 17^{1976} \pmod{23} \\ &\quad \vdots \\ &= \left(17^{22}\right)^{91} \cdot 17^{18} \pmod{23} \\ &= (1)^{91} \cdot 17^{18} \pmod{23}, && \text{using } x^{p-1} = 1 \pmod p \\ &= 17^{18} \pmod{23} \end{aligned}$$

- BTW, in general you can conclude that

$$x^n = x^{n \bmod p-1} \pmod{p},$$

for any integer  $n$  and any integer  $x$  that is not divisible by  $p$

- Now you can compute  $17^{18} \pmod{23}$  result using repeated squaring technique

$$17^1 = 17 \pmod{23}$$

$$17^2 = 13 \pmod{23}$$

$$17^4 = 8 \pmod{23}$$

$$17^8 = 18 \pmod{23}$$

$$17^{16} = 2 \pmod{23}$$



- Now, we have

$$\begin{aligned}17^{18} &= 17^{16+2} \pmod{23} \\ &= 17^{16} \cdot 17^2 \pmod{23} \\ &= 2 \cdot 13 \pmod{23} \\ &= 3 \pmod{23}\end{aligned}$$

- Therefore, we conclude that

$$17^{2020} = 17^{18} = 3 \pmod{23}.$$

That is our answer!