

# Lecture 12: $k$ -wise Independent Hash Function Family

# Intuition

- We want to design a hash function family  $\mathcal{H}$  from a field  $\mathbb{F}$  to  $\mathbb{F}$
- For  $h \xleftarrow{\$} \mathcal{H}$ , we want the first  $k$  inputs to be answered randomly

## Definition ( $k$ -wise Independence)

Let  $\mathcal{H} = \{h_1, \dots, h_t\}$  be a family of hash functions such that  $h_i: \mathcal{D} \rightarrow \mathcal{R}$ , where  $\mathcal{D}$  is the domain and  $\mathcal{R}$  is the range. For any distinct  $x_1, \dots, x_k \in \mathcal{D}$  and any  $y_1, \dots, y_k \in \mathcal{R}$  we have the following guarantee.

$$\mathbb{P} \left[ h(x_1) = y_1, \dots, h(x_k) = y_k : h \xleftarrow{s} \mathcal{H} \right] = \frac{1}{|\mathcal{R}|^k}$$

# Elaborating the Definition I

- From the definition, we can compute the probability that  $h(x_1) = y_1, \dots, h(x_{k-1}) = y_{k-1}$  (note we dropped the constraint that  $h(x_k) = y_k$ )
- For any  $x_k \in \mathcal{D}$ , we can simplify the probability as follows:

$$\begin{aligned} & \mathbb{P} \left[ h(x_1) = y_1, \dots, h(x_{k-1}) = y_{k-1} : h \stackrel{s}{\leftarrow} \mathcal{H} \right] \\ &= \sum_{y_k \in \mathcal{R}} \mathbb{P} \left[ h(x_1) = y_1, \dots, h(x_{k-1}) = y_{k-1}, f(x_k) = y_k : h \stackrel{s}{\leftarrow} \mathcal{H} \right] \\ &= \sum_{y_k \in \mathcal{R}} \frac{1}{|\mathcal{R}|^k} = \frac{1}{|\mathcal{R}|^{k-1}} \end{aligned}$$

- Remark: This proof-technique should seem similar to the solution to one of your HW01 problems!

## Elaborating the Definition II

- Proceeding inductively, we can prove the following statement.  
For all  $i \in \{1, 2, \dots, k\}$ , we have

$$\mathbb{P} \left[ h(x_1) = y_1, \dots, h(x_i) = y_i : h \stackrel{s}{\leftarrow} \mathcal{H} \right] = \frac{1}{|\mathcal{R}|^i}$$

- Using these guarantees, we can prove the following statements:

$$\mathbb{P} \left[ h(x_1) = y_1 : h \stackrel{s}{\leftarrow} \mathcal{H} \right] = \frac{1}{|\mathcal{R}|}$$

$$\mathbb{P} \left[ h(x_2) = y_2 | h(x_1) = y_1 : h \stackrel{s}{\leftarrow} \mathcal{H} \right] = \frac{1}{|\mathcal{R}|}$$

⋮

$$\mathbb{P} \left[ h(x_k) = y_k | h(x_1) = y_1, \dots, h(x_{k-1}) = y_{k-1} : h \stackrel{s}{\leftarrow} \mathcal{H} \right] = \frac{1}{|\mathcal{R}|}$$

## Elaborating the Definition III

- The definition of  $k$ -wise independent function insists that the probability of a random  $h \stackrel{\$}{\leftarrow} \mathcal{H}$  simultaneously mapping  $x_1 \mapsto y_1, x_2 \mapsto y_2, \dots, x_k \mapsto y_k$  is  $1/|\mathcal{R}|^k$
- The first input is uniformly randomly answered: We have proved above that  $h(x_1) = y_1$  is  $\frac{1}{|\mathcal{R}|}$ .
- The second input is uniformly randomly answered (even conditioned on the first input's answer): We have proved above that  $h(x_2) = y_2$  conditioned on  $h(x_1) = y_1$  is  $\frac{1}{|\mathcal{R}|}$ .
- Similarly, for any  $i \in \{1, \dots, k\}$ , the  $i$ -th input is uniformly randomly answered (even conditioned on the previous  $(i - 1)$  inputs' answers): We have proved above that  $h(x_i) = y_i$  conditioned on  $h(x_1) = y_1, \dots, h(x_{i-1}) = y_{i-1}$  is  $\frac{1}{|\mathcal{R}|}$
- Summarizing: It says that the first  $k$  inputs to the function are answered independently and uniformly at random

## Example Construction

- Let  $\mathcal{D} = \mathcal{R} = \mathbb{F}$ , where  $\mathbb{F}$  is a field
- The hash function family is defined as follows:

$$\mathcal{H} = \left\{ h_{a_0, a_1, \dots, a_{k-1}} : a_0, a_1, \dots, a_{k-1} \in \mathbb{F} \right\}$$

where the hash function  $h_{a_0, a_1, \dots, a_{k-1}} : \mathbb{F} \rightarrow \mathbb{F}$  is defined as follows

$$h_{a_0, a_1, \dots, a_{k-1}}(X) = a_0 + a_1X + a_2X^2 + \dots + a_{k-1}X^{k-1}$$

- Intuitively,  $\mathcal{H}$  is the set of all polynomials of degree  $< k$  with coefficients in the field  $\mathbb{F}$
- You have already proved in HW01 that this hash function family is  $k$ -wise independent!

- Let  $\mathcal{D} = \mathbb{F}^m$  and  $\mathcal{R} = \mathbb{F}$ , where  $\mathbb{F}$  is a field
- The hash function family is defined as follows:

$$\mathcal{H} = \{h_{a_1, \dots, a_m} : a_1, \dots, a_m \in \mathbb{F}\}$$

where the hash function  $h_{a_1, \dots, a_m} : \mathbb{F}^m \rightarrow \mathbb{F}$  is defined as follows

$$h_{a_1, \dots, a_m}(x_1, \dots, x_m) = a_1x_1 + \dots + a_mx_m$$

- Is this  $k$ -wise independent hash function family?