## Homework 4

- 1. (5 + 15 points) Let  $L = \mathbb{QR}_n$  and consider the 3-round ZK protocol taught in the class to prove that  $x \in L$ . We had seen that this protocol has completeness 1 and soundness 1/2.
  - (a) Formally write the protocol that sequentially runs this protocol *t*-times.
  - (b) This new protocol has completeness 1 and soundness  $1/2^t$ . Prove that this protocol is also ZK.
- 2. (5 + 15 points) In this question we shall define *witness indistinguishable proofs*, referred to as WI proofs, and their relation to ZK proofs.
  - (a) Consider a language  $L \in \mathsf{NP}$ . Let  $R_L$  be the set of all (x, w) such that w is a witness for  $x \in L$ .

A WI proof is a proof system with the following intuitive property. Consider  $x \in L$  and any two distinct witnesses w and w' such that  $(x, w) \in R_L$  and  $(x, w') \in R_L$ . The distribution of the transcripts produced by the interaction between an honest prover and any arbitrary efficient verifier when the prover uses the witness w is computationally indistinguishable from the case when the prover uses the witness w'. Formalize this intuition to define witness indistinguishable proofs.

(b) Prove that every ZK proof is also a WI proof.