

1.(b)

WX \ YZ	00	01	11	10
00	0	0	0	1
01	1	1	1	1
11	1	1	1	1
10	1	0	1	0

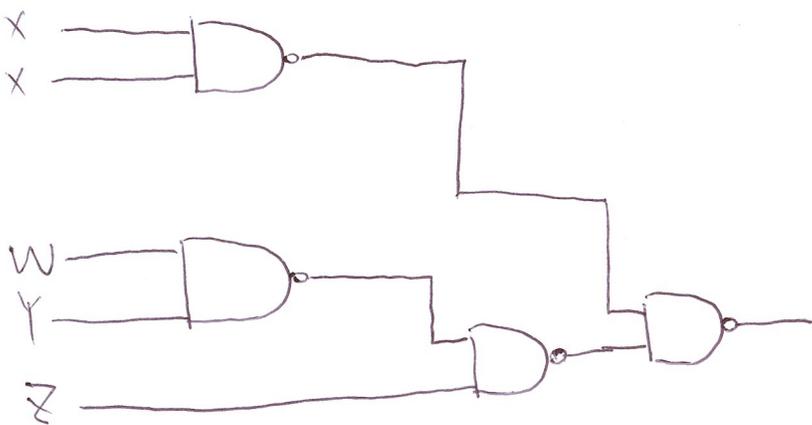
1.(c)

$$x + w'y'z' + wy'z' + wyz$$

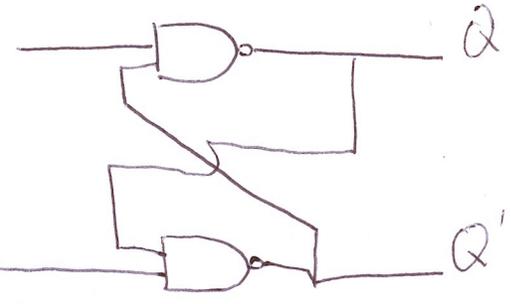
~~$$= (w \text{ NAND } y \text{ NAND } z) \text{ NAND } (x \text{ NAND } x)$$~~

$$= (w \text{ NAND } y \text{ NAND } z) \text{ NAND } (x \text{ NAND } x)$$

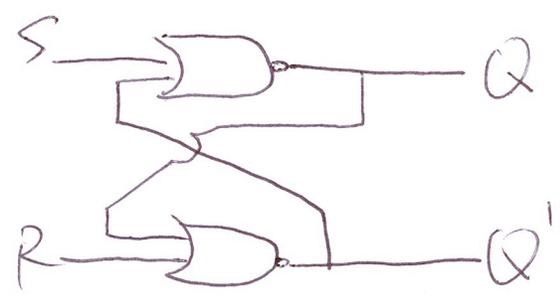
1.(d)



a)



b)



3. `int a = 5 // a is stored in data`
`int b[50] // b is bss`
`int main() // stored in text`
`{ int x // x stored stack`
`int * p = (int *) malloc(sizeof(int));`
`// stored in heap`

① The loader is a program that is used to run an executable file in a process.

- ② Before the program starts running, the loader allocates space for all the sections of the executable file (text, data, bss etc)
- ③ It loads into memory the executable and shared libraries
- ④ It also writes (resolves) any values in the executable to point to the functions/variables in the shared libraries
- ⑤ Once memory image is ready, the loader jumps to the `_start` entry point that calls `init()` of all libraries and initializes static constructors. Then it calls `main()` and the program begins
- ⑥ `_start` also calls `exit()` when `main()` returns.

5. "Dynamic Linker" is the part that loads and links the shared libraries for an executable when it is executed. Linking is often referred to as a process that is performed at compile time of the executable while a dynamic linker is in actuality a special loader that loads external shared libraries into a running process and then binds those shared libraries dynamically to the running process.

"Static linker" copies all library routines used in the program into executable file. It needs more space and memory but it is faster.

6. 1001010×1101

$$\begin{array}{r}
 1001010 \\
 \underline{1101} \\
 1001010 \\
 1001010 \\
 1001010 \\
 \hline
 1111000010 \\
 \begin{array}{cccccccc}
 2^9 & 2^8 & 2^7 & 2^6 & 2^5 & 2^4 & 2^3 & 2^2
 \end{array}
 \end{array}$$

$$2^1 + 2^6 + 2^7 + 2^8 + 2^9 = 962$$

7. $1011 \overline{)100101001}$

$$\begin{array}{r}
 11011 \\
 1011 \overline{)100101001} \\
 \underline{1011} \\
 1111 \\
 \underline{1011} \\
 10000 \\
 \underline{1011} \\
 1011 \\
 \underline{1011} \\
 0
 \end{array}$$

$$11011 = 27$$

8. Val in binary 0011 1111 1110 1000 0000 0000 0000 0000

9. Pipe line allows the execution of instruction in parallel
 Pipe stall occur if the next instruction depends on the result of the previous instruction.

Avoid: A programmer can delay the use of results by

- ① reordering the instructions
- ② Avoid introducing unnecessary branches.
- ③ Delay references to result register.

10. -58

58 : 00111010

-58 : 11000110

78 : 01001110

78	01001110	
+(-58)	11000110	
	<hr/>	
	00010000	= 20.

```
11. int isBigEndian() {
    int k = 8;
    char *p = (char*) &k;
    if (*p == 8) return 0;
    return 1;
}
```

12. Harvard Architecture: has physically separate signals and storage for code and data. It is possible to access program memory and data memory simultaneously.

Von Newman Architecture: have shared signals and memory for code and data. Thus, the program can be easily modified by itself since it is stored in read-write memory.

CISC stand for complex instruction set computer.
 In common, CISC chips are relatively slow (compared to RISC chips) per instruction, but use little (less than RISC) instructions. RISC is cheaper and faster and RISC puts a greater burden on the software.

14. Program Memory (KB) : 32
 RAM (KB) : 2
 Clock speed : 20M
~~#~~ Number of I/O ports : 23
 Number of A/D ports : 6.

```

lds r20, b
lds r21, (b)+1
lds r18, a
lds r19, (a)+1
add r20, r18
abc r21, r19
sts (c)+1, r25
sts c, r24
ref
  
```

```

lds r20, b
lds r21, (b)+1
lds r18, a
lds r19, (a)+1
mul r20, r18
movw r24, r0
mul r20, r19
add r25, r0
mul r21, r18
add r25, r0
clr r1
sts (c)+1, r25
sts c, r24
ref
  
```

16.

```
.text  
.globl addarray  
.type addarray, @function.
```

```
addarray :  
    movl (%esi), %eax  
    movl %edi, %edx
```

```
while: cmpl  
    cmpl $0, %edx  
        jne loop  
    ref
```

```
loop: addl  
    subl $1, %edx  
    add $4, %ediesi  
    add (%esi), %eax  
    jmp while.
```

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```
str1:      .string "%d"
str2:      .string "%d\n"
print1:    .string "a=%d\n"
print2:    .string "b=%d\n"
print3:    .string "max=%d\n"
print4:    .string "min=%d\n"
print5:    .string "avg=%d\n"

.text
.globl main
main:
    pushq   %rbp
    movq   %rsp, %rbp
    subq   $80, %rsp
    movq   $0, -8(%rbp)

    movq   -8(%rbp), %rax
    leaq   -32(%rbp,%rax,8), %rsi
    movq   $str1, %rdi
    call   scanf

    addq   $1, -8(%rbp)

    movq   -8(%rbp), %rax
    leaq   -32(%rbp,%rax,8), %rsi
    movq   $str1, %rdi
    call   scanf

    movq   -32(%rbp), %rsi
    movq   $print1, %rdi
    movq   $0, %rax
    call   printf

    movq   -24(%rbp), %rsi
    movq   $print2, %rdi
    movq   $0, %rax
    call   printf

    movq   -32(%rbp), %rdx
    movq   -24(%rbp), %rcx

    cmpq   %rdx, %rcx
    jge   min
    jmp   max

max:
    movq   -32(%rbp), %rsi
    movq   $print3, %rdi
    movq   $0, %rax
    call   printf

    movq   -24(%rbp), %rsi
    movq   $print4, %rdi
    movq   $0, %rax
    call   printf

    movq   -24(%rbp), %rax
    movq   -32(%rbp), %rdx
    addq   %rdx, %rax
    shrq   $1, %rax
```

```
movq $print5, %rdi
movq %rax, %rsi
movq $0, %rax
call printf
```

```
leave
ret
```

min:

```
movq -24(%rbp), %rsi
movq $print3, %rdi
movq $0, %rax
call printf
```

```
movq -32(%rbp), %rsi
movq $print4, %rdi
movq $0, %rax
call printf
```

```
movq -24(%rbp), %rax
movq -32(%rbp), %rdx
addq %rdx, %rax
shrq $1, %rax
```

```
movq $print5, %rdi
movq %rax, %rsi
movq $0, %rax
call printf
```

```
leave
ret
```