

# Quantum Voting and Bypassing Gibbard-Satterthwaite's Impossibility Theorem

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## Abstract

The Gibbard–Satterthwaite (GS) theorem is a central impossibility result for single-winner voting on ordinal (ranking) ballots: when  $|A| > 2$ , every onto and strategyproof rule is dictatorial. We revisit this no-go phenomenon under a ballot model motivated by indecisive voters and quantum information: a voter may submit not a single ranking, but a state over rankings (a probability mixture, or more generally a density operator on the permutation basis). This raises the key question: how should “preference” and “manipulation” be interpreted when ballots are uncertain/quantum?

We use a projection-and-trace semantics for pairwise comparisons: a ballot’s support for  $x > y$  is  $\text{Tr}(\Pi_{x>y}\rho)$ . This yields sharp and unsharp (0/1 and nonzero) preference types and a corresponding “support-flip” notion of strategic manipulation. Within this framework we formulate a natural Quantum GS conjecture and disprove it by counterexample using Quantum Condorcet Voting (QCV) plus a natural winner readout.

## Classical Baseline: GS on Point Rankings

Classical GS is a theorem about point ballots: each voter reports one linear order  $R_i \in \mathcal{L}(A)$ . A social choice rule maps  $(R_1, \dots, R_n)$  to a winner in  $A$ , and strategyproofness is defined relative to that point domain. For  $|A| > 2$ , “onto + strategyproof” forces dictatorship.

**Our shift:** keep the primitive objects ordinal (rankings) but enlarge the ballot domain from points to distributions/states over  $\mathcal{L}(A)$ . This isolates a concrete source of brittleness in GS: the theorem does not automatically carry over when ballots encode uncertainty at the level of full rankings and when incentives are lifted using a semantics appropriate to that domain.

## Ballot Model: Density Operators over the Ranking Basis

Let  $\mathcal{L}(A)$  be the set of linear orders on  $A$ . Define the ranking Hilbert space

$$R \cong \mathbb{C}^{|\mathcal{L}(A)|}, \quad \{ |R\rangle : R \in \mathcal{L}(A) \} \text{ (ranking basis).}$$

A voter’s ballot is a density operator  $\rho_i \in D(R_i)$ . A ballot profile can be taken as a joint state

$$\rho \in D(R_1 \otimes \dots \otimes R_n),$$

which allows correlated ballots at the level of the formalism.

**Indecisive voting as a restriction.** If we restrict ballots to classical mixtures over basis rankings (equivalently, diagonal/“classical” structure in the ranking basis), we recover “indecisive ballots”: probability distributions over rankings rather than over candidates.

## Support Semantics: Projectors for Pairwise Comparisons

For a pair  $(x, y)$ , let

$$S_{x>y} := \text{span}\{ |R\rangle : x >_R y \}, \quad \Pi_{x>y} \text{ projector onto } S_{x>y}.$$

Define the ballot’s support for  $x > y$  as

$$s_i(x > y) := \text{Tr}(\Pi_{x>y} \rho_i).$$

Interpretation:  $s_i(x > y)$  is the probability mass (more generally, the trace weight) that voter  $i$ ’s ballot assigns to rankings where  $x > y$ .

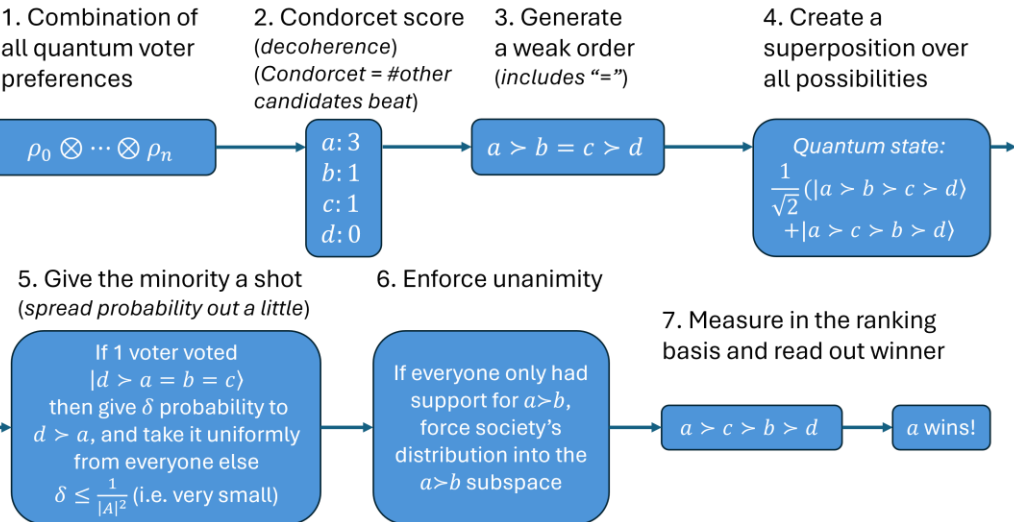
This induces sharp/unsharp preference types used throughout:

- **Strong +:**  $s_i(x > y) = 1$  (ballot supports only  $x > y$ );
- **Strong –:**  $s_i(x > y) = 0$  (ballot forbids  $x > y$ );
- **Weak:**  $s_i(x > y) > 0$  (ballot allows  $x > y$ ).

**Onto** (the relevant notion here): for each candidate  $a$ , there exists a profile for which the social outcome assigns support only to  $a$  as winner (so “ $|A| > 2$  has bite”).

## Open directions

First, strengthen incentives beyond the support semantics (e.g.,  $p$ -threshold preferences or notions of moving the societal ballot “closer” to a voter’s ballot) and ask when a GS-style dictatorship reappears. Second, test generality: do other quantum rules (e.g., Quantum Majority Rule) and other welfare-to-choice readouts behave like QCVNE? Third, map the boundary: what domain/axiom restrictions recover the classical regime, and which other theorems (Sen, Müller–Satterthwaite) or non-ranked ballots (e.g., combined approval) admit analogous translations?



**Fig. 1: QCV Pipeline.** QCV maps ballots  $\rho_1, \dots, \rho_n$  over the ranking basis to a societal ranking state. Two steps drive the GS counterexample: Step 5 (“minority shot”) ensures any pairwise relation with nonzero individual support retains nonzero societal support; Step 6 (“unanimity enforcement”) projects unanimously certain relations to certainty at the social level. A natural top-candidate readout then produces a single-winner distribution from societal support on “ $aa$  is top” subspaces. Theorem X: the resulting rule is onto, QIC (support-flip), and non-dictatorial, refuting QGS for  $|A| > 2$ .

## Incentives: Manipulation as Support Flips

When ballots are uncertain/quantum, the notion of manipulation must specify what counts as a “better” social outcome. We adopt a minimal, qualitative notion aligned with sharp/unsharp semantics.

A voter manipulates if, by misreporting  $\rho_i \rightarrow \rho_i'$ , they can force a socially relevant support flip that improves with respect to their preference type on some pair  $(x, y)$ :

- **Strong + exploit:** raise society from “not certain” to certain ( $< 1 \rightarrow 1$ );
- **Strong– exploit:** reduce society from “some support” to no support ( $> 0 \rightarrow 0$ );
- **Weak exploit:** create support where there was none ( $0 \rightarrow > 0$ ).

A rule is **Quantum Incentive Compatible (QIC)** if no single voter can achieve any beneficial support flip (pairwise, and after winner readout, at the level of winners).

## Conjecture: A Quantum Gibbard–Satterthwaite Analogue

**Conjecture (Quantum Gibbard–Satterthwaite, QGS).** Every QIC voting rule that is onto more than two alternatives ( $|A| > 2$ ) must be a quantum dictatorship (i.e., one voter controls the societal support structure, in the sharp/unsharp sense).

The purpose of QGS is to test whether the GS dictatorship conclusion survives once ballots are states over rankings and incentives are interpreted via support.

## Counterexample: Quantum Condorcet Voting

The counterexample comes from Quantum Condorcet Voting (QCV), which is naturally described as a protocol on ranking states. Fig. 1 shows the pipeline, but only two steps matter for incentives:

**Minority shot (Step 5):** if any voter has  $s_i(x > y) > 0$ , the societal state is forced to have  $\text{Tr}(\Pi_{x>y}\rho_{soc}) > 0$ .

**Unanimity enforcement (Step 6):** if all voters have  $s_i(x > y) = 1$ , the societal state is projected so  $\text{Tr}(\Pi_{x>y}\rho_{soc}) = 1$ .

Any attempted misreport that would change support for a pairwise relation is blocked by either step 5 or 6. To obtain a single-winner rule, apply the natural top-candidate readout. (This is the “Natural Extension” in the paper.)

## Theorem 3.4: Disproof of the QGS Conjecture

For  $|A| > 2$ , the voting rule obtained by QCV plus the natural top-candidate readout is simultaneously:

1. QIC (no beneficial support-flip manipulation),
2. onto, and
3. non-dictatorial (sharp and unsharp).

Hence the Quantum GS conjecture is false. Restricting ballots to classical mixtures yields the same conclusion for indecisive voting.

## References & Contact

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