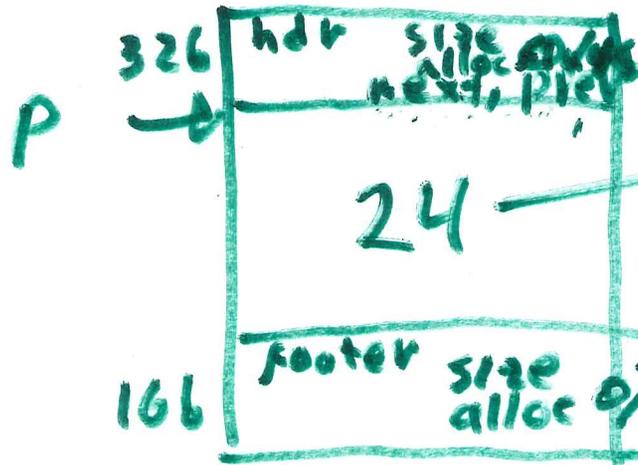


Lab 1 malloc

①

```
int MemSize = 20;  
char *p = malloc(MemSize);  
= block
```



Memory aligned size =
 $(\text{MemSize} + 7) \& \sim 7;$

Memory size is aligned to next 8 byte boundary.

that is because a long or double have to be stored in 8 byte boundary.

free(p)

- We need the header to store the size to know how many bytes to return to malloc free list.

Header (32 bytes)

size - size of block including header and footer (8 bytes)

alloc flags - 0 = free 1 = alloc 2 = sentinel (8 bytes)

next, prev - Used when object is free, to add to free list.
(8 + 8) = 16 bytes

Footer (16 bytes)

(2)

Size - Size of block including header + footer (8 bytes)

Alloc flags - 0 = free 1 = alloc 2 = sentinel (8 bytes)

The smallest object is

data MemSize = 1 byte \rightarrow 8 bytes aligned sizes

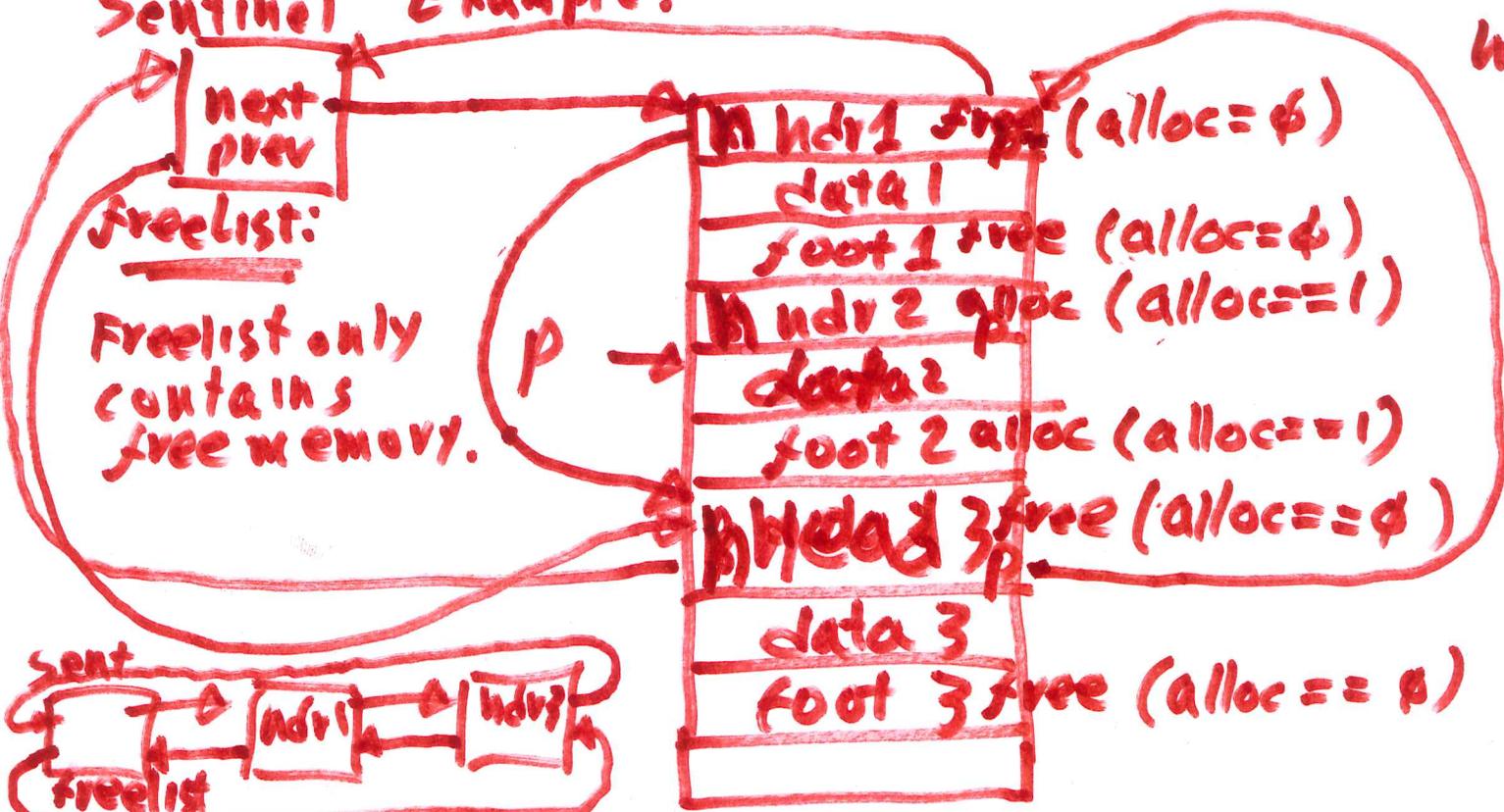
Header 32 bytes \rightarrow 32

Footer 16 bytes \rightarrow 16

56 bytes

The footer is used to coalesce blocks when they are freed.

Sentinel Example:



hdr & footer are also called "boundary tags".

- We have two separate data structures
- Boundary tags
 - Double linked list.

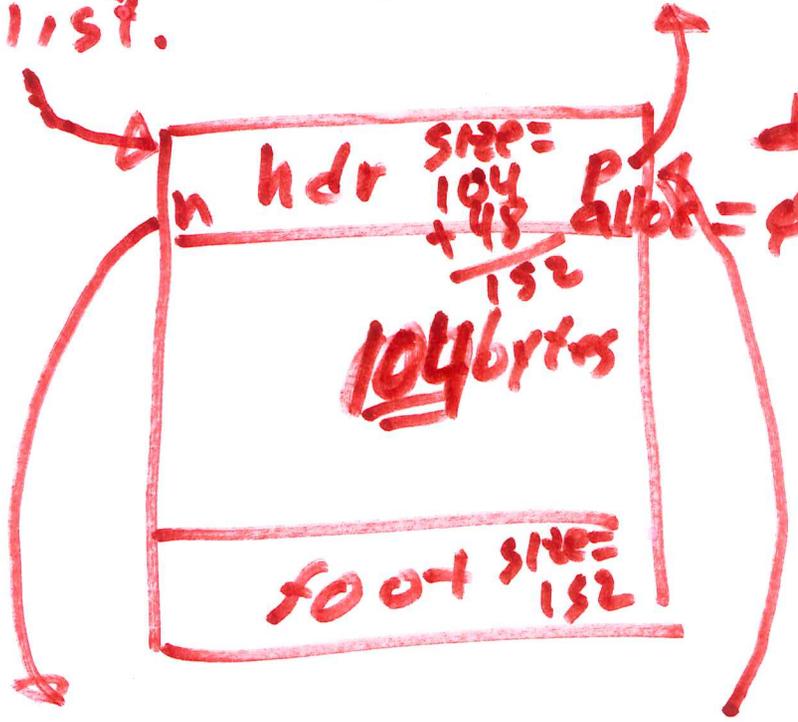
`free(p)`

Check if object before `p` is free. If free then coalesce it.

p = malloc (memSize)

- Find first block large enough to satisfy request.
- Try split object. If remainder is smaller than smallest usable object i.e. $\text{sizeof(Header)} + \text{sizeof(footer)} + 8$, then don't split and return whole object. else split object return remainder to freelist.

p = malloc(20)



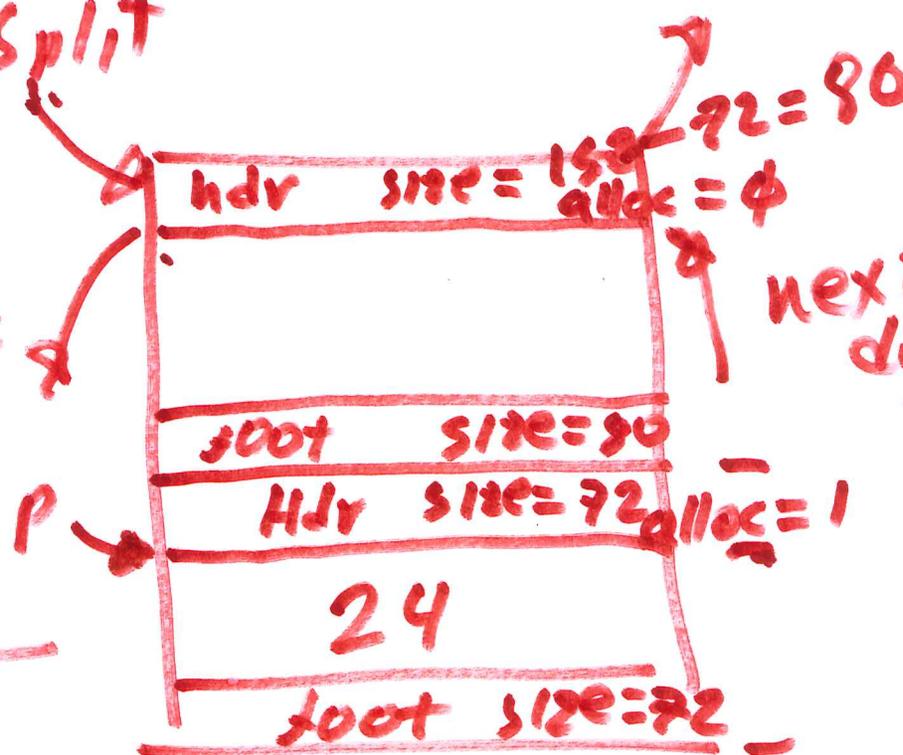
object that needs to be split

After Split

(5)

Mem Aligned =
 24
 + Header 32
 + Foot 16

 72



next & prev
 do not need
 to be change
 if we return
 the lower
 part after
 split.

6

free(void *p) {

Header *hdr =

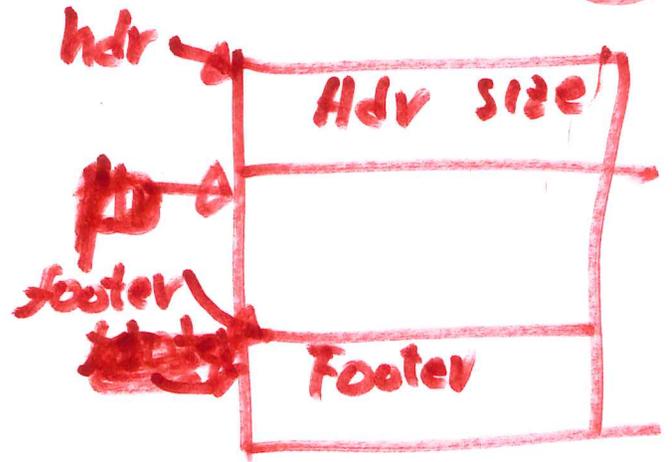
(Header *) ((char *) p

- sizeof(Header))

Footer *footer =

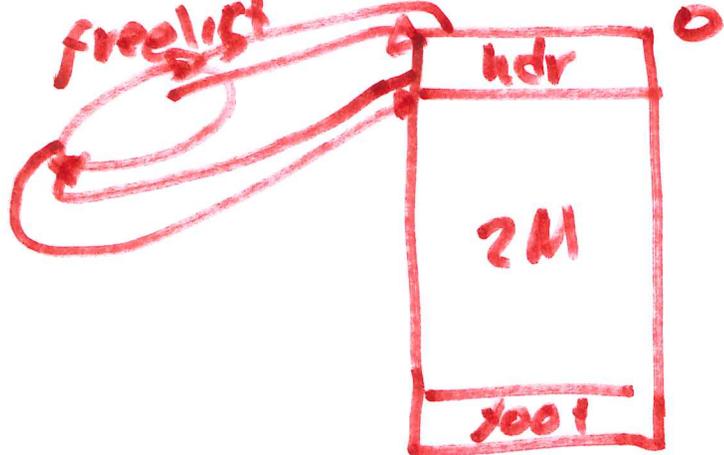
(Footer *) ((char *) hdr

+ hdr->size - sizeof(Footer))

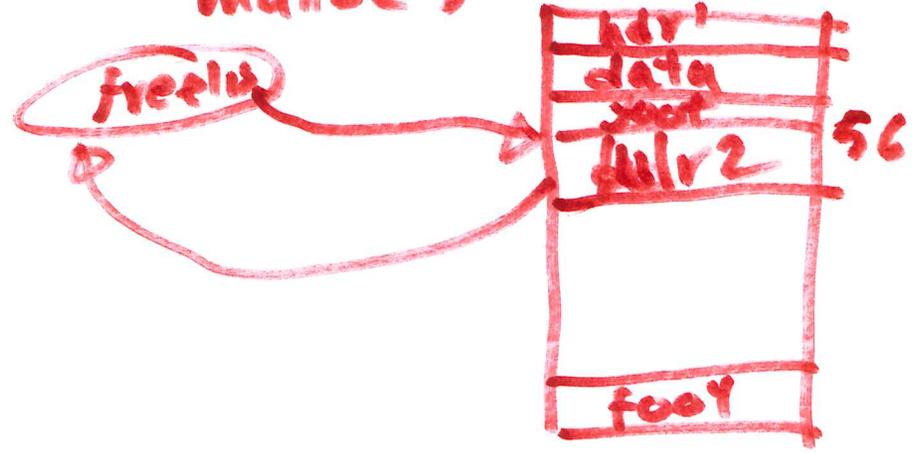


4

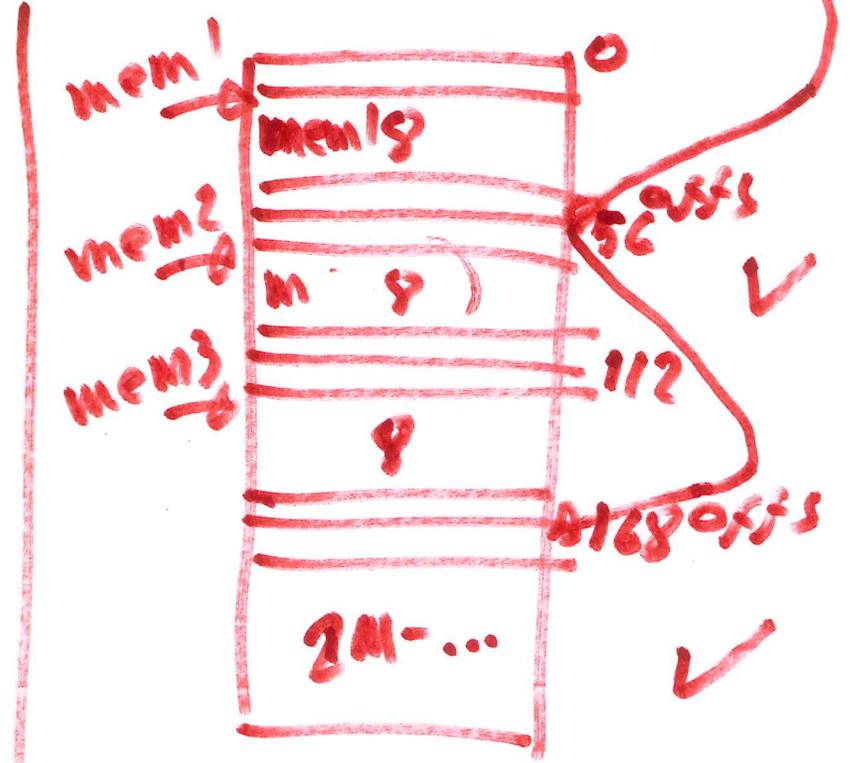
Initially
freelist



malloc 8



freelist



free(mem2)