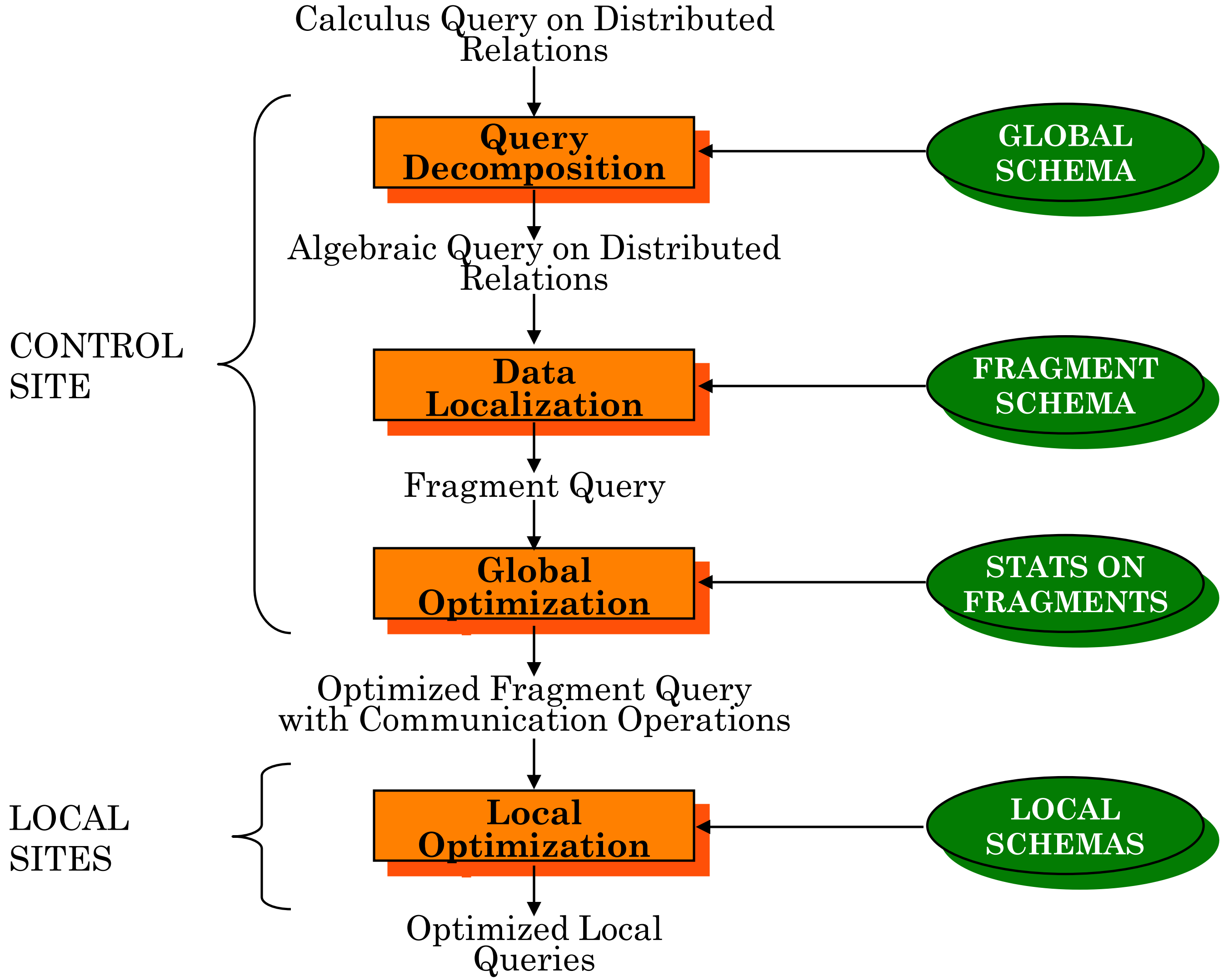
Outline

- Introduction
- Background
- Distributed DBMS Architecture
- Distributed Database Design
- Distributed Query Processing
  - Query Processing Methodology
  - Distributed Query Optimization
- Distributed Transaction Management (Extensive)
- Building Distributed Database Systems (RAID)
- Mobile Database Systems
- Privacy, Trust, and Authentication
- Peer to Peer Systems

Useful References

- Textbook *Principles of Distributed Database Systems*, Chapter 7

Distributed Query Processing Methodology

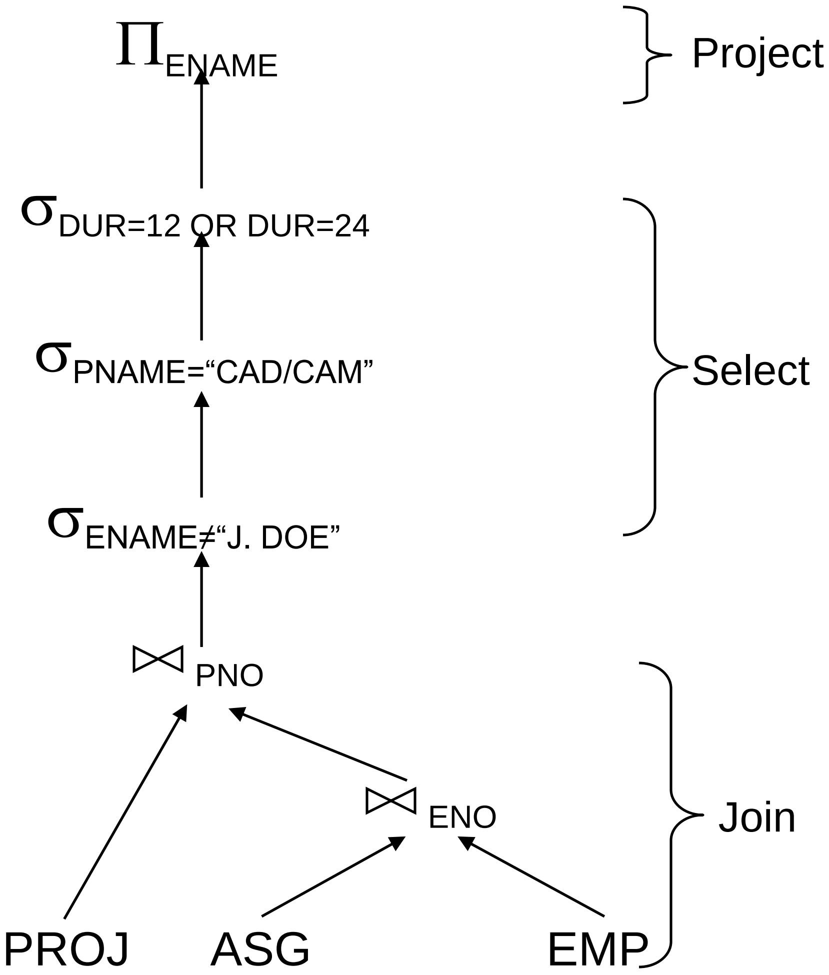


Restructuring

- Convert relational calculus to relational algebra
- Make use of query trees
- Example

  Find the names of employees other than J. Doe who worked on the CAD/CAM project for either 1 or 2 years.

```
SELECT   ENAME
FROM     EMP, ASG, PROJ
WHERE    EMP.ENO = ASG.ENO
AND      ASG.PNO = PROJ.PNO
AND      ENAME ≠ "J. Doe"
AND      PNAME = "CAD/CAM"
AND      (DUR = 12 OR DUR = 24)
```



Restructuring – Transformation Rules

- Commutativity of binary operations
  - $R \times S \Leftrightarrow S \times R$
  - $R \bowtie S \Leftrightarrow S \bowtie R$
  - $R \cup S \Leftrightarrow S \cup R$
- Associativity of binary operations
  - $(R \times S) \times T \Leftrightarrow R \times (S \times T)$
  - $(R \bowtie S) \bowtie T \Leftrightarrow R \bowtie (S \bowtie T)$
- Idempotence of unary operations
  - $\Pi_{A'}(\Pi_{A''}(R)) \Leftrightarrow \Pi_{A'}(R)$
  - $\sigma_{p_1(A_1)}(\sigma_{p_2(A_2)}(R)) = \sigma_{p_1(A_1) \wedge p_2(A_2)}(R)$

  where $R[A]$ and $A' \subseteq A$, $A'' \subseteq A$ and $A' \subseteq A''$
- Commuting selection with projection

Restructuring – Transformation Rules

- Commuting selection with binary operations
  - $\sigma_{p(A)}(R \times S) \Leftrightarrow (\sigma_{p(A)}(R)) \times S$
  - $\sigma_{p(A_i)}(R \bowtie_{(A_j,B_k)} S) \Leftrightarrow (\sigma_{p(A_i)}(R)) \bowtie_{(A_j,B_k)} S$
  - $\sigma_{p(A_i)}(R \cup T) \Leftrightarrow \sigma_{p(A_i)}(R) \cup \sigma_{p(A_i)}(T)$

  where A_i belongs to R and T
- Commuting projection with binary operations
  - $\Pi_C(R \times S) \Leftrightarrow \Pi_{A'}(R) \times \Pi_{B'}(S)$
  - $\Pi_C(R \bowtie_{(A_j,B_k)} S) \Leftrightarrow \Pi_{A'}(R) \bowtie_{(A_j,B_k)} \Pi_{B'}(S)$
  - $\Pi_C(R \cup S) \Leftrightarrow \Pi_C(R) \cup \Pi_C(S)$

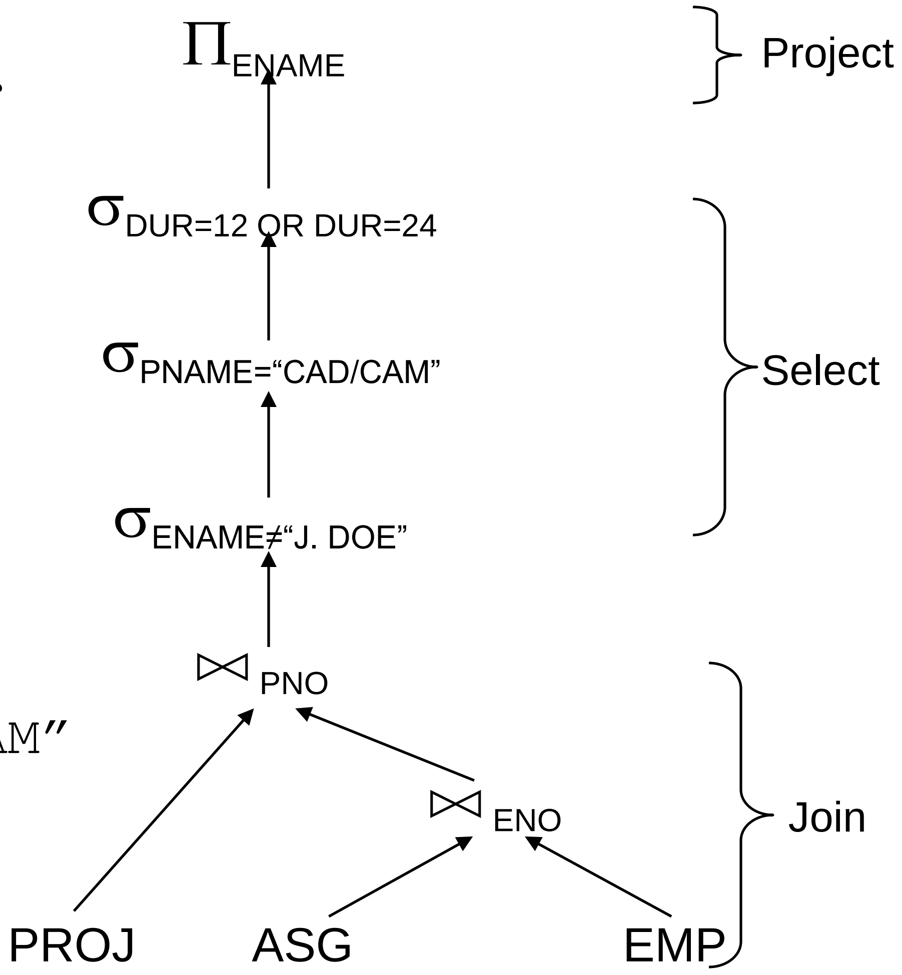
  where $R[A]$ and $S[B]$; $C = A' \cup B'$ where $A' \subseteq A$, $B' \subseteq B$

Example

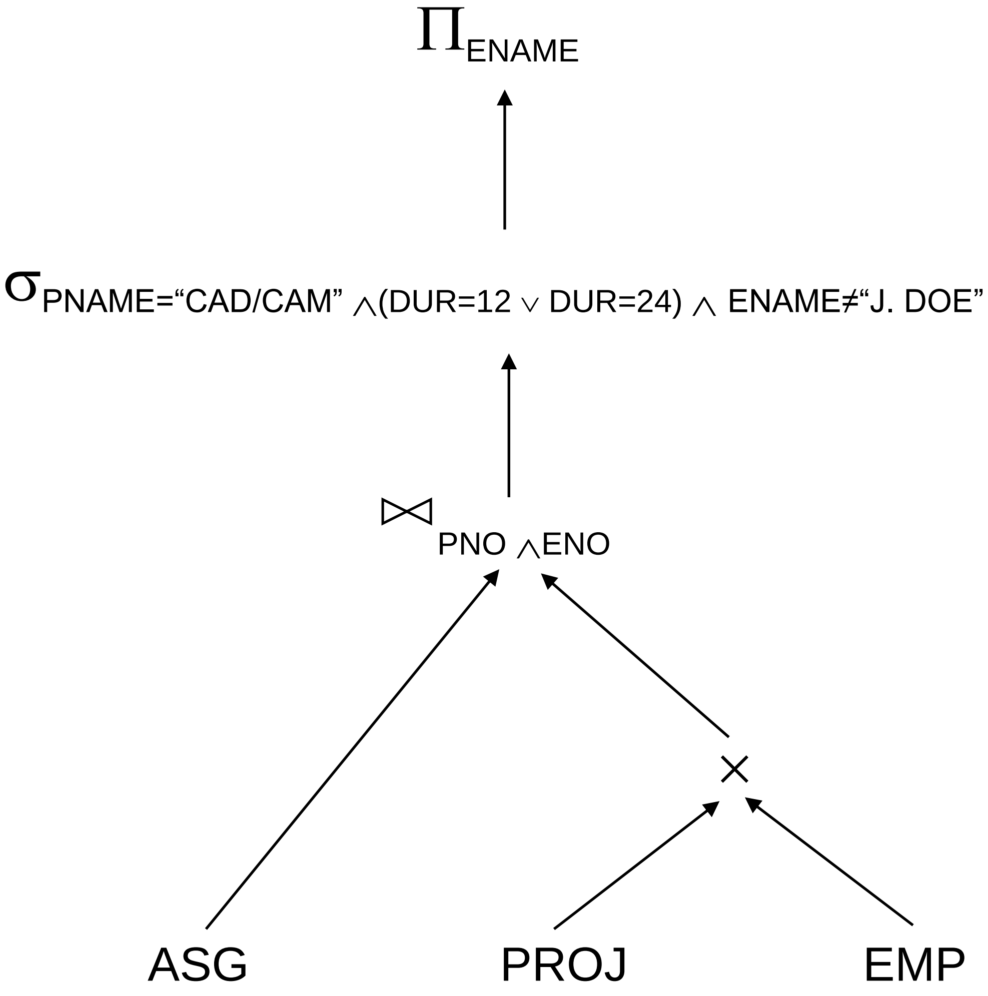
Recall the previous example:

Find the names of employees other than J. Doe who worked on the CAD/CAM project for either one or two years.

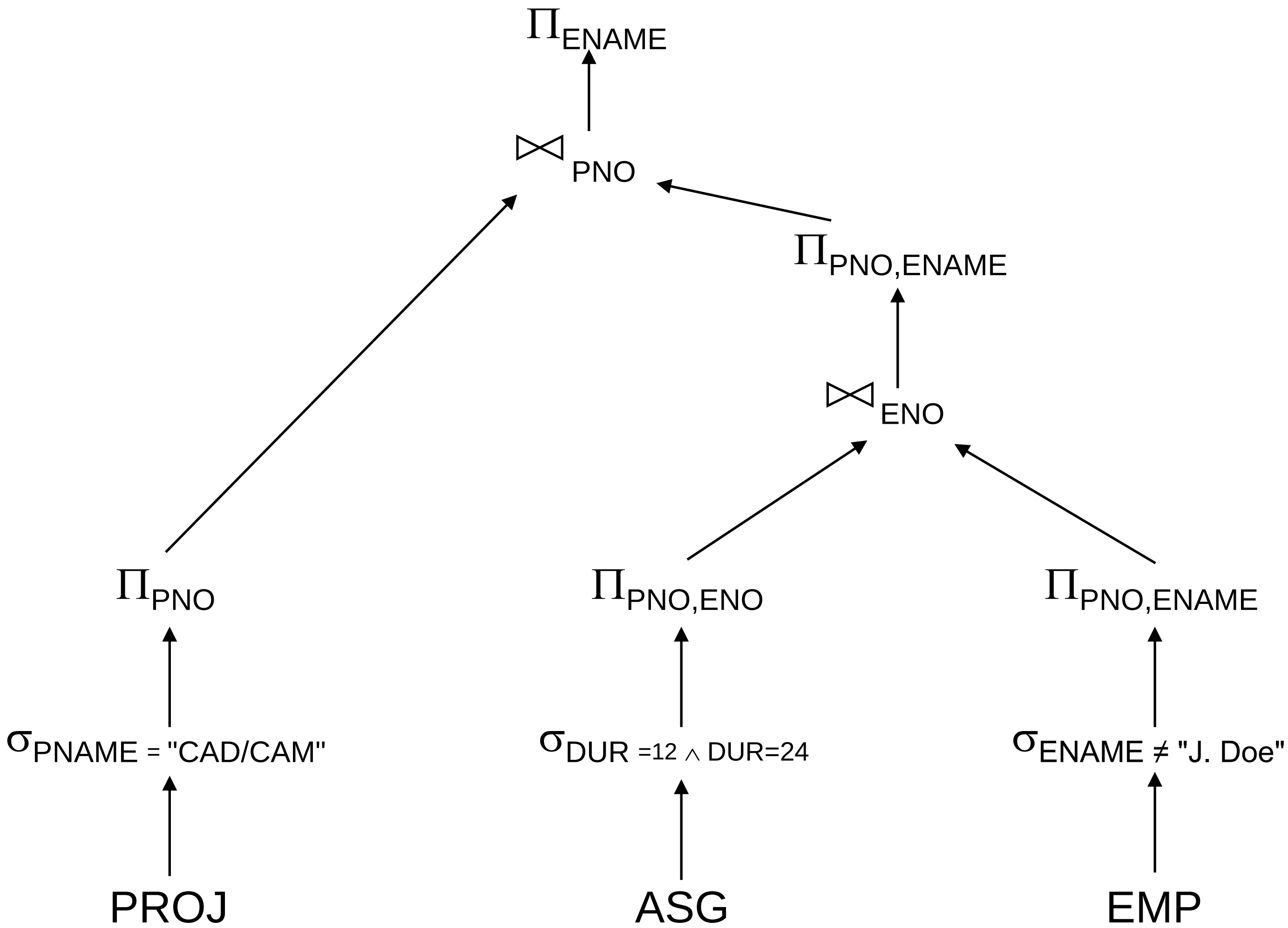
```
SELECT   ENAME
FROM     PROJ, ASG, EMP
WHERE    ASG.ENO=EMP.ENO
AND      ASG.PNO=PROJ.PNO
AND      ENAME≠"J. Doe"
AND      PROJ.PNAME="CAD/CAM"
AND      (DUR=12 OR DUR=24)
```



Equivalent Query



Restructuring



Cost Functions

- Total Time (or Total Cost)
  - Reduce each cost (in terms of time) component individually
  - Do as little of each cost component as possible
  - Optimizes the utilization of the resources

$\Downarrow$

Increases system throughput

- Response Time
  - Do as many things as possible in parallel
  - May increase total time because of increased total activity

Total Cost

Summation of all cost factors

Total cost = CPU cost + I/O cost + communication cost

CPU cost = unit instruction cost ∗ no.of instructions

I/O cost = unit disk I/O cost ∗ no. of disk I/Os

communication cost = message initiation + transmission

Total Cost Factors

- Wide area network
  - message initiation and transmission costs high
  - local processing cost is low (fast mainframes or minicomputers)
  - ratio of communication to I/O costs = 20:1
- Local area networks
  - communication and local processing costs are more or less equal
  - ratio = 1:1.6

Response Time

Elapsed time between the initiation and the completion of a query

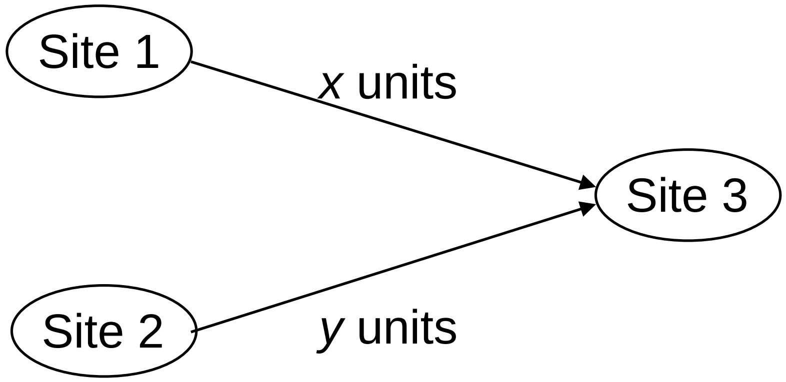
Response time = CPU time + I/O time + communication time

CPU time = unit instruction time ∗ no. of sequential instructions

I/O time = unit I/O time ∗ no. of sequential I/Os

communication time = unit msg initiation time ∗ no. of sequential msg + unit transmission time ∗ no. of sequential bytes

Example



Assume that only the communication cost is considered

Total time = 2 $*$ message initialization time + unit transmission time $*$ $(x+y)$

Response time = max {time to send x from 1 to 3, time to send y from 2 to 3}

time to send x from 1 to 3 = message initialization time + unit transmission time $*$ x

time to send y from 2 to 3 = message initialization time + unit transmission time $*$ y