Question 1. (1.5 point)

Consider the following schedules, where R, W, C stands for ‘Read’, ‘Write’ and ‘Commit’ respectively and subscripts refers to transactions (i.e., \( R_i(X) \) means transaction \( T_i \) read \( X \)):

**Schedule 1:** \( R_1(X), W_1(X), R_3(X), W_3(X), R_1(Y), W_1(Y), C_1, W_3(Y) C_3, R_2(Y), W_2(Y) C_2 \)

**Schedule 2:** \( R_2(X), W_2(X), R_3(Y), W_3(Y), R_3(Z), W_3(Z), C_3, R_2(Z), W_2(Z), C_2, R_1(X), W_1(X), C_1 \)

**Schedule 3:** \( W_1(A), W_2(B), W_3(C), R_1(X), R_2(X), R_1(Y), W_1(X), C_1, W_2(Y), C_2, W_3(Y), C_3 \)

For each schedule answer the following question:

(a) Is this schedule conflict serializable? If yes, provide an equivalent serial schedule of transactions.
(b) Does the schedule avoid cascading aborts? Explain your answer.
(c) Assume 2PL protocol where operations are skipped if the transaction cannot acquire a lock and deadlocks are handled using a wait-for graph, like Project 4. Furthermore, once a transaction commits, all the locks held are released and the operations that were on hold are retried – according to the priority in the schedule. Is the schedule possible under this 2PL protocol (i.e. all transactions can be committed)? Explain your answer. If the 2PL protocol leads to a deadlock, give the wait-for graph. Note: operations that were skipped will be re-tried later.

Question 2. (1.5 point)

(a) Discuss the different types of transaction failures. Explain each failure.
(b) What is the system log used for?
(c) What are checkpoints and why are they important?