# Purdue University Computer Science Department CS44800 Homework 4 Fall 2019

## Due: November 11, 2019, 11:59PM

Reference:

## • Fundamentals of Database Systems by Elmasri & Navathe, 7<sup>th</sup> Edition Total Points: **3 points** Submit via Blackboard

Provide concise but to the point explanation to support your answers. If possible, use bullets to organize the ideas in your answer.

## Question 1. (0.5 points)

Consider the universal relation  $R = \{A, B, C, D, E, F, G, H, I, J\}$  and the set of functional dependencies F.

$$F = \{\{A, B\} \rightarrow \{C, D\}, \\ \{D\} \rightarrow \{E, F, G\}, \\ \{F, G\} \rightarrow \{H\}, \\ \{A\} \rightarrow \{I\}, \\ \{A, B\} \rightarrow \{E, G\}, \\ \{A, I\} \rightarrow \{I, J\}\}\}$$

- (a) What is the closure of the set {A, D}?
- (b) What is the key for R?
- (c) What is the minimum cover for the set of functional dependencies F?
- (d) Decompose R into 2NF.
- (e) Decompose R into 3NF.

# Question 2. (0.75 points)

- (a) What are the possible sources of the information that defines the functional dependencies that hold among the attributes of a relation schema? Discuss two possible sources. Explain your answer.
- (b) What is meant by the closure of a set of functional dependencies? Illustrate with an example.
- (c) What is the lossless (or nonadditive) join property of a decomposition? Why is it important?

#### **Question 3. (0.75 points)**

- (a) Discuss the four levels of isolation in SQL.
- (b) Discuss the types of violation that isolation levels prevent. Give an example for each violation.
- (c) Discuss the concept of snapshot isolation and its effect on the phantom read problem.

## **Question 4. (0.50 points)**

Consider the following four schedules for transactions T1, T2 and T3:

a. r1(X); r3(X); w1(X); r2(X); w3(X); w2(X)
b. r1(Y); r3(Y); w3(X); w1(X); r2(X); w2(Y)
c. w2(X); r1(Y); w2(Y); r2(X); w2(Y); r3(Y)
d. r3(X); r2(X); w3(X); r1(Y); w1(X); r2(Y)
e. r3(X); r2(X); r1(X); w3(X); w1(X); w2(X)

- (a) For each schedule, specify whether or not the schedule is conflict serializable and why. To determine if a schedule is conflict serializable you can either use the precedence graph or swap conflicting operations.
- (b) Then, for each serializable schedule, give the equivalent serial schedule.

## **Question 5. (0.50 points)**

Consider the following actions taken by transaction T1 on a database with objects X and Y:

## **R1(X), W1(X), R1(Y), W1(Y)**

Without altering the order of operations of T1, give example schedules with an additional Transactions T2 (with 4 operations accessing both database objects X and Y) that fulfill the following conditions:

- (a) Give an example schedule using two transactions (T1 and T2) which is conflict serializable.
- (b) Give an example schedule using two transactions (T1 and T2) which is not conflict serializable. Explain why the schedule is not conflict serializable.