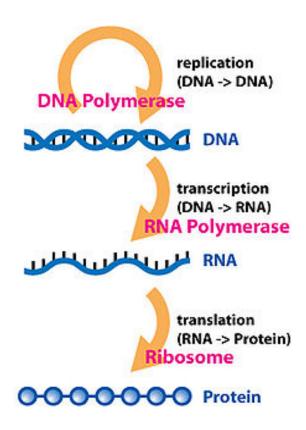
Introduction to Modeling and Algorithms in Life Sciences

Ananth Grama
Purdue University
http://www.cs.purdue.edu/homes/ayg

Acknowledgements

- To various sources, including Prof. Michael Raymer, Wiki sources (pictures), and other noted attributions.
- To the US National Science Foundation and the Center for Science of Information.

Central Dogma of Molecular Biology



Central Dogma of Molecular Biology

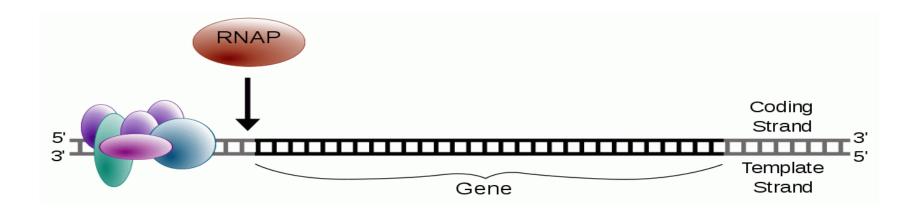
- Mostly valid with some exceptions:
 - Reverse Transcription: Retroviruses such as Feline Leukemia, HIV
 - RNA Replication: RNA to RNA transfer in viruses
 - Direct Translation: DNA to Protein (typically in cell fragments)

Protein Synthesis

- Transcription: a DNA molecule is converted into a complementary strand of RNA
- This RNA is also called messenger RNA (mRNA) since it acts as an intermediary between DNA and the Ribosomes
- Ribosomes are parts of cell that synthesize proteins from mRNA

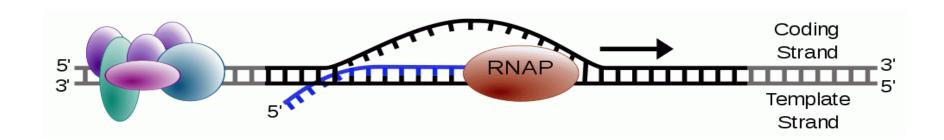
Transcription: Initiation

- Transcription is affected by an enzyme called RNA Polymerase (RNAP)
- RNAP binds to core promoters (in the presence of transcription factors) and pre-initiates the transcription process
- The essentials for pre-initiation include the core promoter sequence, transcription factors, DNA Helicase, RNA Polymerase, and activators/ repressors.



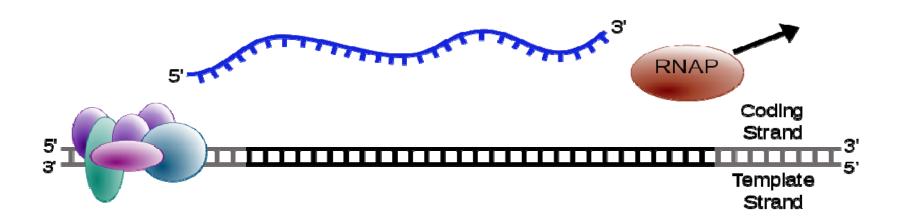
Transcription: Elongation

- One strand of DNA is used as a template for RNA synthesis
- RNAP traverses the template, using base-pair complementarity to synthesize RNA

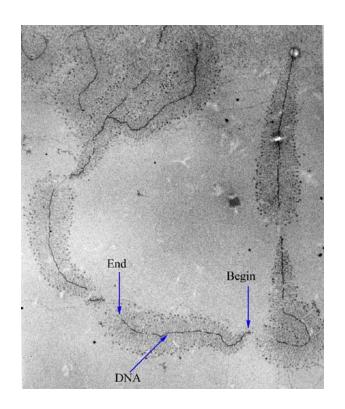


Transcription: Termination

 The transcript is cleaved (typically induced by a termination sequence on the template) followed by a template-independent addition of As at its new 3' end, in a process called polyadenylation.

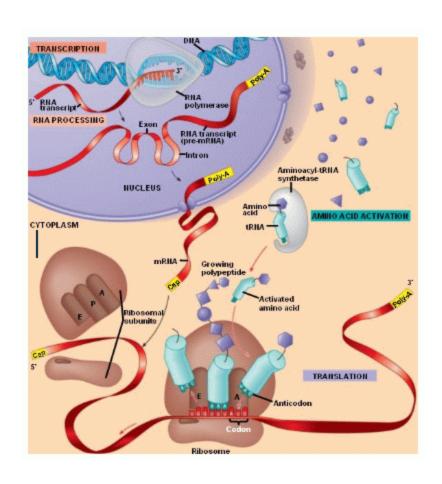


Transcription: Snapshot



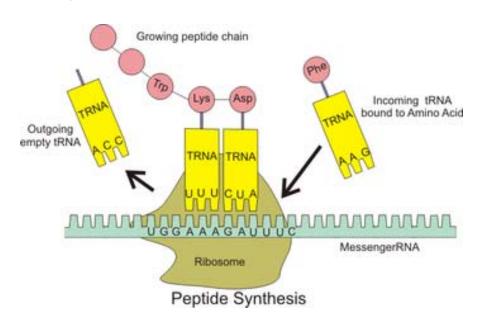
Micrograph of gene transcription of ribosomal RNA illustrating the growing primary transcripts. "Begin" indicates the 5' end of the coding strand of DNA, where new RNA synthesis begins; "end" indicates the 3' end, where the primary transcripts are almost complete. [Trepte, 2005]

Eukaryotic Transcription



Synthesizing Proteins: Translation

- mRNA is decoded by the Ribosome to produce specific proteins (polypeptide chains)
- Polypeptide chains fold to make active proteins
- The amino acids are attached to <u>transfer RNA</u> (tRNA)
 molecules, which enter one part of the ribosome and bind to the
 messenger RNA sequence.



- Each combination of 3 nucleotides on mRNA is called a <u>codon.</u>
- Each codon specifies a particular <u>amino acid</u> that is to be placed in the polypeptide chain.
- Using the mRNA as a template, the ribosome traverses each codon of the mRNA, pairing it with the appropriate amino acid provided by a tRNA. Molecules of <u>transfer RNA</u> (tRNA) contain a complementary <u>anticodon</u> on one end and the appropriate amino acid on the other.

Outgoing

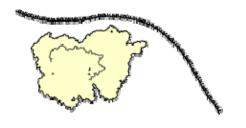
Growing peptide chain

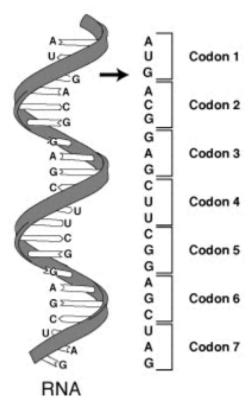
Ribosome

Peptide Synthesis

Incoming tRNA bound to Amino Acid

MessengerRNA





Ribonucleic acid

- Since three mRNA nucleotides form a codon, and each nucleotide can be selected from among four bases, there are 64 possible combinations
- Of there, three are stop codons: UAA, UAG, UGA
- The codon AUG serves as a start codon, in addition to coding the amino acid methionine
- Since 61 codons (64 3) code 20 aminoacids, there is considerable redundancy in the code.

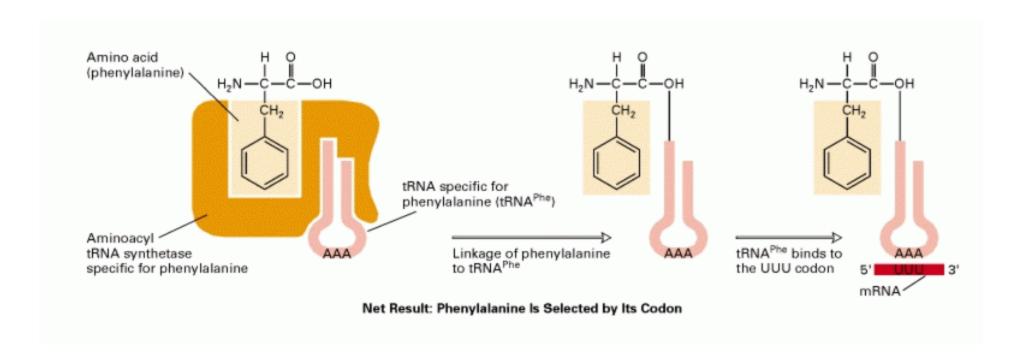
1st base

		U		С		Α		G		
2nd base	U	UUU UUC UUA UUG	Phenylalanine Phenylalanine Leucine Leucine	UCU UCC UCA UCG	Serine Serine Serine Serine	UAU UAC UAA UAG	Tyrosine Tyrosine Stop Stop	UGU UGC UGA UGG	Cysteine Cysteine Stop Tryptophan	U C A G
	С	CUU CUC CUA CUG	Leucine Leucine Leucine Leucine	CCU CCC CCA CCG	Proline Proline Proline Proline	CAU CAC CAA CAG	Histidine Histidine Glutamine Glutamine	CGU CGC CGA CGG	Arginine Arginine Arginine Arginine	U C A G
	A	AUU AUC AUA AUG	Isoleucine Isoleucine Isoleucine Methionine (Start)	ACU ACC ACA ACG	Threonine Threonine Threonine Threonine	AAU AAC AAA AAG	Asparagine Asparagine Lysine Lysine	AGU AGC AGA AGG	Serine Serine Arginine Arginine	U C A G
	G	GUU GUC GUA GUG	Valine Valine Valine Valine	GCU GCC GCA GCG	Alanine Alanine Alanine Alanine	GAU GAC GAA GAG	Aspartic Acid Aspartic Acid Glutamic Acid Glutamic Acid	GGU GGC GGA GGG	Glycine Glycine Glycine Glycine	U C A G

Nonpolar, aliphatic Polar, uncharged

Aromatic

Positively charged Negatively charged



- Amino-acids are bonded through a covalent peptide bond
- The carboxyl group of one molecule reacts with the amino group of the other molecule, thereby releasing a molecule of water (H₂O)

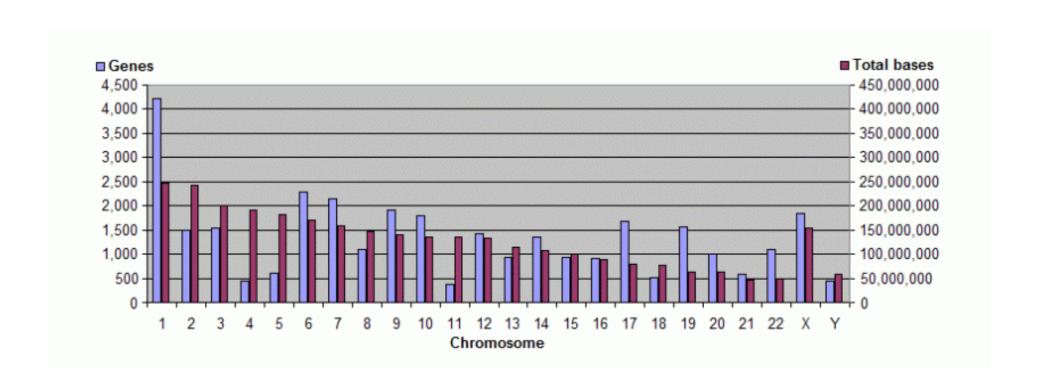
$$R$$
 OH + H_2N R' H_2O

Some Numbers

Human DNA has:

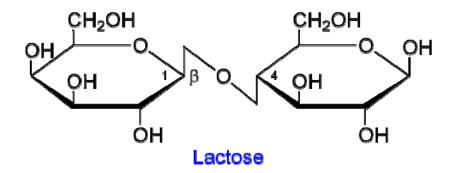
- 3 billion base pairs
- The length of DNA in a cell is 1m!
- This is packed into a nucleus of 3 10 microns
- Each chromosome (46 in all) is about 2 cm on average.

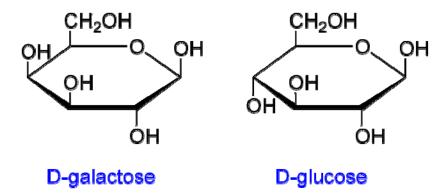
Some Numbers



Putting it all Together

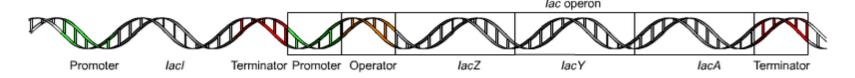
- Models for Lactose Intolerance
 - Roughly 80% of people (about 4B) have varying levels of lactose intolerance
 - These people are unable to break down lactose to smaller sugars
 - Lactose passes into the intestines, where it is broken down by the resident bacteria
 - This results in acute intestinal discomfort





- beta-galactosidase (LacZ), an intracellular enzyme (protein) cleaves the disaccharide lactose into glucose and galactose
- These smaller molecules can be metabolized within the cell.

- Beta-galactosidase is expressed by the lacZ gene
 - lacZ is part of the lac operon, comprised of lacZ, lacY, and lacA genes in E. Coli
 - lacY encodes #-galactoside permease (LacY), a membrane-bound transport protein that pumps lactose into the cell.
 - lacA encodes #-galactoside transacetylase (LacA), an enzyme that transfers an acetyl group from acetyl-CoA to beta-galactosides.
- So, should beta-galactosidase be expressed all the time?
 Would this be an effective use of limited cellular resources?



- It was experimentally observed that when the bacteria is starved of lactose, very low levels of betagalactosides are observed in the cell.
- Conversely, when the <u>level of glucose in the cell is</u> <u>low, high levels of beta-galactosides are observed.</u>
- This must imply <u>other forms of control</u> on the expression of genes!

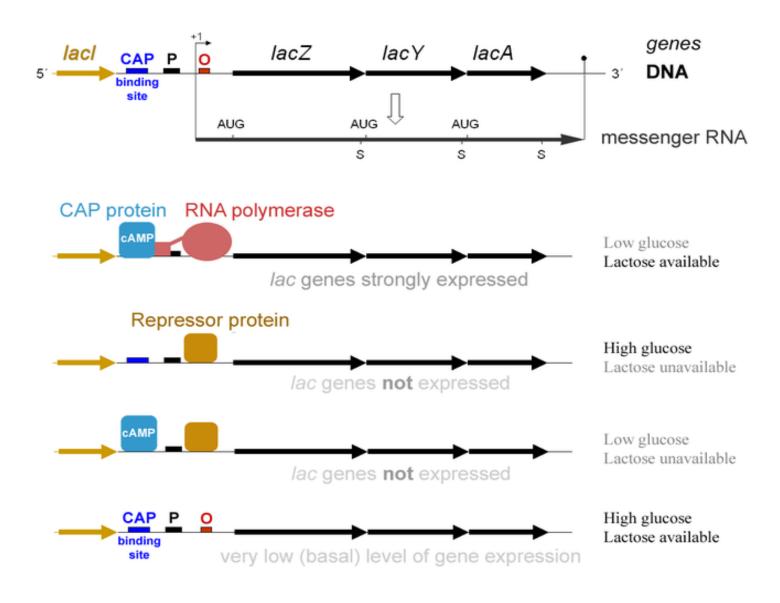
- Negative control: When levels of lactose are low, lacZ expression should be supressed.
- Negative control is affected by the lacl gene (a few hundred basepairs upstream from the lac operon).
- Recall that the ribosome is responsible for transcription.
- In E.Coli, this ribosome consists of a complex of five proteins.
- One of these proteins recognizes a promoter region for the lac operon, causing the ribosome to bind to the DNA and initiate transcrption.

- Negative control is exerted by the lacl gene (few hundred bases upstream of the lac operon).
- lacl constantly (and in very low concentrations) expresses the lac repressor protein.
- lac repressor preferentially binds to the promoter site for the lac operon, preventing the RNAP from binding and transcribing the lac operon.
- In this way, no beta-galactoside is expressed.

- When lactose is present, it binds to the lac repressor protein, causing a change in its conformation
- This prevents it from binding to the promoter site for the lac operon
- This results in the RNAP binding and transcribing the lac operon!

- Positive Feedback: when glucose levels in the cell are low, we want to generate significant amounts of beta-galactosidase, so as to be able to break down any amount of lactose present.
- Cyclic adenosine monophosphate (cAMP) is a signal molecule whose prevalence is inversely proportional to that of glucose.
- cAMP binds to the the Catabolite activator protein (CAP)
- This allows CAP to bind to the CAP binding site upstream of the promoter for the lac operon.
- This in turn facilitates the RNAP to bind and transcribe the operon.

The lac Operon and its Control Elements

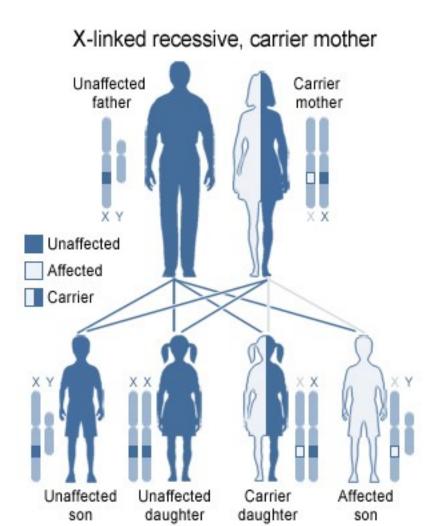


- So what causes lactose intolerance anyway?
 - The lack of beta-galactoside in cells.
- What does one do about it?
 - Take lactase enzymes (beta-galactosides) before lactose!

One more example: Haemophilia

- Genetic disorder that impairs blood coagulation
- Haemophilia A is the most common form, effecting 1 in 5000 to 10000 male births
- Haemophilia A is a recessive X-linked genetic disorder involving a lack of functional clotting Factor VIII and represents 80% of haemophilia cases.
- In females, since there are two X chromosomes, a defect on one of the chromosomes may not manifest itself in the form of a disease.

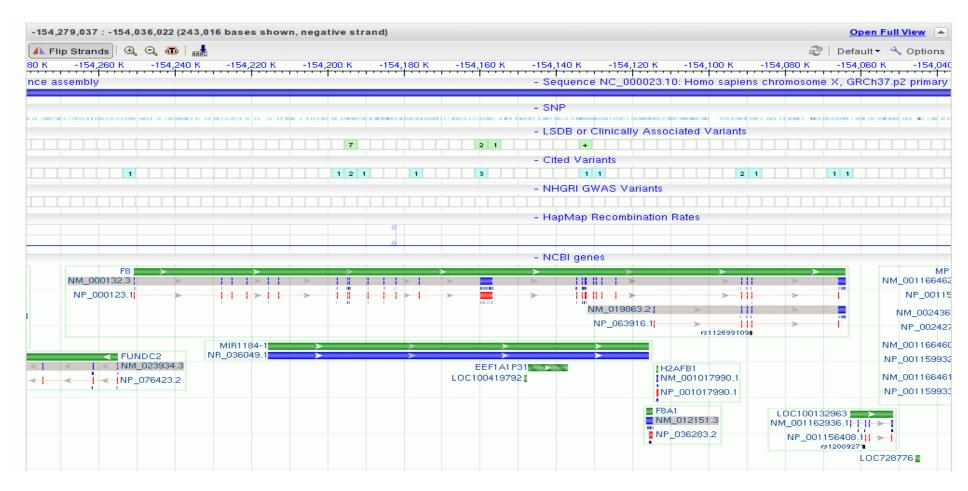
Haemophilia-A



U.S. National Library of Medicine

Haemophilia-A

 Haemophilia-A is caused by clotting factor VIII deficiency. Factor VIII is encoded by the F8 gene.



Haemophilia-A

- Through a sequence of complexes and cofactors, Factor VIII results in the generation of fibrin, which causes the clot.
- Factor VIII, concentrated from plasma is given to haemophiliacs to restore haemostasis.

What we have learnt thus far?

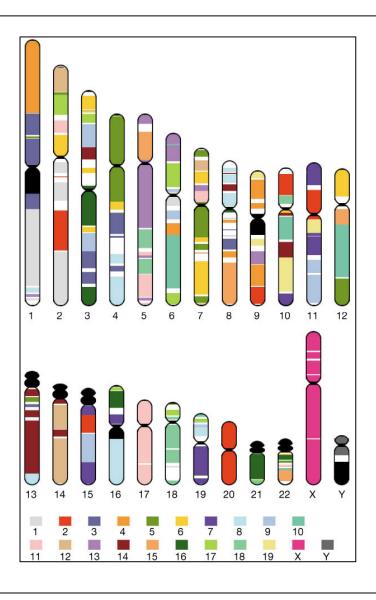
- Central dogma of biology
- The role of sequences (DNA, RNA, proteins)
- The ability to translate across DNA, RNA, and proteins
- The role of genes and proteins in disorders/ functions.

Analyzing Sequences

Sequences: An Evolutionary Perspective

- Evolution occurs through a set of modifications to the DNA
- These modifications include point mutations, insertions, deletions, and rearrangements
- Seemingly diverse species (say mice and humans) share significant similarity (80-90%) in their genes
- The locations of genes may themselves be scrambled

Chromosomal Rearrangements



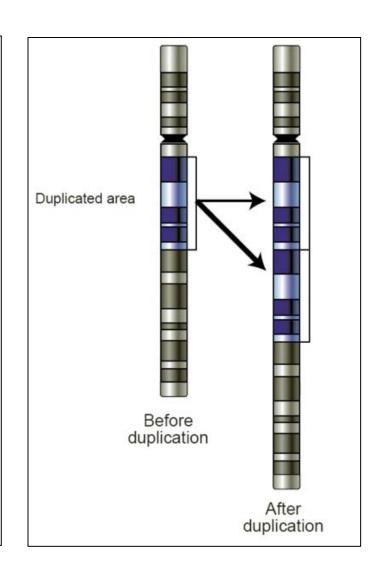
Mouse genome mappings to human genome.

Mouse Genome

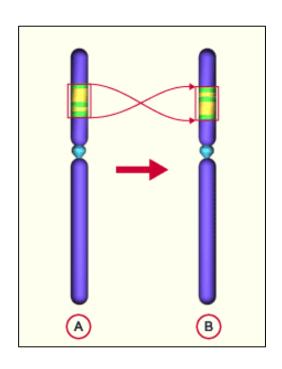
- Mouse genome 2.5 Gb vs human 2.9 Gb
- Can identify <u>regions of synteny</u> between mouse and human for 90% of genome.
- Both genomes have ~30,000 genes
- 99% of mouse genes have a human homolog (and vice versa)
- Some genes appear to have evolved more quickly than random chance (immunity and reproduction).

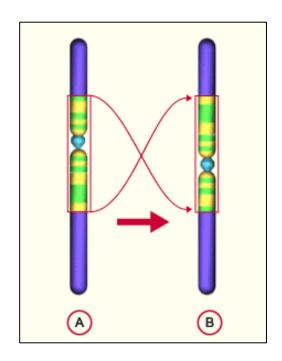
Gene Duplication

- Gene duplication has important evolutionary implications
- Duplicated genes are not subject to evolutionary pressures
- Therefore they can accumulate mutations faster (and consequently lead to specialization)



Inversions

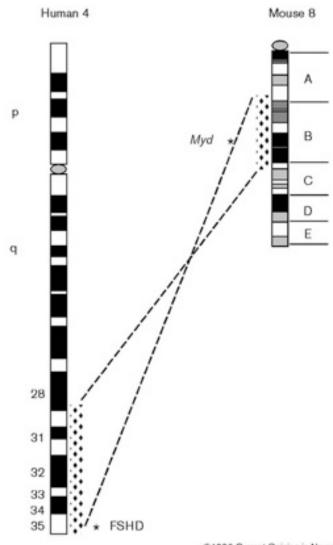




Para and pericentric inversions

Transposition

A group of conserved genes appears in a transposed fashion at a different location



©1996 Current Opinion in Neurology

Comparing Sequences

 Define distance between two sequences as the number of mutations that would result in the second string, starting from the first

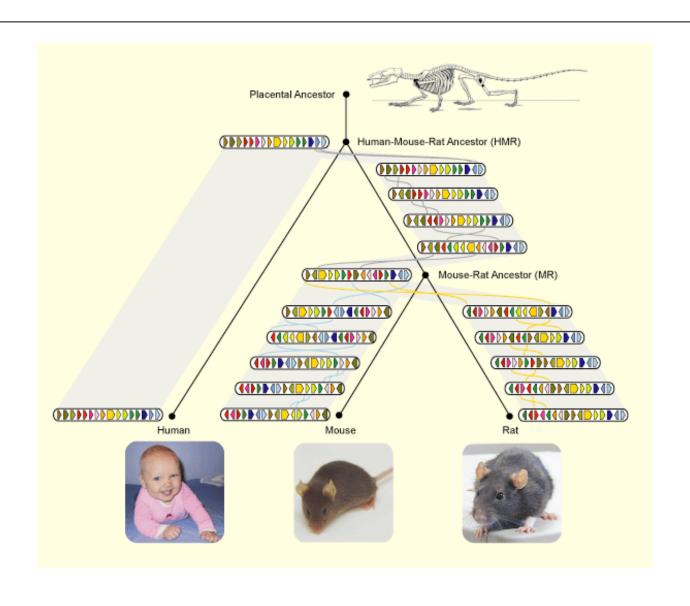
ACGGCGTGCTTTAGAACATAG

AAGGCGTGCTTTAGAACATAG

AAGGCGTGC**G**TTAGAACATAG

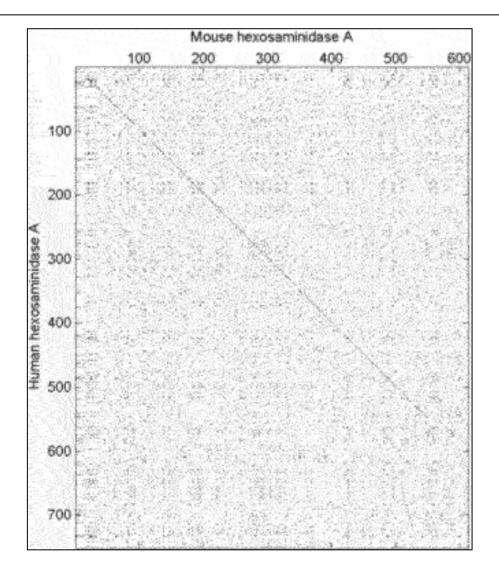
ACGGCGTGCGT<u>A</u>AG<u>G</u>ACA<u>A</u>TAG

Evolution and Edit Distances



Plotting Genome Rearrangements

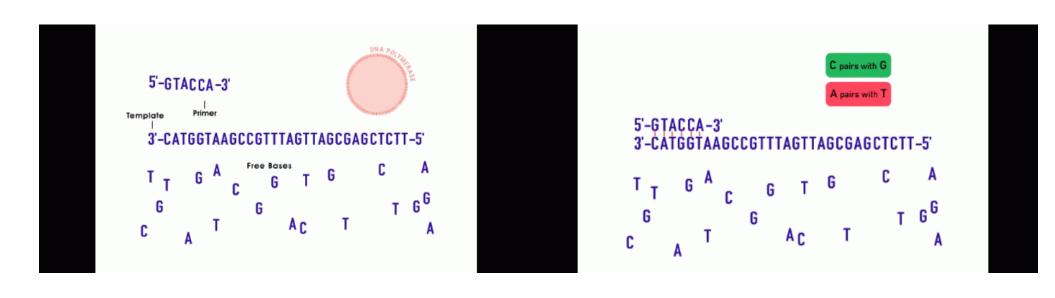
Diagonals imply direct alignment Reverse diagonals imply inverse alignment

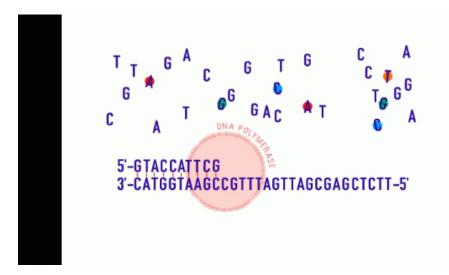


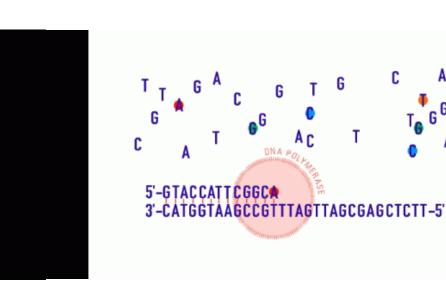
Genomic Sequences

- Sanger Sequencing
- Next-Generation Sequencing
 - Illumina Solexa
 - Helicos
 - Solid
 - Roche/454

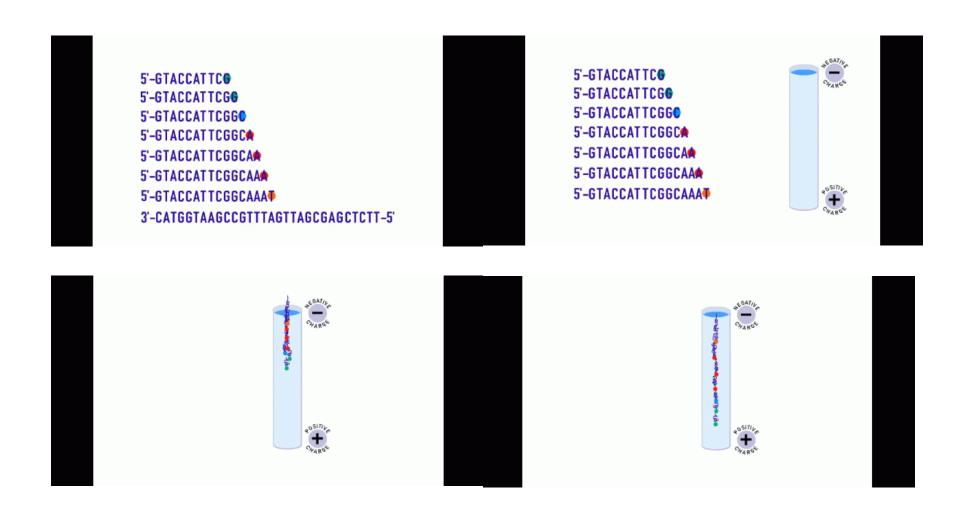
Sanger Sequencing



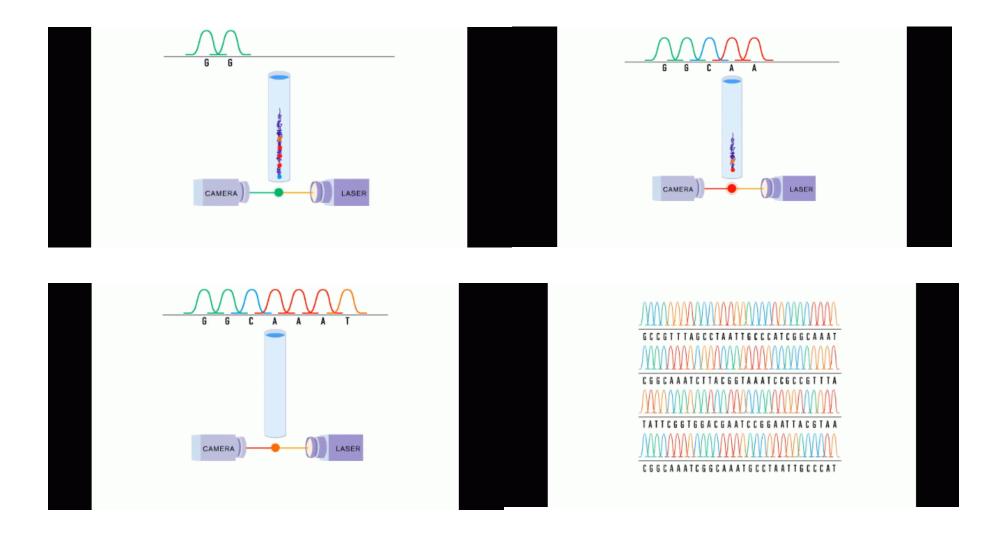




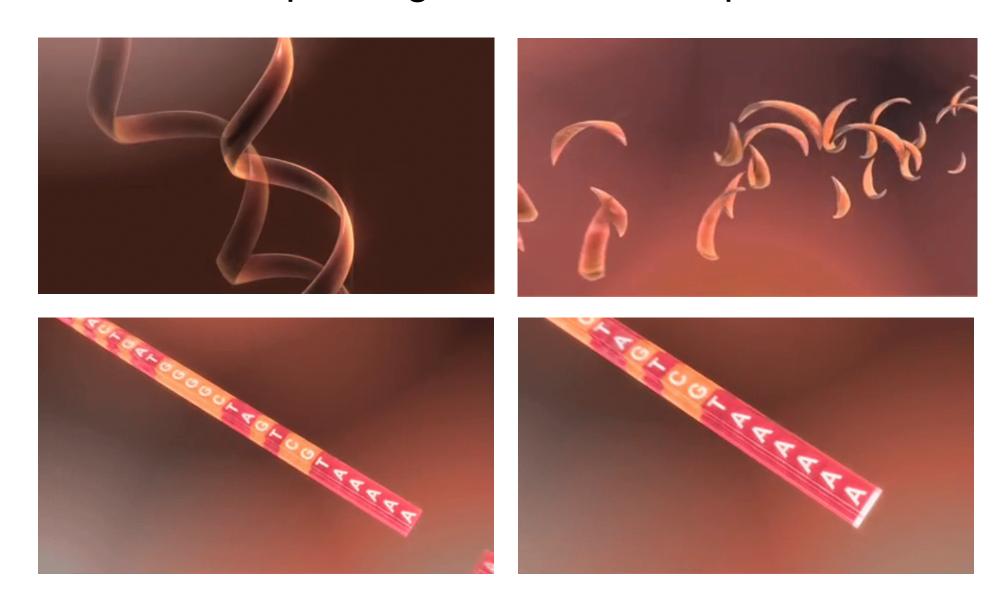
Sanger Sequencing



Sanger Sequencing

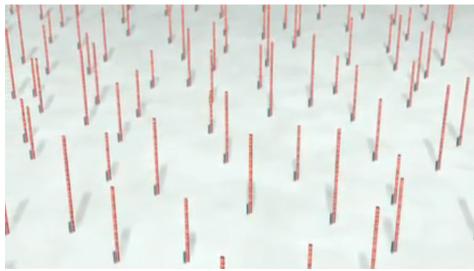


HeliScope Single Molecule Sequencer

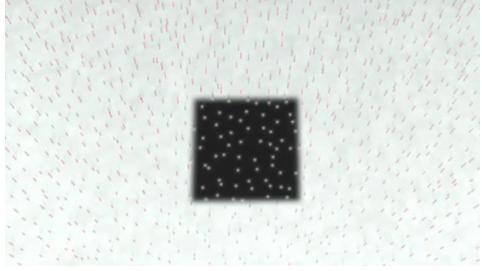


HeliScope Single Molecule Sequencer





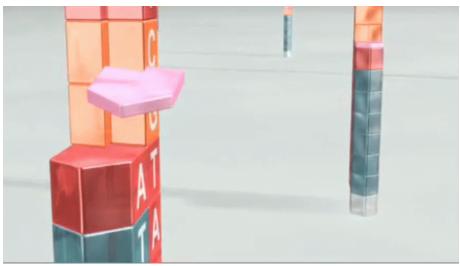




HeliScope Single Molecule Sequencer

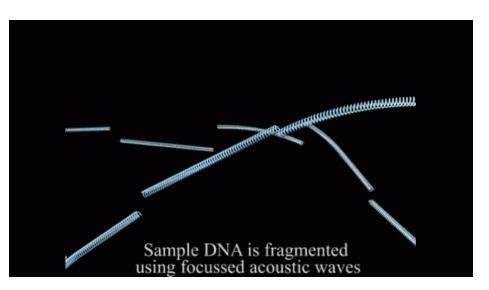


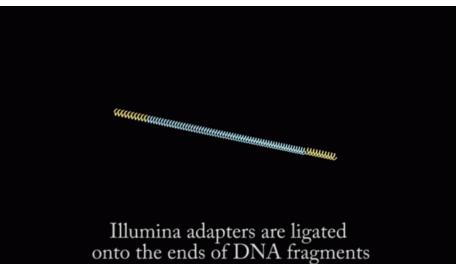


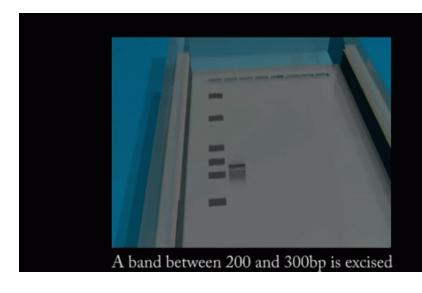


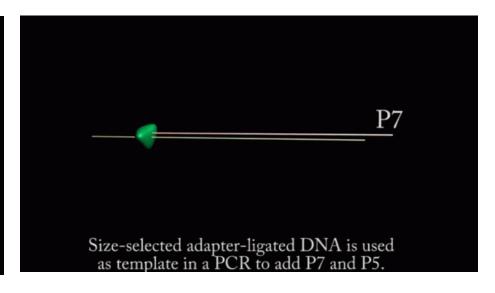


Illumina Solexa

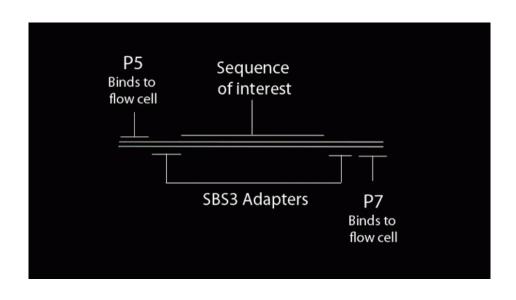


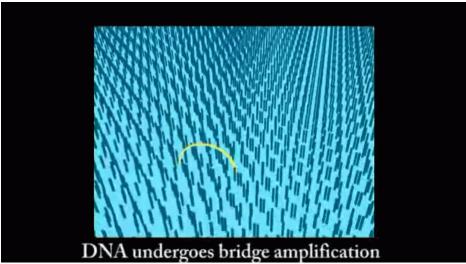


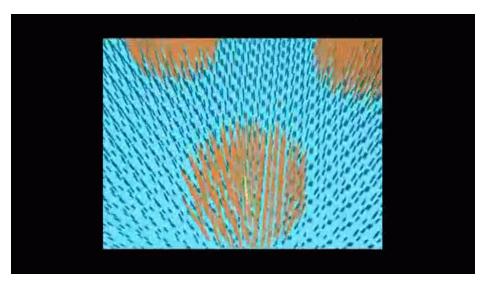


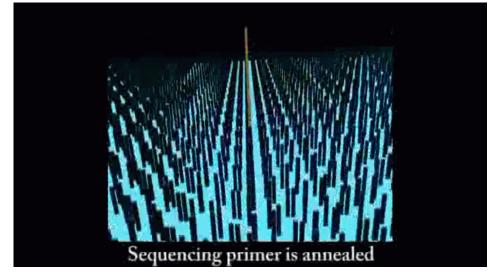


Illumina Solexa

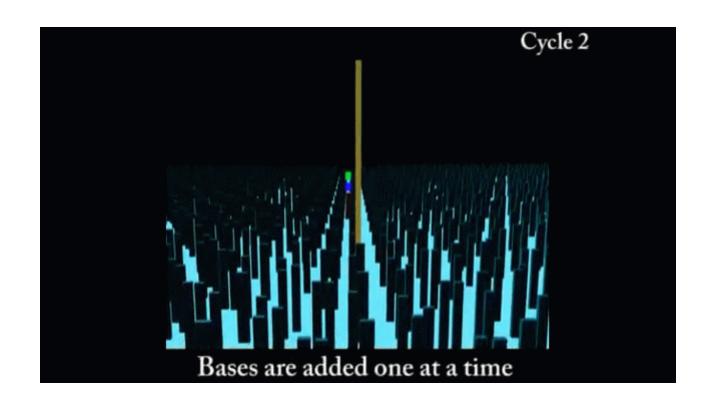








Illumina Solexa



Dealing with Reads

- From fluorescence to nucleotides (Phread)
- Error correction
- Mapping to reference genomes
- Assembly

Error Correction

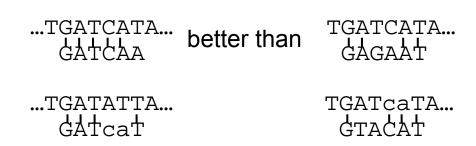
- NGS reads range from 50 300 bps (constantly changing)
- Error rates range from 1 3%
- Errors are not uniformly distributed over the read
- Correcting errors is a critical step before mapping/ assembly

Error Correction

- Needed coverage on the genome
- k-mer based error correction
- Suffix Trees

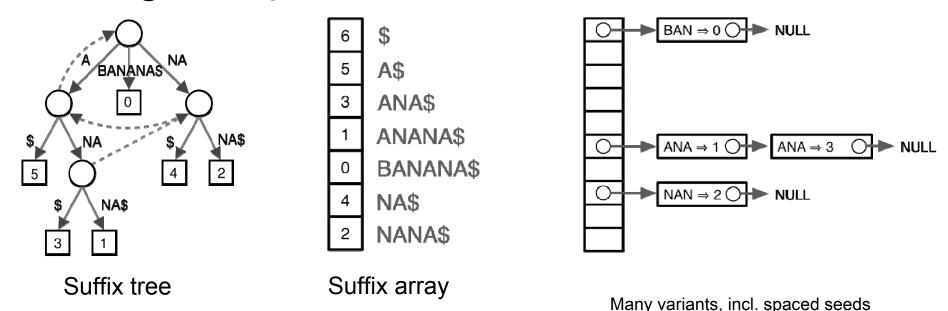
Short Read Alignment

- Given a reference and a set of reads, report at least one "good" local alignment for each read if one exists
 - Approximate answer to: where in genome did read originate?
- What is "good"? For now, we concentrate on:
 - Fewer mismatches is better
 - Failing to align a low-quality base is better than failing to align a high-quality base



Indexing

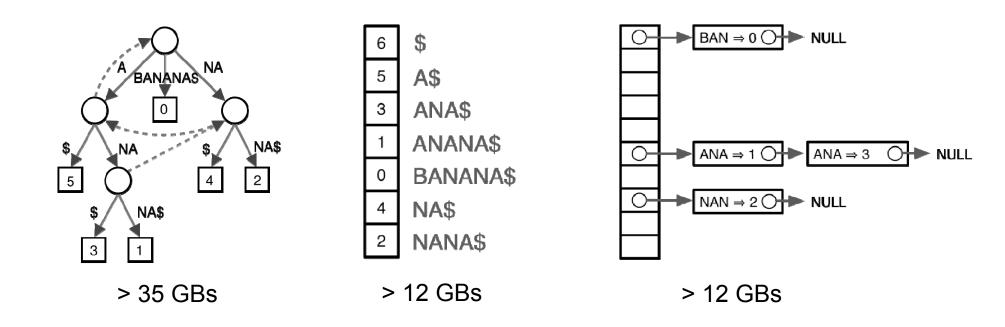
- Genomes and reads are too large for direct approaches like dynamic programming
- Indexing is required



Choice of index is key to performance

Indexing

Genome indices can be big. For human:

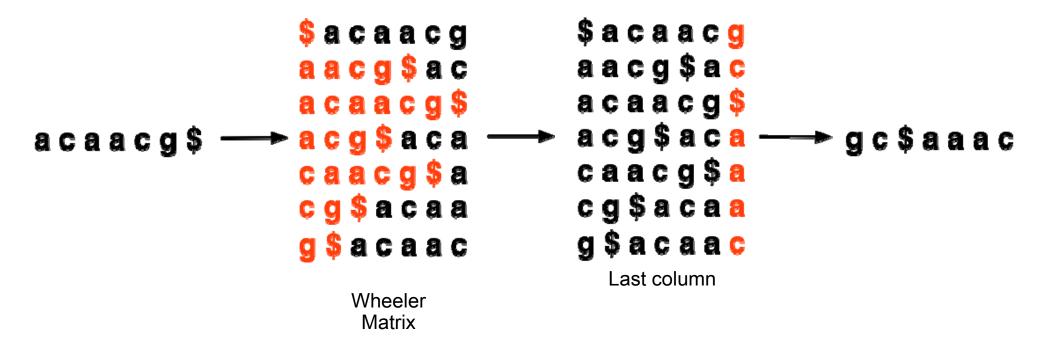


- Large indices necessitate painful compromises
 - 1. Require big-memory machine
 - 2. Use secondary storage

- 3. Build new index each run
- Subindex and do multiple passes

Burrows-Wheeler Transform

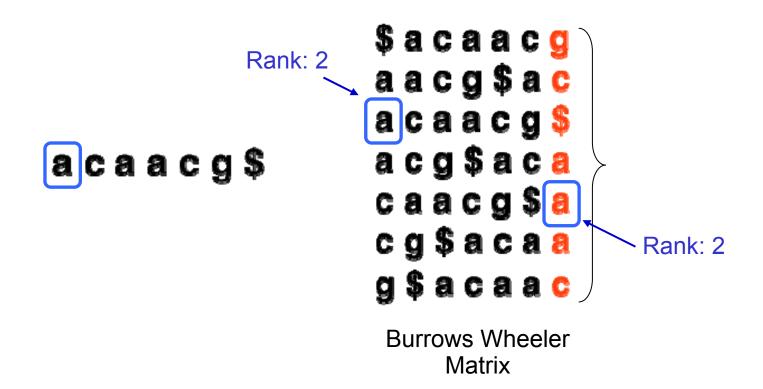
Reversible permutation used originally in compression



- Once BWT(T) is built, all else shown here is discarded
 - Matrix will be shown for illustration only

Burrows-Wheeler Transform

- Property that makes BWT(T) reversible is "LF Mapping"
 - ith occurrence of a character in Last column is same text occurrence as the ith occurrence in First column



Burrows-Wheeler Transform

- To recreate T from BWT(T), repeatedly apply rule:
 - T = BWT[LF(i)] + T; i = LF(i)
 - Where LF(i) maps row i to row whose first character corresponds to i's last per LF Mapping



Could be called "unpermute" or "walk-left" algorithm

FM Index

Ferragina & Manzini propose "FM Index" based on BWT

- Observed:
 - LF Mapping also allows exact matching within T
 - LF(i) can be made fast with checkpointing
 - ...and more (see FOCS paper)

Ferragina P, Manzini G: Opportunistic data structures with applications. FOCS. IEEE Computer Society; 2000.

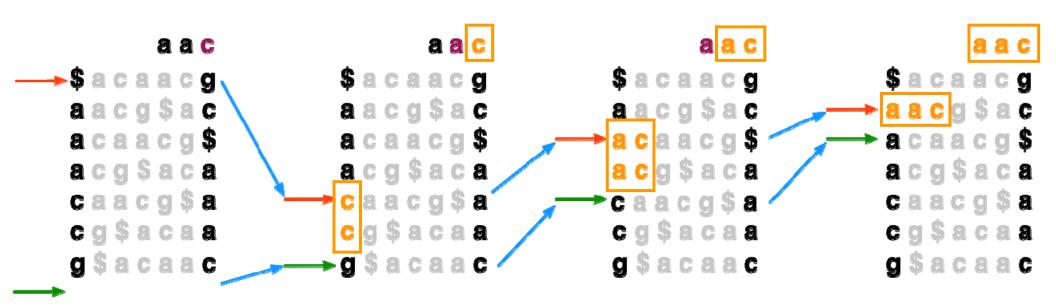
Ferragina P, Manzini G: An experimental study of an opportunistic index. *SIAM symposium on Discrete algorithms*. Washington, D.C.; 2001.

Exact Matching with FM Index

- To match Q in T using BWT(T), repeatedly apply rule:
 - top = LF(top, qc); bot = LF(bot, qc)
 - Where qc is the next character in Q (right-to-left) and LF(i, qc) maps row i
 to the row whose first character corresponds to i's last character as if it
 were qc

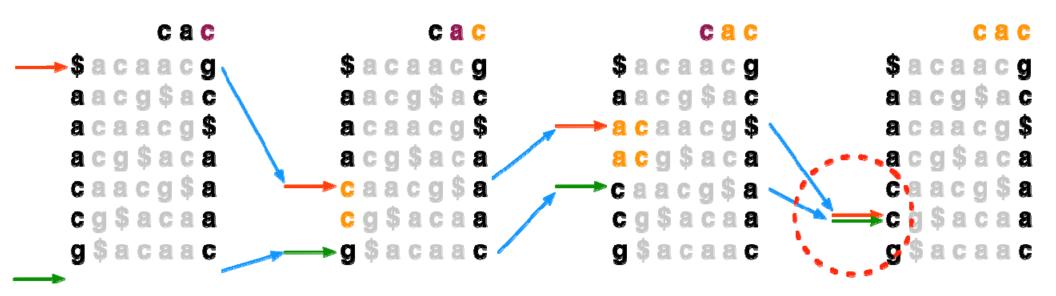


Exact Matching with FM Index



 In progressive rounds, top & bot delimit the range of rows beginning with progressively longer suffixes of Q

Exact Matching with FM Index



 If range becomes empty (top = bot) the query suffix (and therefore the query) does not occur in the text

Backtracking

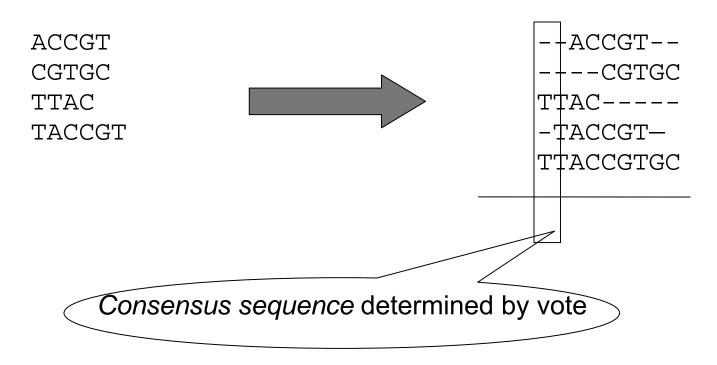
Consider an attempt to find Q = "agc" in T = "acaacg":

```
$acaacg
aacg$ac
acaacg$
acaacg$
acaacg$
acg$aca
acg$aca
caacg$a
caacg$a
caacg$a
caacg$a
cg$acaa
```

 Instead of giving up, try to "backtrack" to a previous position and try a different base

Sequencing

Find maximal overlaps between fragments:



Quality Metrics

The *coverage* at position *i* of the <u>target</u> or <u>consensus</u> sequence is the number of fragments that overlap that position

Target:	
	No coverage

Two contigs

Quality Metrics

Linkage – the degree of overlap between fragments

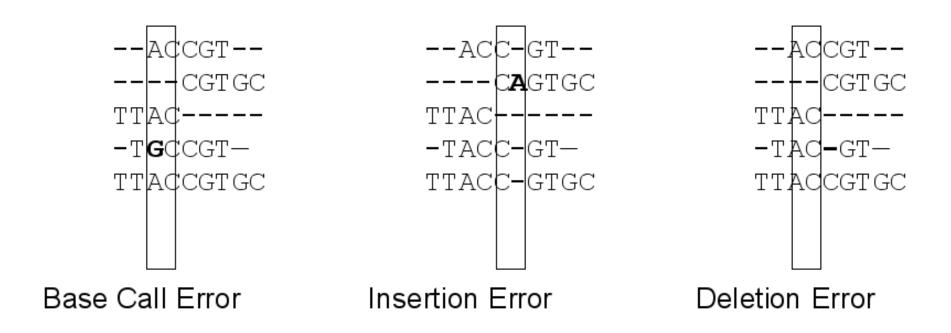
Target:				
		_		
_		_		
_	_			_

Perfect coverage, poor average linkage poor minimum linkage

Real World Complications

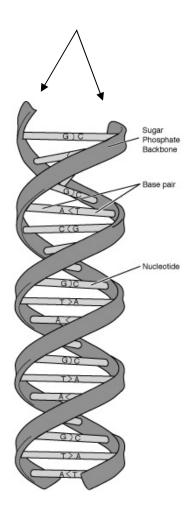
Base call errors

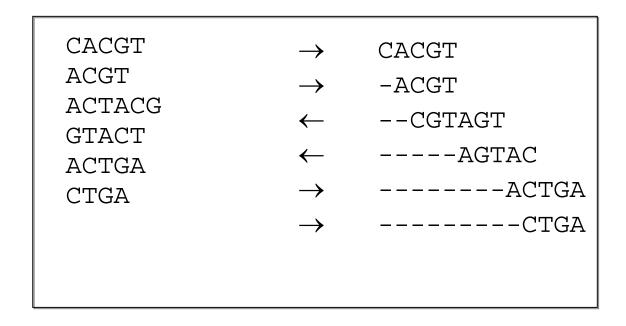
Chimeric fragments, contamination (e.g. from the vector)



Unknown Orientation

A fragment can come from either strand





Sequence Alignment Models

Shortest common superstring

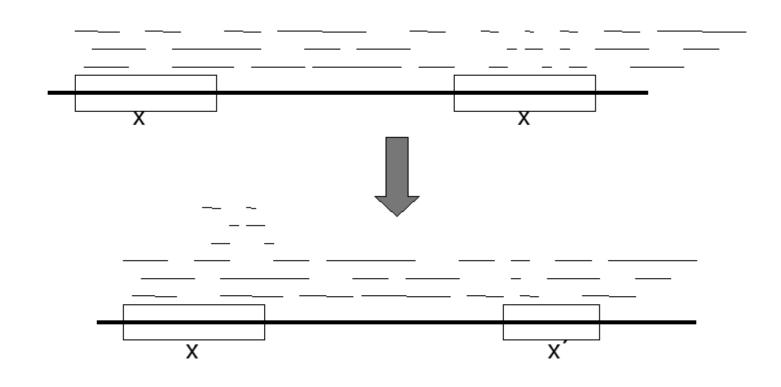
Input: A collection, F, of strings (fragments)

Output: A shortest possible string S such that for every $f \in \mathcal{F}$, S is a superstring of f.

Example:

```
\mathcal{F} = \{ACT, CTA, AGT\}
S = ACTAGT
```

Problems with the SCS model



- Directionality of fragments must be known
- No consideration of coverage
- Some simple consideration of linkage
- No consideration of base call errors

Reconstruction

Deals with errors and unknown orientation

Definitions

f is an approximate substring of S at error level ε when $d_{\varepsilon}(f, S) \leq \varepsilon \times 1$

 $\mid f \mid$

 d_s = substring edit distance:

Reconstruction

Input: A collection, \mathcal{F} , of strings, and a tolerance level, ε

Output: Shortest possible string, S, such that for every $f \in \mathcal{F}$:

sible string,
$$S$$
, such that for every $f \in \mathcal{F}$:

 $\min \ d_s \mathcal{J}, S \mathcal{J}d_s \mathcal{J}, S \mathcal{S} \mathcal{I} \mathcal{I}$

Reconstruction Example

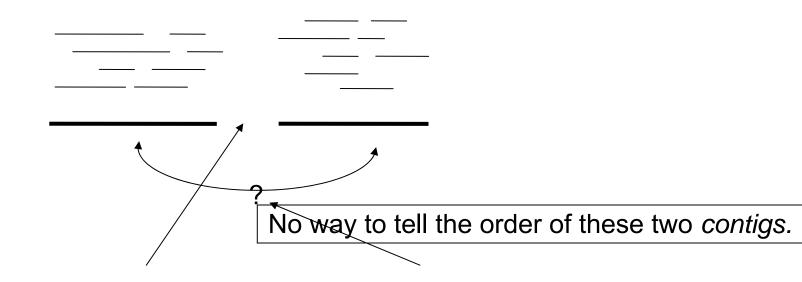
Input: $\mathcal{F} = \{ ATCAT, GTCG, CGAG, TACCA \}$ ε = 0.25 ATCAT **Output:** ATGAT -----CGAC -CGAG **GTCG** ---TAC**C**A ACGATACGAC $d_s(CGAG, ACGATACGAC) = 1$ $= 0.25 \times 4$ So this output is OK for $\mathcal{E} = 0.25$

Limitations of Reconstruction

- Models errors and unknown orientation
- Doesn't handle repeats
- Doesn't model coverage
- Only handles linkage in a very simple way
- Always produces a single contig

Contigs

Sometimes you just can't put all of the fragments together into one *contig*uous sequence:



No way to tell how much sequence is missing between them.

Multicontig

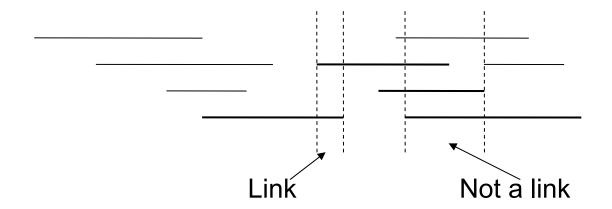
Definitions

A layout, \mathcal{L} , is a multiple alignment of the fragments

Columns numbered from 1 to $|\mathcal{L}|$

Endpoints of a fragment: I(f) and r(f)

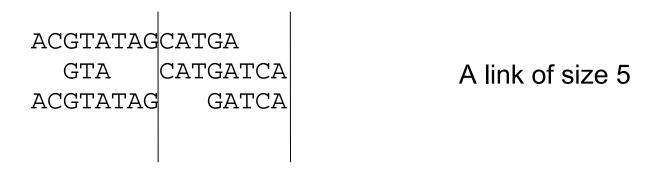
An overlap is a *link* where no other fragment completely covers the overlap



Multicontig

More definitions

The size of a link is the number of overlapping positions



The weakest link is the smallest link in the layout

A t-contig has a weakest link of size t

A collection, \mathcal{F} , admits a t-contig if a t-contig can be constructed from the fragments in \mathcal{F}

Perfect Multicontig

Input: \mathcal{F} , and t

Output: a minimum number of collections, C_i , such that every C_i admits a t-contig

Let \mathcal{F} = {GTAC, TAATG, TGTAA}

t = 3

--TAATG
TGTAA--

GTAC

t = 1

TGTAA-----TAATG------GTAC

Handling errors in Multicontig

- •The *image* of a fragment is the portion of the consensus sequence, *S*, corresponding to the fragment in the layout
- •S is an ε -consensus for a collection of fragments when the edit distance from each fragment, f, and its image is at most $\varepsilon \times |f|$

TATAGCAT**C**AT

CGT**C** CATGATCA

ACG**G**ATAG G**TC**CA

ACGTATAGCATGATCA

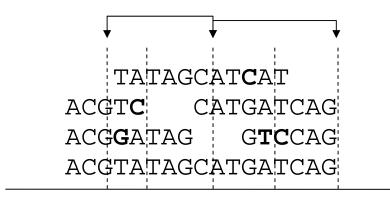
An ε -consensus for ε = 0.4

Definition of Multicontig

Input: A collection, \mathcal{F} , of strings, an integer $t \geq 0$, and an error tolerance ε between 0 and 1 Output: A partition of \mathcal{F} into the minimum number of collections C_i such that every C_i admits a t-contig with an ε -consensus

Example of Multicontig

Let ε = 0.4, t = 3



Assembly Algorithms

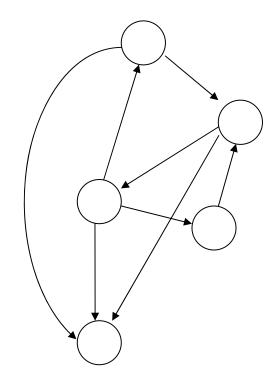
Most of the algorithms to solve the fragment assembly problem are based on a graph model

A graph, G, is a collection of edges, e, and vertices, v.

Directed or undirected

Weighted or unweighted

We will discuss representations and other issues shortly...

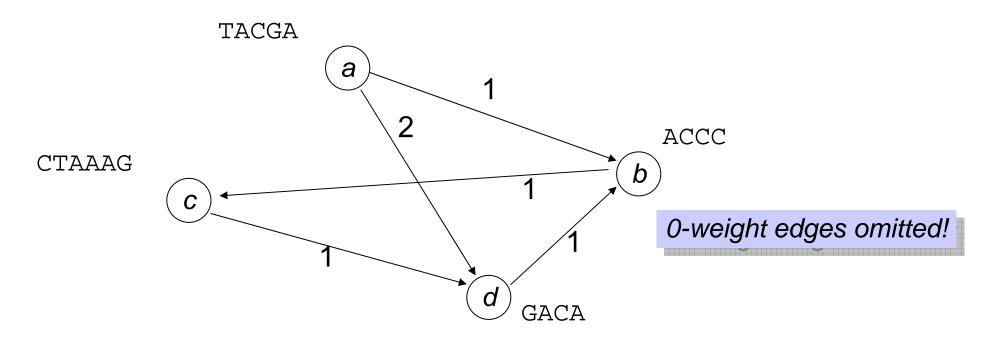


A directed, unweighted graph

The Maximum Overlap Graph

Overlap multigraph

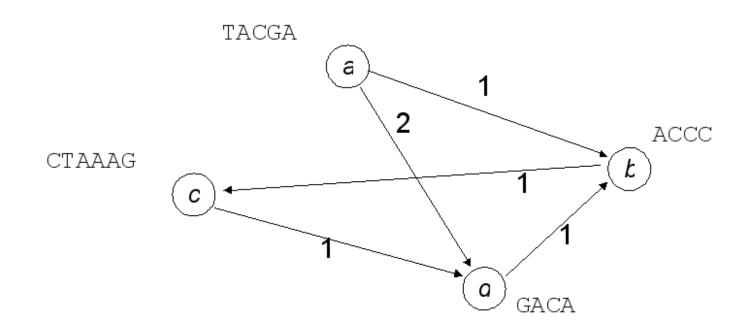
Each directed edge, (u,v) is weighted with the length of the maximal overlap between a suffix of u and a prefix of v



Paths and Layouts

The path *dbc* leads to the alignment:

```
GACA-----
---ACCC-----
-----CTAAAG
```



Superstrings

Every path that covers every node is a superstring Zero weight edges result in alignments like:

Higher weights produce more overlap, and thus shorter strings

The *shortest common superstring* is the highest weight path that covers every node

Graph formulation of SCS

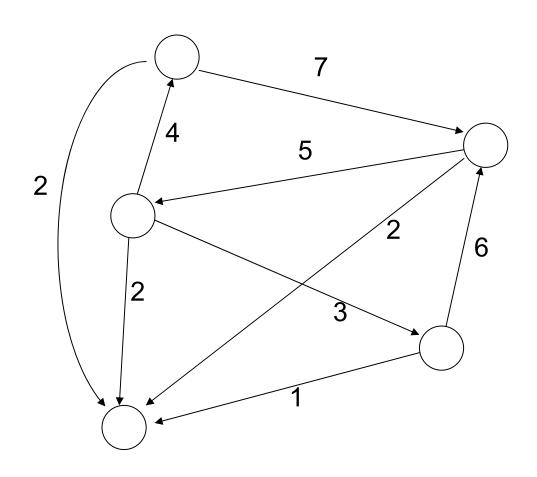
Input: A weighted, directed graph

Output: The highest-weight path that touches

every node of the graph

NP Hard, Use Greedy Approximation

Greedy Example



So we have sequences now!

- Find genes in sequences.
- Query: AGTACGTATCGTATAGCGTAA

What does it do?

- Find similar gene in other species with known function and reason from it
- Align sequences with known genes
- Find the gene with the "best" match

Sequence Alignment

 Point mutations can be easily handled: ACGTCTGATACGCCGTATAGTCTATCT
 ACGTCTGATTCGCCCCTATCGTCTATCT

Insertions and deletions (InDels) are harder!

ACGTCTGATACGCCGTATAGTCTATCT CTGATTCGCATCGTCTATCT

ACGTCTGAT**A**CGCCGTAT**A**GTCTATCT ----CTGAT**T**CGC---AT**C**GTCTATCT

Sequence Alignment: Scoring

```
Match score: +1
```

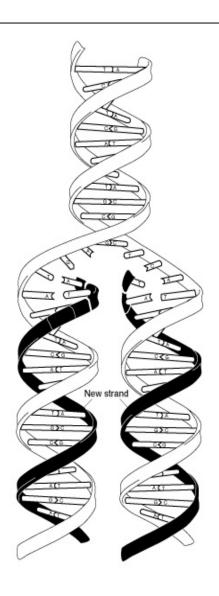
```
ACGTCTGATACGCCGTATAGTCTATCT
```

Matches:
$$18 \times (+1)$$

Gaps:
$$7 \times (-1)$$

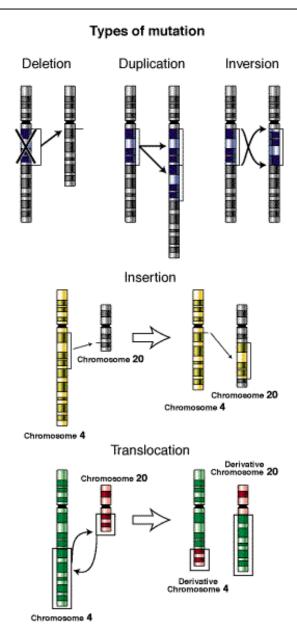
Sequence Alignment: Scoring

- Prior to cell division, all the genetic instructions must be "copied" so that each new cell will have a complete set
- DNA polymerase is the enzyme that copies DNA
 - Synthesizes in the 5' to 3' direction



Sequence Alignment: Scoring

- Environmental factors
- Mistakes in replication or repair



Deletions in Sequences

Codon deletion:

ACG ATA GCG TAT GTA TAG CCG...

- Effect depends on the protein, position, etc.
- Almost always deleterious
- Sometimes lethal
- Frame shift mutation:

ACG ATA GCG TAT GTA TAG CCG...

ACG ATA GCG ATG TAT AGC CG?...

Almost always lethal

Insertions/ Deletions

 It is very difficult to determine whether an *InDel* is an insertion in one gene, or a deletion in another, unless ancestry is known:

ACGTCTGATACGCCGTATCGTCTATCT ACGTCTGAT---CCGTATCGTCTATCT

Insertions/ Deletions

- We want to find alignments that are evolutionarily likely.
- Which of the following alignments is more likely?

```
ACGTCTGATACGCCGTATAGTCTATCT
ACGTCTGAT-----ATAGTCTATCT
ACGTCTGATACGCCGTATAGTCTATCT
AC-T-TGA--CG-CGT-TA-TCTATCT
```

 Initiating a gap must cost more than extending an existing gap! (why?)

Alignments

- Match/mismatch score: +1/+0
- Origination/length penalty: −2/−1

ACGTCTGATACGCCGTATAGTCTATCT

- ----CTGAT**T**CGC---AT**C**GTCTATCT
- Matches: 18 × (+1)
- Mismatches: 2 × 0
- Origination: $2 \times (-2)$
- Length: $7 \times (-1)$

Optimal Alignments

- Finding optimal alignment hard:
 ACGTCTGATACGCCGTATAGTCTATCT
 CTGAT---TCG-CATCGTC--T-ATCT
- C(27,7) gap positions = ~888,000 possibilities
- Dynamic programming: The Smith Waterman algorithm

Optimal Alignments

An Example:

ACTCG ACAGTAG

- Match: +1
- Mismatch: 0
- Gap: −1

Dynamic Programming

- Each sequence along one axis
- Mismatch penalty multiples in first row/column
 0 in [0,0]

Dynamic Programming

- Vertical/Horiz. move: Score + (simple) gap penalty
- Diagonal move: Score + match/mismatch score
- Take the MAX of the three possibilities

		A	C	T	C	G
	0	-1	-2	-3	-4	-5
A	-1	1				
C	-2					
A	-3					
G	-4					
T	-5					
A	-6					
G	-7					

Dynamic Programming

		a	С	t	С	g
	0	-1	-2	-3	-4	-5
а	-1	1	0	-1	-2	-3
С	-2	0	2	1	0	-1
а	-3	-1	1	2	1	0
g	-4	-2	0	1	2	2
t	-5	-3	-1	1	1	2
а	-6	-4	-2		1	1
g	-7	-5	-3	-1	0	2

Optimal Alignment

Trace back from the maximum value to the origin.

		a	С	t	С	g
	0	-1	-2	-3	-4	-5
а	-1	1	0	-1	-2	-3
С	-2	0	2	1	0	-1
a	-3	-1	1	2	1	0
g	-4	-2	0	1	2	2
t	-5	-3	-1	1	1	2
а	-6	-4	-2	0	1	1
g	-7	-5	-3	-1	O [*]	2

Paths Correspond to Alignments

- = GAP in left sequence
- = ALIGN both positions
- Corresponding alignment (start at the end):

Score
$$= +2$$

Semi-Global Alignments

Suppose we are aligning:

GCG GGCG

Which one is biologically relevant?

G-CG -GCG

GGCG GGCG

 Semi-global alignment allows gaps at the ends for free.

Semi-global alignment

 Semi-global alignment allows gaps at the ends for free.

		g	С	g
	, 0	0	0	0
g	Q	1	0	1
g	0	1	1	1
С	0	0	2	1
g	0	1	1	3

Local alignment

- Global alignment score entire alignment
- Semi-global alignments allow unscored gaps at the beginning or end of either sequence
- Local alignment find the best matching subsequence
- CG**ATG**AA**ATG**GA
- This is achieved through a 4th alternative at each position in the table: zero.

Local alignment

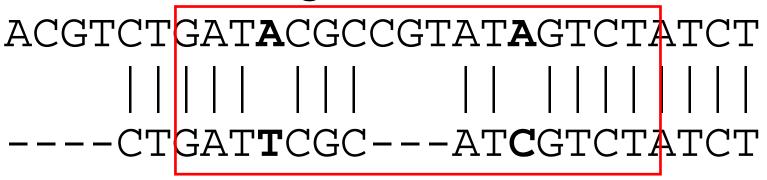
• Mismatch = -1 this time

		С	g	a	t	g
	0	-1	-2	-3	-4	-5
а	-1	0	0	0	0	0
а	-2	0	0	. 1	0	0
а	-3	0	0	1	0	0
t	-4	0	0	0	2	1
g	-5	0	1	0	1	3
g	-6	0	1	0	0	2
а	-7	0	0	2	1	1

CG**ATG** AA**ATG**GA

Optimal Sub-alignments

Consider the alignment:



 Is it true that the alignment in the boxed region must be optimal?

A Greedy Strategy

Consider this pair of sequences

GAGC

CAGC

GAP = -1

Match = +1

Mismatch = -2

Greedy Approach:

G

or

G

or

G

Leads to

GAGC---

Better:

GACG

---CAGC

CACG

Divide and Conquer

Suppose we are aligning:

ACTCG ACAGTAG

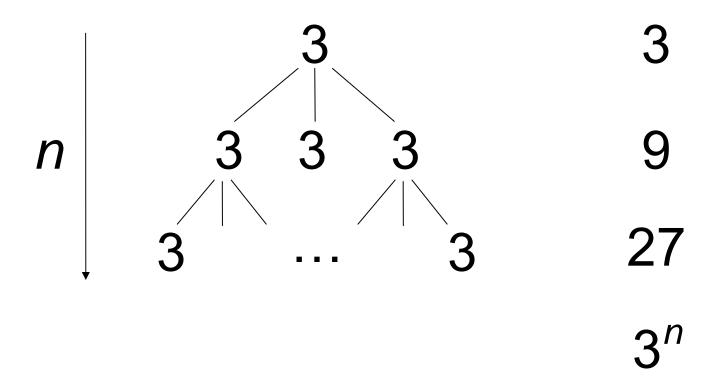
First position choices:

A A	+1	CTCG CAGTAG
A -	-1	CTCG ACAGTAG
– А	-1	ACTCG CAGTAG

Complexity of RecurseAlign

 What is the recurrence equation for the time needed by RecurseAlign?

$$T \square n \equiv 3T \square n - 1 \square \square 3$$



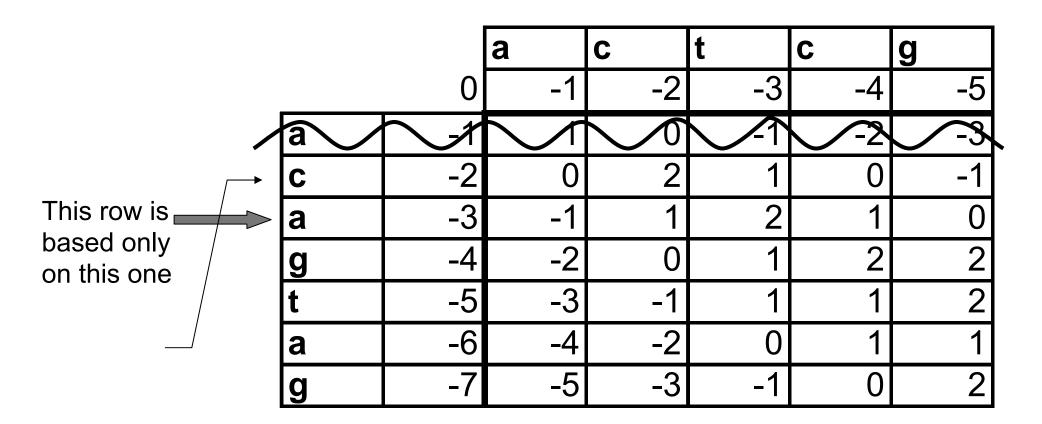
Dynamic Programming

		a	С	t	С	g
	0	-1	-2	-3	-4	-5
a	-1	1	0	-1	-2	-3
С	-2	0	2	1	0	-1
a	-3	-1	1	2	1	0
g	-4	-2	0	1	2	2
t	-5	-3	-1	1	1	2
a	-6	-4	-2	0	1	1
g	-7	-5	-3	-1	0	2

 This is possible for any problem that exhibits optimal substructure (Bellman's principle of optimality)

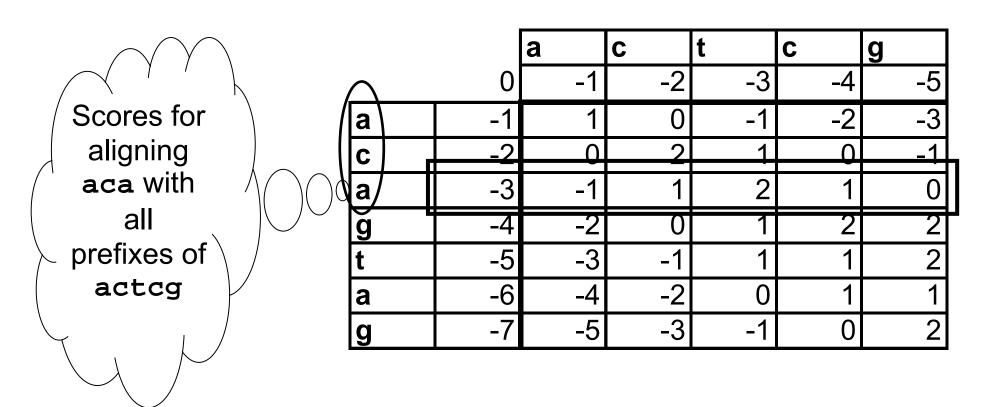
Space Complexity

 Note that we can throw away the previous rows of the table as we fill it in:



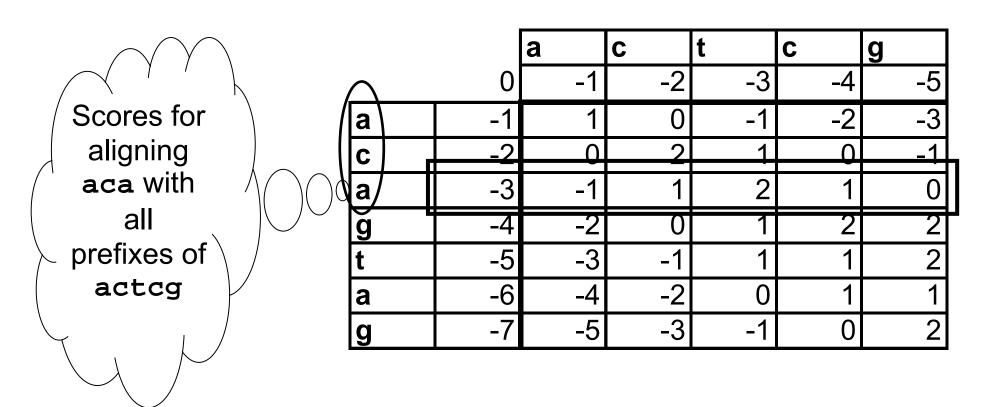
Space Complexity

• **Each row** of the table contains the scores for aligning a **prefix** of the left-hand sequence with **all prefixes** of the top sequence:



Space Complexity

• **Each row** of the table contains the scores for aligning a **prefix** of the left-hand sequence with **all prefixes** of the top sequence:



So Where Does i Line Up?

- Find out where i aligns to the bottom sequence
 - Needs two vectors of scores

```
i
↓
s: ACGCTATGCTCATAG

t: CGACGCTCATCG
```

- Assuming i lines up with a character:
 alignscore = align(ACGCTAT, prefix(t)) + score(G, char from t)
 + align(CTCATAG, suffix(t))
- Which character is best?
 - Can quickly find out the score for aligning ACGCTAT with every prefix of t.

So where does *i* line up?

But, i may also line up with a gap

↓ S: ACGCTAT**G**CTCATAG

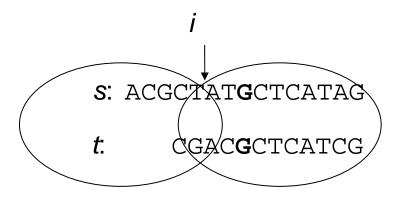
t. CGACGCTCATCG

Assuming i lines up with a gap:

alignscore = align(ACGCTAT, prefix(t)) + gapscore + align(CTCATAG, suffix(t))

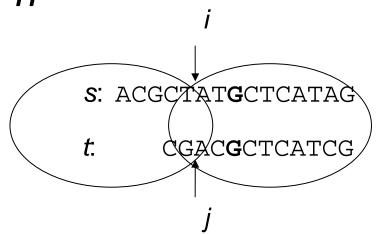
Recursive Call

- Fix the best position for I
- Call align recursively for the prefixes and suffixes:



Time Complexity

- Let len(s) = m and len(t) = n
- Space: 2m
- Time:
 - Each call to build similarity vector = m'n'
 - First call + recursive call:



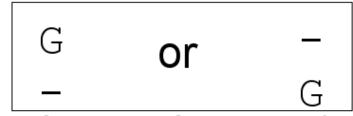
$$T m, n \leq \frac{mn}{2} - \frac{mn}{2} + T \frac{m}{2}, j + T \frac{m}{2}, n - j$$

$$mn + mj + m - j$$

$$2 mn$$

General Gap Penalties

- Suppose we are no longer using simple gap penalties:
 - Origination = -2
 - Length = -1
- Consider the last position of the alignment for ACGTA with ACG
- We can't determine the score for



unless we know the previous positions!

Scoring Blocks

Now we must score a block at a time

- A block is a pair of characters, or a maximal group of gaps paired with characters
- To score a position, we need to either start a new block or add it to a previous block

Alignment Algorithm

- Three tables
 - a scores for alignments ending in char-char blocks
 - b scores for alignments ending in gaps in the top sequence
 (s)
 - c scores for alignments ending in gaps in the left sequence
 (t)
- Scores no longer depend on only three positions, because we can put any number of gaps into the last block

The Recurrences

The Optimal Alignment

- The optimal alignment is found by looking at the maximal value in the lower right of all three arrays
- The algorithm runs in $O(n^3)$ time
 - Uses $O(n^2)$ space

Searching in Sequence Databases: BLAST

Database Searching

- How can we find a particular short sequence in a database of sequences (or one HUGE sequence)?
- Problem is identical to local sequence alignment, but on a much larger scale.
- We must also have some idea of the significance of a database hit.
 - Databases <u>always</u> return some kind of hit, how much attention should be paid to the result?

BLAST

- BLAST: <u>Basic Local Alignment Search Tool</u>
- An approximation of the Dynamic Programming algorithm
- Sacrifices some search sensitivity for speed

Scoring Matrices

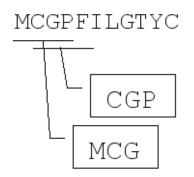
- DNA
 - Identity
 - Transition/Transversion

- Proteins
 - PAM
 - BLOSUM

	Α	R	N	D	С	Q	Е	G	Н	I	L	K	M	F	Р	S	Т	W	Υ	V
Α	2																			
R	-2	6																		
N	0	0	2																	
D	0	-1	2	4																
С	-2	-4	-4	-5	4															
Q	0	1	1	2	-5	4														
Е	0	-1	1	3	-5	2	4													
G	1	-3	0	1	-3	-1	0	5												
Н	-1	2	2	1	-3	3	1	-2	6											
ı	-1	-2	-2	-2	-2	-2	-2	-3	-2	5										
L	-2	-3	-3	-4	-6	-2	-3	-4	-2	2	6									
K	-1	3	1	0	-5	1	0	-2	0	-2	-3	5								
M	-1	0	-2	-3	-5	-1	-2	-3	-2	2	4	0	6							
F	-4	-4	-4	-6	-4	-5	-5	-5	-2	1	2	-5	0	9						
Р	1	0	-1	-1	-3	0	-1	-1	0	-2	-3	-1	-2	-5	6					
S	1	0	1	0	0	-1	0	1	-1	-1	-3	0	-2	-3	1	3				
Т	1	-1	0	0	-2	-1	0	0	-1	0	-2	0	-1	-2	0	1	3			
W	-6	2	-4	-7	-8	-5	-7	-7	-3	-5	-2	-3	-4	0	-6	-2	-5	17		
Υ	-3	-4	-2	-4	0	-4	-4	-5	0	-1	-1	-4	-2	7	-5	-3	-3	0	10	
V	0	-2	-2	-2	-2	-2	-2	-1	-2	4	2	-2	2	-1	-1	-1	0	-6	2	4

The BLAST algorithm

- Break the search sequence into words
 - -W=3 for proteins, W=12 for DNA



MCG, CGP, GPF, PFI, FIL, ILG, LGT, GTY, TYC

 Include in the search all words that score above a certain value (T) for any search

word

MCG CGP
MCT MGP ...
MCN CTP

This list can be computed in linear time

... ..

BLAST Algorithm

- Search for the words in the database
 - Word locations can be precomputed and indexed
 - Searching for a short string in a long string
 - Regular expression matching: FSA
- HSP (High Scoring Pair) = A match between a query word and the database
- Find a "hit": Two non-overlapping HSP's on a diagonal within distance A
- Extend the hit until the score falls below a threshold value, X

The BLAST Search Algorithm

query word (W = 3) GSVEDTTGSQSLAALLNKCKTPQGQRLVNQWIKQPLMDKNRIEERLNLVEAFVEDAELRQTLQEDL Query: PQGPEG 15 PRG 14 neighborhood PKG 14 13 PNG words PDG 13 PHG neighborhood PMG PSG score threshold POA (T = 13)PON etc....

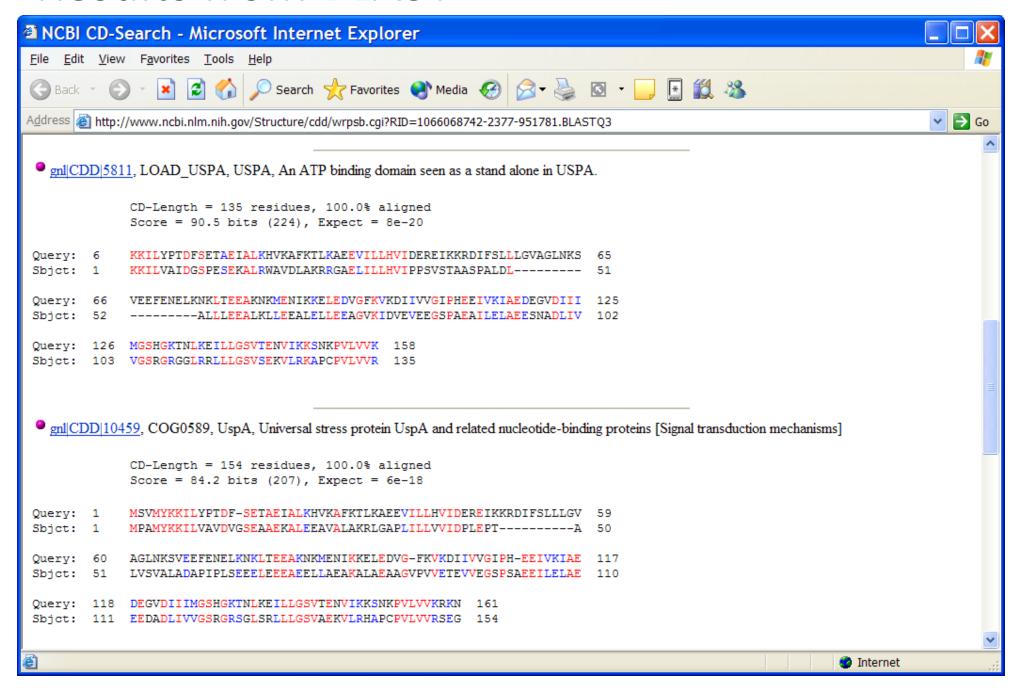
Query: 325 SLAALLNKCKTPQGQRLVNQWIKQPLMDKNRIEERLNLVEA 365

+LA++L+ TP G R++ +W+ P+ D + ER + A

Sbjct: 290 TLASVLDCTVTPMGSRMLKRWLHMPVRDTRVLLERQQTIGA 330

High-scoring Segment Pair (HSP)

Results from BLAST



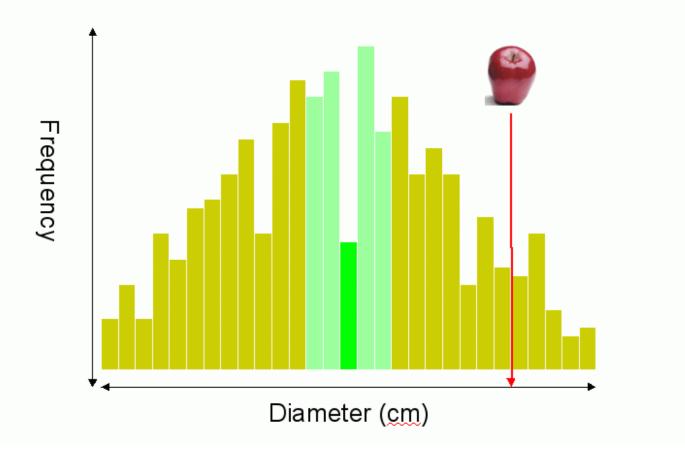
Search Significance Scores

A search will always return some hits.

- How can we determine how "unusual" a particular alignment score is?
 - ORF's
 - Assumptions

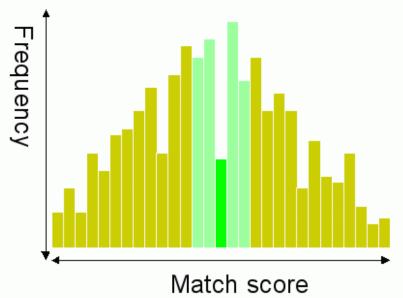
Significance from a Distribution

 I have an apple of diameter 5". Is that unusual?



Is a Match Significant?

- Match scores for aligning my sequence with random sequences.
- Depends on:
 - Scoring system
 - Database
 - Sequence to search for
 - Length
 - Composition



 How do we determine the random sequences?

Generating "random" sequences

Random uniform model:

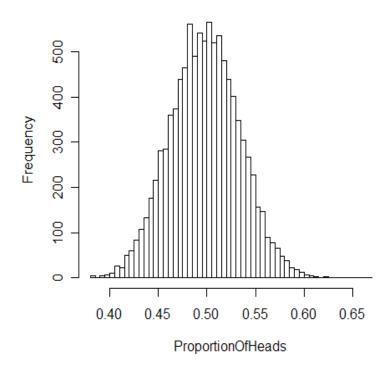
$$P(G) = P(A) = P(C) = P(T) = 0.25$$

- Doesn't reflect nature
- Use sequences from a database
 - Might have genuine homology
 - We want unrelated sequences
- Random shuffling of sequences
 - Preserves composition
 - Removes true homology

Sums of Distributions

- The mean of n random (i.i.d.) events tends towards a Normal/ Gaussian.
 - Example: Throw n dice and compute the mean.
 - Distribution of means:

Histogram of ProportionOfHeads



The Extreme Value Distribution

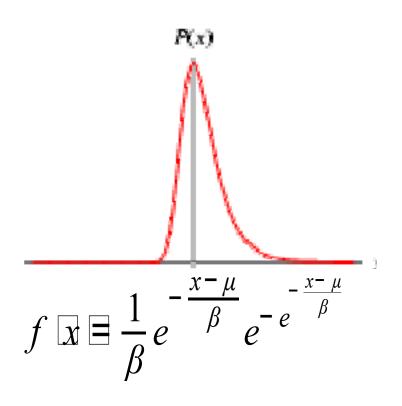
 This means that if we get the match scores for our sequence with n other sequences, the mean would follow a Gaussian distribution.

 The maximum of n (i.i.d.) random events tends towards the extreme value distribution as n grows large.

Gaussian/ Extreme Value Distributions

Extreme Value:

Gaussian:



$$f = \frac{1}{\sigma} \frac{e^{-|x-\mu|^2}}{2\sigma^2}$$

Computing P-values

- If we can estimate β and μ , then we can determine, for a given match score x, the probability that a random match with score x or greater would have occurred in the database.
- For sequence matches, a scoring system and database can be parameterized by two parameters, K and λ , related to β and μ .
 - It would be nice if we could compare hit significance without regard to the database and scoring system used!

Bit Scores

• Expected number of hits with score ≥ *S*:

$$E = Kmn e^{-\lambda s}$$

- Where m and n are the sequence lengths
- Normalize the raw score using:

$$S' = \frac{\lambda S - \ln K}{\ln 2}$$

- Obtains a "bit score" S', with a *standard set of units*.
- The new E-value is: $E = mn \ 2^{-S}$

P values and E values

- Blast reports *E*-values
- E = 5, E = 10 versus P = 0.993 and P = 0.99995
- When E < 0.01 P-values and E-values are nearly identical

BLAST Parameters

- Lowering the neighborhood word threshold (T) allows more distantly related sequences to be found, at the expense of increased noise in the results set.
- Raising the segment extension cutoff (X) returns longer extensions for each hit.
- Changing the minimum E-value changes the threshold for reporting a hit.

Aligning Protein Sequences

Sequence Alignments Revisited

- Scoring nucleotide sequence alignments was easier
 - Match score
 - Possibly different scores for transitions and transversions
- For amino acids, there are many more possible substitutions
- How do we score which substitutions are highly penalized and which are moderately penalized?
 - Physical and chemical characteristics
 - Empirical methods

Scoring Mismatches

- Physical and chemical characteristics
 - V → I Both small, both hydrophobic, conservative substitution, small penalty
 - V → K Small → large, hydrophobic → charged, large penalty
 - Requires some expert knowledge and judgement
- Empirical methods
 - How often does the substitution V → I occur in proteins that are known to be related?
 - Scoring matrices: PAM and BLOSUM

PAM Matrices

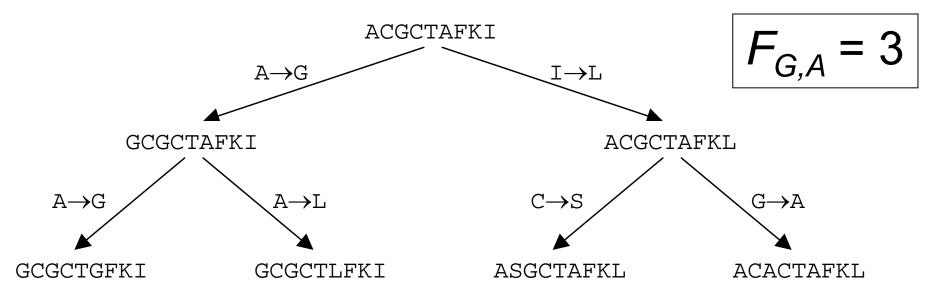
- PAM = "Point Accepted Mutation" interested only in mutations that have been "accepted" by natural selection
- Starts with a multiple sequence alignment of very similar (>85% identity) proteins. Assumed to be homologous
- Compute the relative mutability, m_i, of each amino acid
 - e.g. m_A = how many times was alanine substituted with anything else?

Relative Mutability

- ACGCTAFKI
 - GCGCTAFKI
 - ACGCTAFKL
 - GCGCTGFKI
 - GCGCTLFKI
 - ASGCTAFKL
 - ACACTAFKL
- Across all pairs of sequences, there are 28
 A → X substitutions
- There are 10 ALA residues, so $m_A = 2.8$

Pam Matrices

Construct a phylogenetic tree for the sequences in the alignment



- Calculate substitution frequences $F_{X,X}$
- Substitutions may have occurred either way, so
 A → G also counts as G → A.

Mutation Probabilities

• $M_{i,j}$ represents the score of $J \rightarrow I$ substitution.

$$M_{ij} = \frac{m_j F_{ij}}{\sum_{i} F_{ij}}$$

$$\frac{\text{ACGCTAFKI}}{\text{ACGCTAFKL}}$$

$$\frac{\text{ACGCTAFKL}}{\text{ACGCTAFKL}}$$

$$\frac{\text{ACGCTAFKL}}{\text{ACGCTAFKL}}$$

$$\frac{\text{ACGCTAFKL}}{\text{ACGCTAFKL}}$$

$$M_{G,A} = \frac{2.025}{2.7 \times 3}$$

The PAM matrix

- The entries, $R_{i,j}$ are the $M_{i,j}$ values divided by the frequency of occurrence, f_i , of residue i.
- $f_G = 10 \,\text{GLY} / 63 \,\text{residues} = 0.1587$
- $R_{G,A} = \log(2.025/0.1587) = \log(12.760) = 1.106$
- The log is taken so that we can add, rather than multiply entries to get compound probabilities.
- Log-odds matrix
- Diagonal entries are $1-m_i$

Interpretation of PAM matrices

- PAM-1 one substitution per 100 residues (a PAM unit of time)
- Multiply them together to get PAM-100, etc.
- "Suppose I start with a given polypeptide sequence M at time t, and observe the evolutionary changes in the sequence until 1% of all amino acid residues have undergone substitutions at time t+n. Let the new sequence at time t+n be called M'. What is the probability that a residue of type j in M will be replaced by i in M'?"

PAM Matrix Considerations

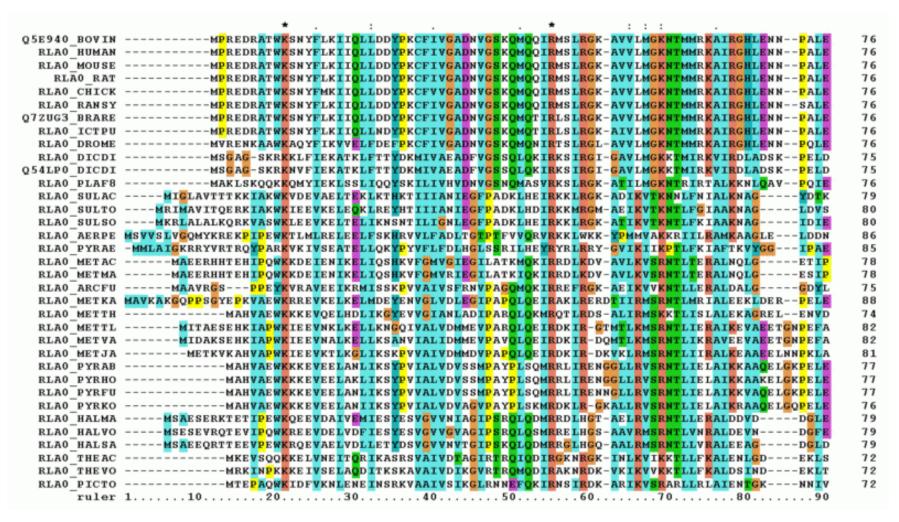
- If $M_{i,j}$ is very small, we may not have a large enough sample to estimate the real probability. When we multiply the PAM matrices many times, the error is magnified.
- PAM-1 similar sequences, PAM-1000 very dissimilar sequences

BLOSUM Matrix

- Starts by clustering proteins by similarity
- Avoids problems with small probabilities by using averages over clusters
- Numbering works opposite
 - BLOSUM-62 is appropriate for sequences of about 62% identity, while BLOSUM-80 is appropriate for more similar sequences.

Multiple Sequence Alignment

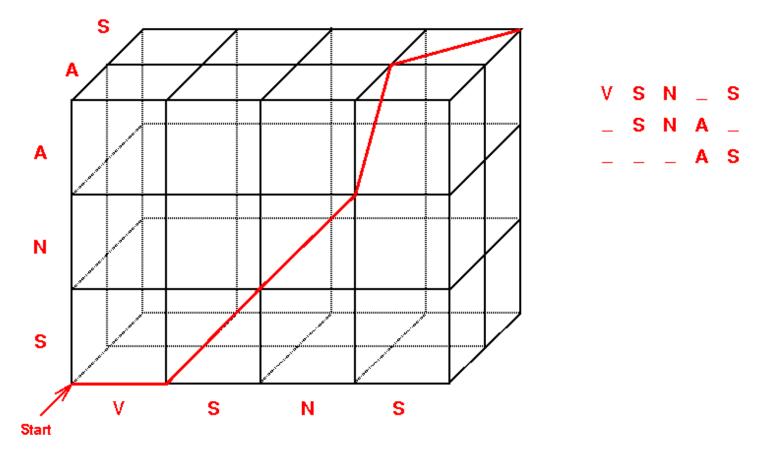
Multiple Alignment



Multiple Alignment

- The alignment of two sequences is relatively straightforward.
- Can we generalize our Dynamic Programming approach to multiple sequences?

Multiple Alignment



Turns out the complexity in exponential in the number of sequences.

Optimal Multiple alignment

- What is a suitable cost measure?
- Sum of scores of pairwise alignments?
- Most of the available multiple alignment
 methods use a progressive approach that
 makes pairwise alignments, averages them
 into a consensus (actually a profile), then adds
 new sequences one at a time to the aligned
 set.

Optimal Multiple alignment

- There can be various rules for building the consensus: simple majority rules, plurality by a specific fraction.
- This is an approximate method! (why?)

Multiple alignment: ClustalW

- CLUSTAL is the most popular multiple alignment program
- Gap penalties can be adjusted based on specific amino acid residues, regions of hydrophobicity, proximity to other gaps, or secondary structure.
- It can re-align just selected sequences or selected regions in an existing alignment
- It can compute phylogenetic trees from a set of aligned sequences.

Editing Multiple Alignments

- There are a variety of tools that can be used to modify and display a multiple alignment.
- An editor can also be used to make modifications by hand to improve biologically significant regions in a multiple alignment created by an alignment program.
- Examples of such editors include MACAW, SeqVu, and GeneDoc.

Multiple Alignments

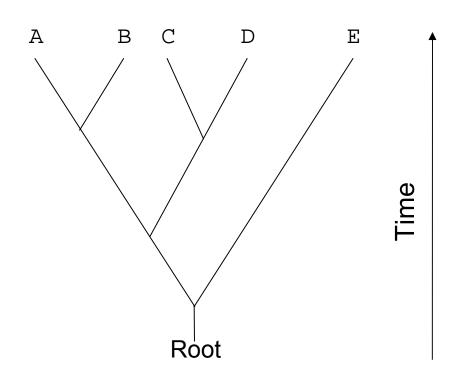
- Multiple Alignments are starting points for calculating phylogenetic trees
- Motifs and Profiles are calculated from multiple alignments

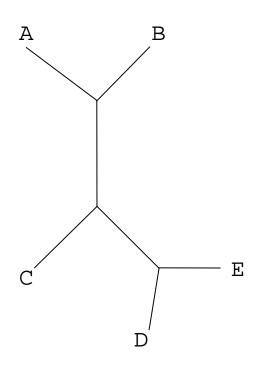
Phylogenetics

Phylogenetic Trees

Hypothesis about the relationship between organisms

Can be rooted or unrooted





Tree proliferation

$$N_R = \frac{2n - 3 \Omega}{2^{n-2} n - 2 \Omega}$$

$$N_U = \frac{2n - 5 \Omega}{2^{n-3} n - 3 \Omega}$$

Species	Number of Rooted Trees	Number of Unrooted Trees
2	1	1
3	3	1
4	15	3
5	105	15
6	34,459,425	2,027,025
7	213,458,046,767,875	7,905,853,580,625
Species	Number of Rooted Trees	Number of Unrooted Trees

Molecular phylogenetics

Specific genomic sequence variations (alleles) are much more reliable than phenotypic characteristics

More than one gene should be considered

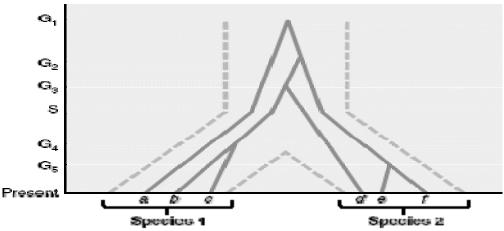


FIGURE 4.4 Individuals may actually appear to be more closely related to members of a species other than their own when only one gene is considered. Gene divergence events (G₁ through G₂) often occur before as well as after speciation events (S). The evolutionary history of gene divergence verilting in the six alleles denoted a through 1 is shown in which lines; speciation G.c., population splitting) is shown by broken lines. Individual d would actually appear to be more closely related to individuals in species 1 if only this locus were considered even though it is a member of species 2.

Distance matrix methods

	10	20	30	40	50
A:	GTGCTGCACGG	CTCAGTATA	GCATTTACCC	TTCCATCTTC	AGATCCTGAA
B:	ACGCTGCACGG	CTCAGTGCG	GTGCTTACCC	TCCCATCTTC	AGATCCTGAA
C:	GTGCTGCACGG	CTCGGCGCA	GCATTTACCC	TCCCATCTTC	AGATCCTATC
D:	GTATCACACGA	CTCAGCGCA	GCATTTGCCC	TCCCGTCTTC	AGATCCTAAA
E:	GTATCACATAG	CTCAGCGCA	GCATTTGCCC	TCCCGTCTTC	AGATCTAAAA

FIGURE 4.5 A five-way alignment of homologous DNA sequences.

Species	A	В	C	D
В	9	_	_	_
C	8	11	_	_
D	12	15	10	_
Species	A	В	C	D

UPGMA

Similar to average-link clustering

Merge the closest two groups

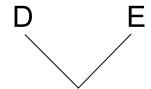
Replace the distances for the new, merged group with the average of the distance for the previous two groups

Repeat until all species are joined

UPGMA Step 1

Species	A	В	C	D
В	9	_	_	_
C	8	11	_	_
D	12	15	10	_
Species	A	В	C	D

Merge D & E

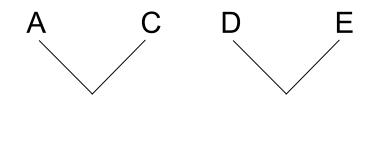


Species	A	В	<u>C</u>
В	9	_	_
C	8	11	_
Species	A	В	C

UPGMA Step 2

Species	A	В	C
В	9	_	_
C	8	11	
Species	A	В	C

Merge A & C

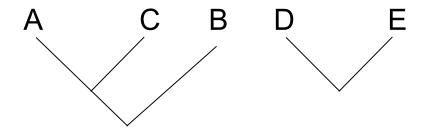


Species	В	AC
AC	10	_
Species	В	AC

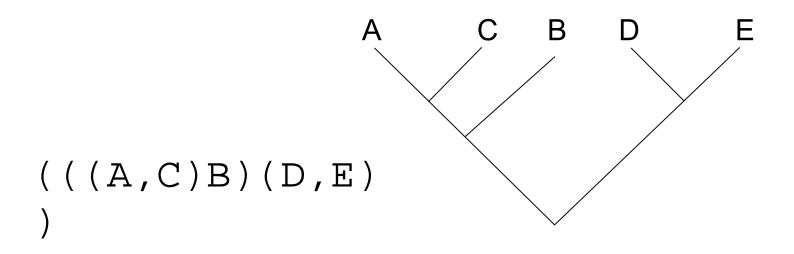
UPGMA Steps 3 & 4

Species	В	AC
AC	10	_
Species	В	AC

Merge B & AC



Merge ABC & DE



Parsimony approaches

Belong to the broader class of <u>character based</u> methods of phylogenetics

GCGGACG

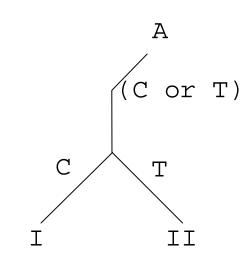
GTGGACG

Emphasize simpler, and thus more likely evolutionary pathways

(C or T)

C T

I



Parsimony methods

Enumerate all possible trees

Note the number of substitutions events invoked by each possible tree

Can be weighted by transition/transversion probabilities, etc.

Select the most parsimonious

Branch and Bound methods

Key problem – number of possible trees grows enormous as the number of species gets large Branch and bound – a technique that allows large numbers of candidate trees to be rapidly disregarded

Requires a "good guess" at the cost of the best tree

Parsimony – Branch and Bound

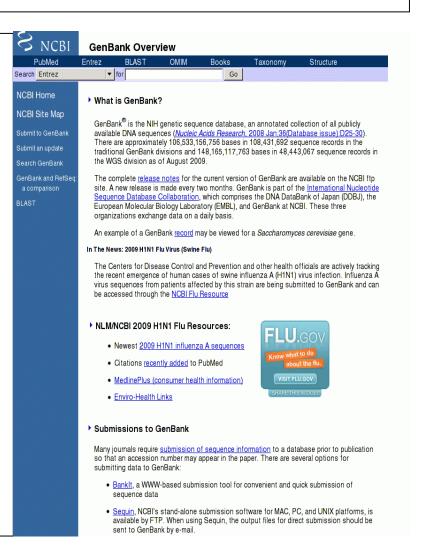
- Use the UPGMA tree for an initial best estimate of the minimum cost (most parsimonious) tree
- Use branch and bound to explore all feasible trees
- Replace the best estimate as better trees are found
- Choose the most parsimonious

Online Resources

Genbank:

http://www.ncbi.nlm.nih.gov/genbank/

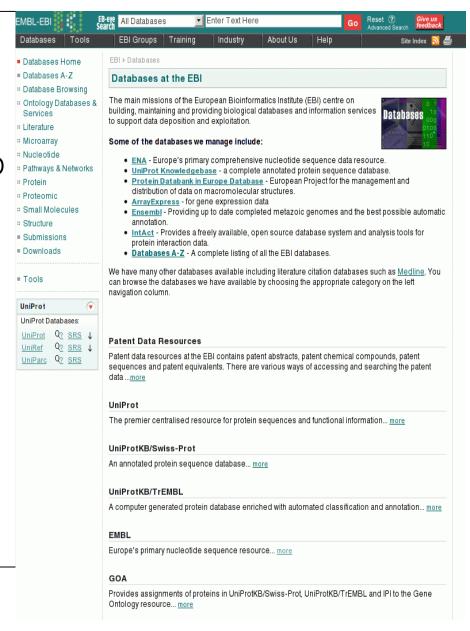
 NIH genetic sequence database, an annotated collection of all publicly available sequences



Online Resources: EBI

European Bioinformatics Institute

- http://www.ebi.ac.uk/D atabases/
- ENA Europe's primary comprehensive nucleotide sequence data resource.
- UniProt Knowledgebase a complete annotated protein sequence database.
- Protein Databank in Europe
 Database European Project for



Online Resources: DDBJ

DDBJ: DNA Data Bank of Japan

- http://www.ddbj.nig.ac.jp/
- Data exchange with EMBL/EBI, GenBank on a daily basis.
- Data across these databases is virtually identical (modulo curation practices)
- Virtually all sequence data in Japan is submitted through DDBJ.

Online Resources: SwissProt/UniProt

- http://ca.expasy.org/sprot/sprot-top.html
- Curated protein sequence databases
- UniProt/TrEMBL: annotated supplement to UniProt of EMBL nucleotides

Online Resources: Structure Databases

- PDB: http://www.rcsb.org/pdb/home/home.do Experimental structures of proteins, nucleic acids and assemblies
- NDB: http://ndbserver.rutgers.edu/ Nucleic acid structures
- SCOP: http://scop.mrc-lmb.cam.ac.uk/scop/ Structural classification of proteins
- Cambridge Structure
 Database: http://www.ccdc.cam.ac.uk/
 structure, visualization and analysis of organic molecules and metal-organic structures

Online Resources

- Motifs in protein structure and/or function PROSITE http://ca.expasy.org/prosite/
- Function
 EC Enzyme database http://ca.expasy.org/enzyme/
- Integrated databases

WIT

Entrez http://www.ncbi.nlm.nih.gov/sites/gquery

Online Resources

Fetching the sequences

- BLAST Search http://blast.ncbi.nlm.nih.gov/Blast.cgi
- Genbank Database Query Form at NCBI
- Entrez at NCBI
- Batch downloads via Batch Entrez
- NCSA Biology Workbench