

GQ: A Graph Toolkit for Multicore Environments

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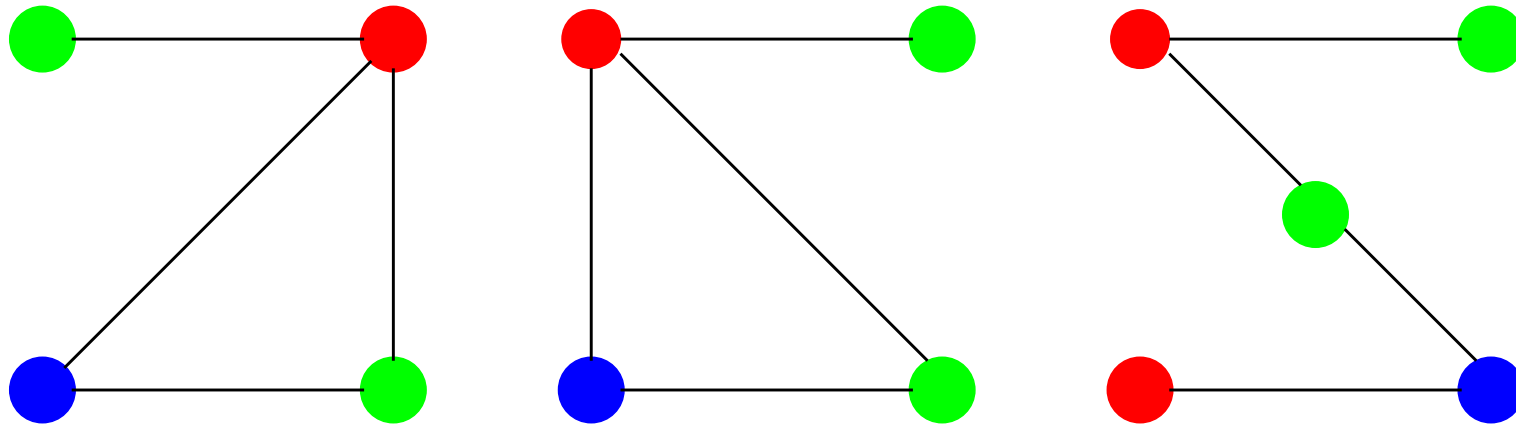
Graph Library Components

- Conservation in Networks
- Alignment of Networks
- Modularity in Networks
- Reputation/Rank
- Graph Grammars and Parsers
- Some Ongoing Work

Conservation in Networks

- Given a collection of networks (say, protein interaction networks belonging to different species), find **sub-networks** that are **common** to an **interesting** subset of these networks (Koyutürk, Grama, & Szpankowski, *ISMB*, 2004)
 - A sub-network is connected.
 - **Frequency**: The number of networks that contain a sub-network, is a coarse measure of **statistical significance**
- Requires solution of the intractable **subgraph isomorphism** problem
- Must be scalable to potentially **large** number of networks
- Networks are **large** (in the range of $10K$ edges and beyond)

Graph Analysis



Network database



Interaction patterns that are common to all networks

Problem Statement

- Given a set of **nodes** V , a set of **edges** E , and a **many-to-many** mapping from V to a set of **ortholog groups** $\mathcal{L} = \{l_1, l_2, \dots, l_n\}$, the corresponding interaction network is a **labeled graph** $G = (V, E, \mathcal{L})$.
 - $v \in V(G)$ is associated with a set of ortholog groups $L(v) \subseteq \mathcal{L}$.
 - $uv \in E(G)$ represents an interaction between u and v .
- S is a **sub-network** of G , i.e., $S \sqsubseteq G$ if there is an **injective** mapping $\phi : V(S) \rightarrow V(G)$ such that for all $v \in V(S)$, $L(v) \subseteq L(\phi(v))$ and for all $uv \in E(S)$, $\phi(u)\phi(v) \in E(G)$.

Computational Problem

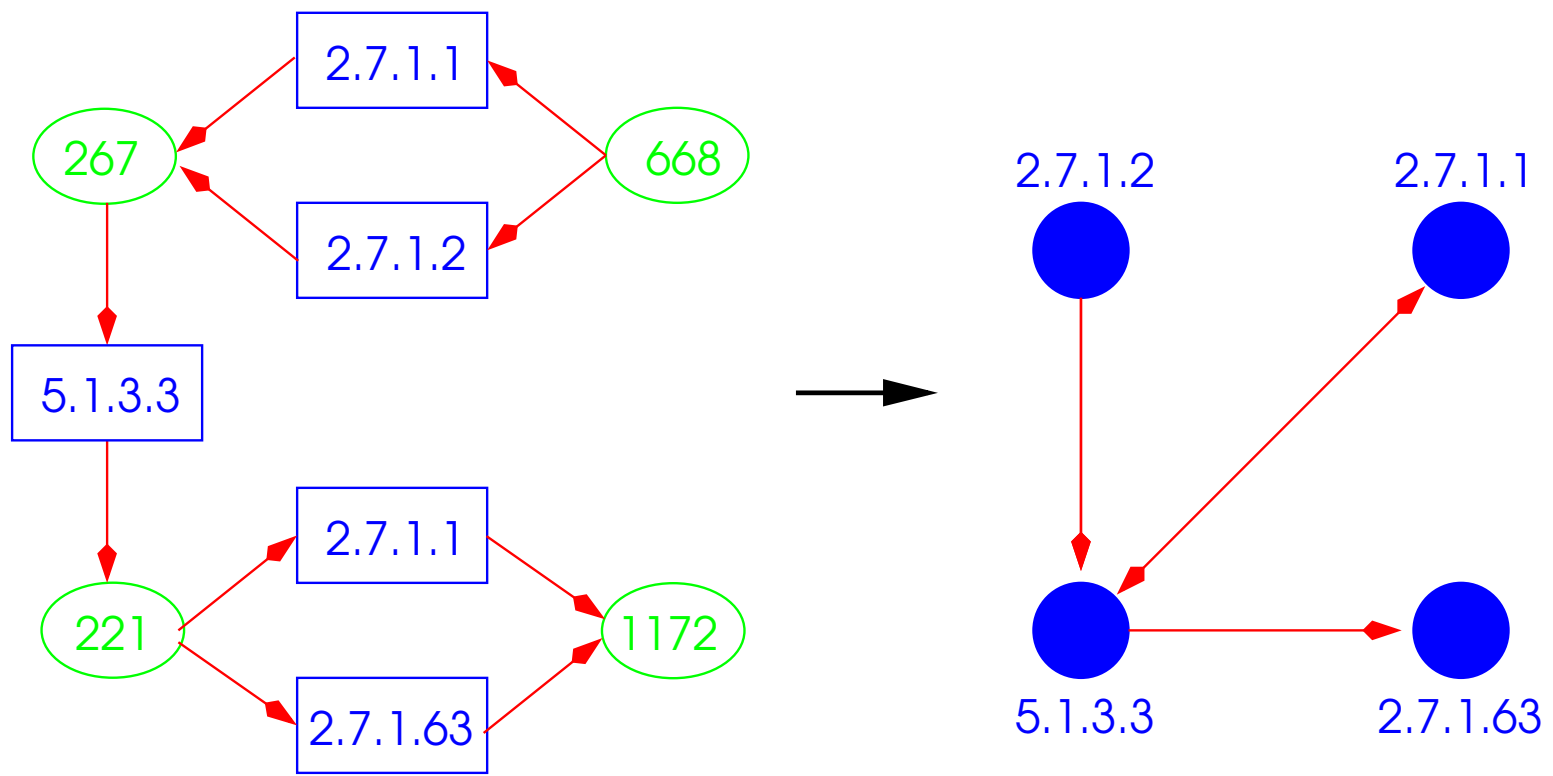
- Conserved sub-network discovery
 - **Instance:** A set of networks $\mathcal{G} = \{G_1 = (V_1, E_1, \mathcal{L}), G_2 = (V_2, E_2, \mathcal{L}), \dots, G_m = (V_m, E_m, \mathcal{L})\}$, and a **frequency** threshold σ^* .
 - **Problem:** Let $H(S) = \{G_i : S \subseteq G_i\}$ be the **occurrence** set of graph S . Find all **connected** subgraphs S such that $|H(S)| \geq \sigma^*$, i.e., S is a **frequent** subgraph in \mathcal{G} and for all $S' \supset S$, $H(S) \neq H(S')$, i.e., S is **maximal**.

Algorithmic Insight: Ortholog Contraction

- Contract orthologous nodes into a single node
- No subgraph isomorphism
 - Graphs are uniquely identified by their edge sets
- Key observation: Frequent sub-networks are preserved \Rightarrow No information loss
 - Sub-networks that are frequent in general graphs are also frequent in their ortholog-contracted representation
 - Ortholog contraction is a powerful pruning heuristic

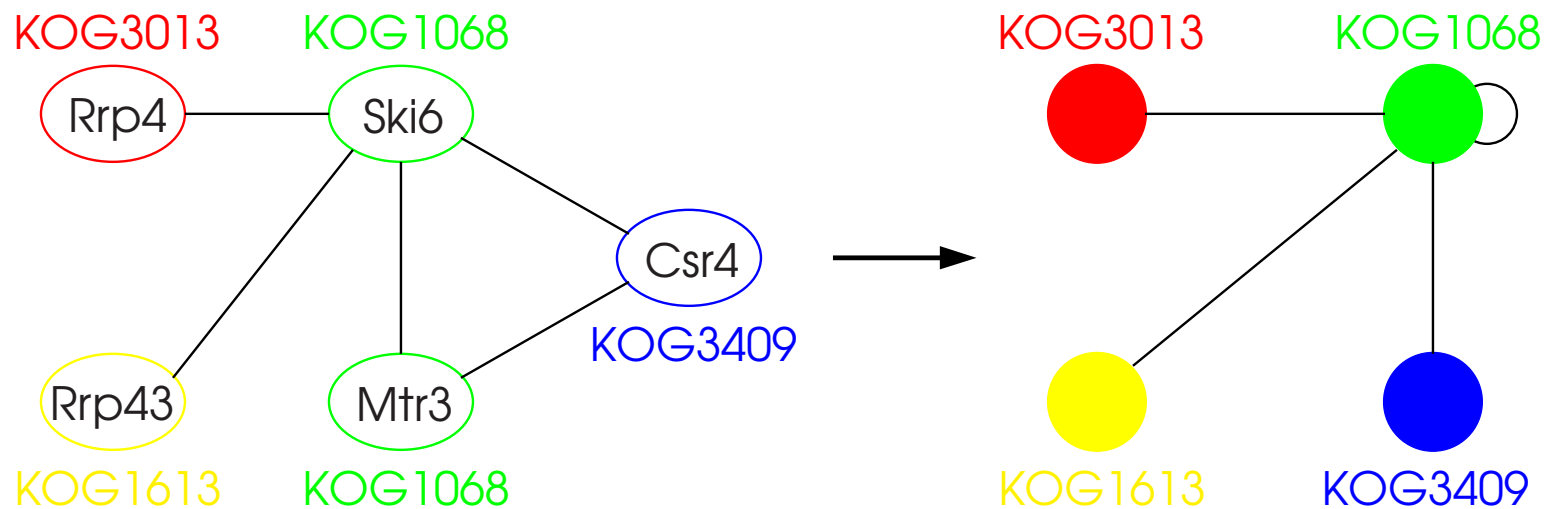
Ortholog Contraction in Real Applications (Metabolic Pathways)

- Directed hypergraph \rightarrow uniquely-labeled directed graph
 - Nodes represent enzymes
 - Global labeling by enzyme nomenclature (EC numbers)
 - A directed edge from one enzyme to the other implies that the second consumes a product of the first



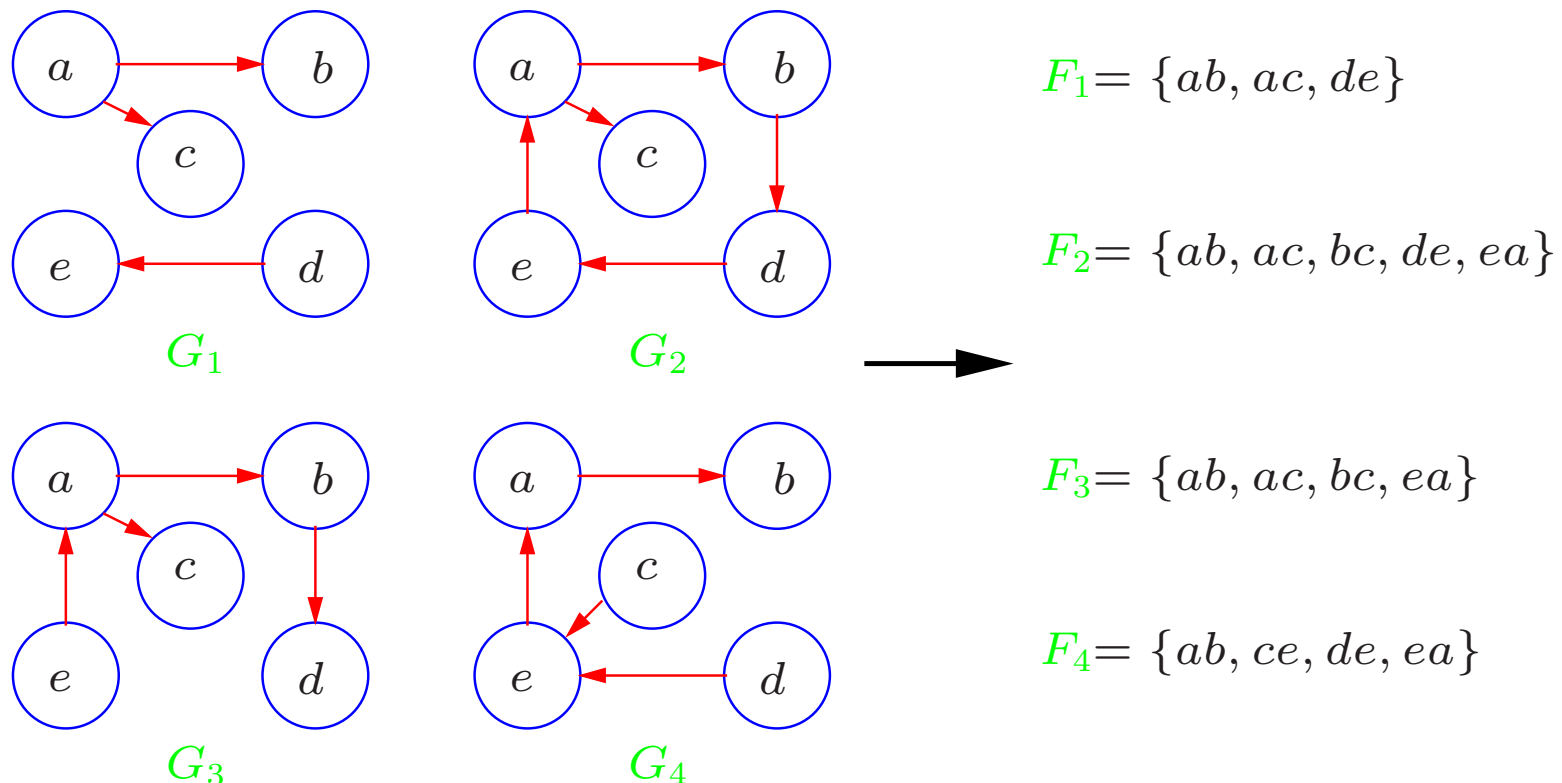
Ortholog Contraction in Real Applications (PPI Networks)

- Interaction between **proteins** → Interaction between **ortholog groups** or **protein families**

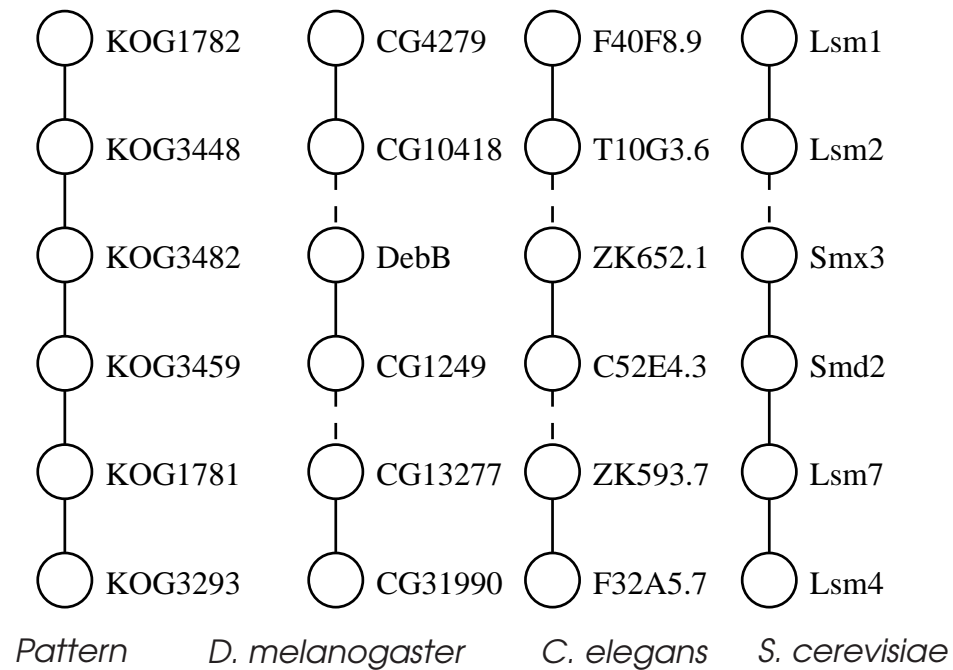


Simplifying the Graph Analysis Problem

- **Observation:** An ortholog-contracted graph is uniquely determined by the set of its edges.
 - Conserved **Sub-network** Discovery Problem \rightarrow Frequent **Edge set** Discovery Problem

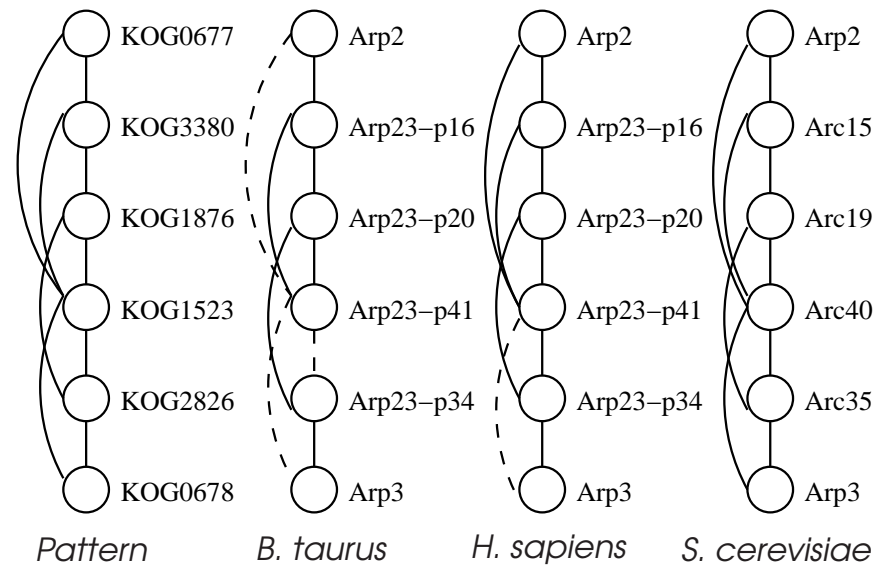


Conserved Protein Interaction Patterns



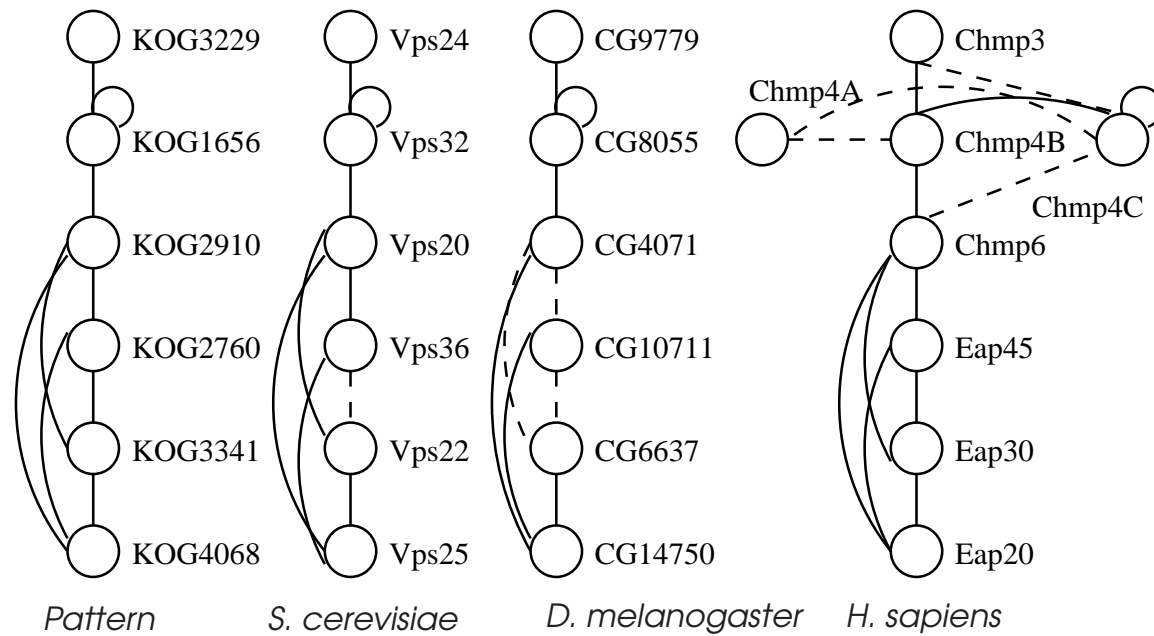
Small nuclear ribonucleoprotein complex ($p < 2e - 43$)

Conserved Protein Interaction Patterns



Actin-related protein Arp2/3 complex ($p < 9e - 11$)

Conserved Protein Interaction Patterns



Endosomal sorting ($p < 1e - 78$)

Runtime Characteristics

Dataset	Minimum Support (%)	Runtime (secs.)	Largest pattern	Number of patterns	Runtime 2 Cores	Runtime 4 Cores
Glutamate	12	0.10	13	39	0.08	0.07
	10	0.29	15	34	0.16	0.10
	8	0.99	15	56	0.58	0.37
Alanine	16	0.06	12	21	0.05	0.07
	12	1.06	16	25	0.57	0.33
	10	1.72	16	34	0.90	0.52

All times on a 2.66 MHz i7 Processor.

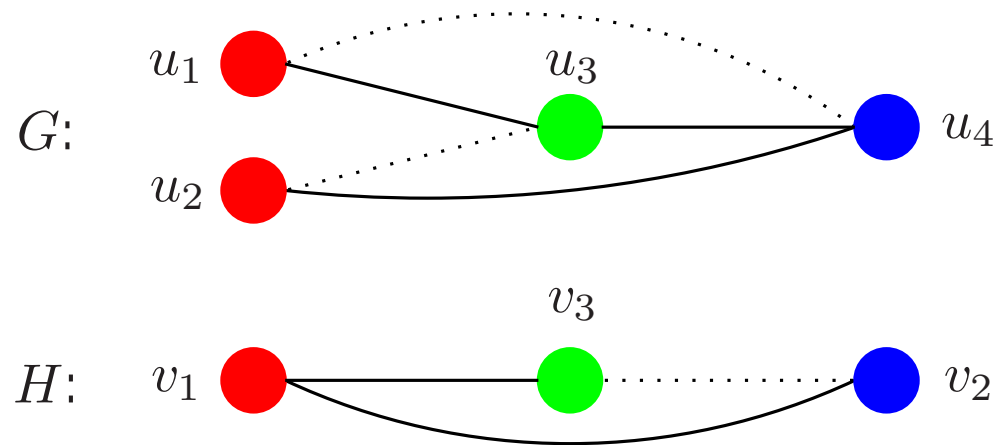
Alignment of Networks

- Given two networks, identify sub-networks that are similar to each other.
 - Optimization function
 - Mathematical modeling
- Existing algorithms
 - PathBLAST aligns pathways (linear chains) to simplify the problem while maintaining biological meaning (Kelley et al., *PNAS*, 2004)
 - NetworkBLAST compares conserved complex model with null model to identify significantly conserved subnets (Sharan et al., *J. Comp. Biol.*, 2005)

Match, Mismatch, and Duplication

- Establishing a Cost Measure

- A **match** $\in \mathcal{M}$ corresponds to two pairs of homologous nodes such that both pairs interact in both networks. A match is associated with **score** μ .
- A **mismatch** $\in \mathcal{N}$ corresponds to two pairs of homologous nodes such that only one pair is interacting. A mismatch is associated with **penalty** ν .
- A **duplication** $\in \mathcal{D}$ corresponds to a pair of homologous nodes in the same network. A duplication is associated with **score** δ .



Pairwise Alignment of Networks as an Optimization Problem

- Alignment score:

$$\sigma(\mathcal{A}(P)) = \sum_{M \in \mathcal{M}} \mu(M) - \sum_{N \in \mathcal{N}} \nu(N) + \sum_{D \in \mathcal{D}} \delta(D)$$

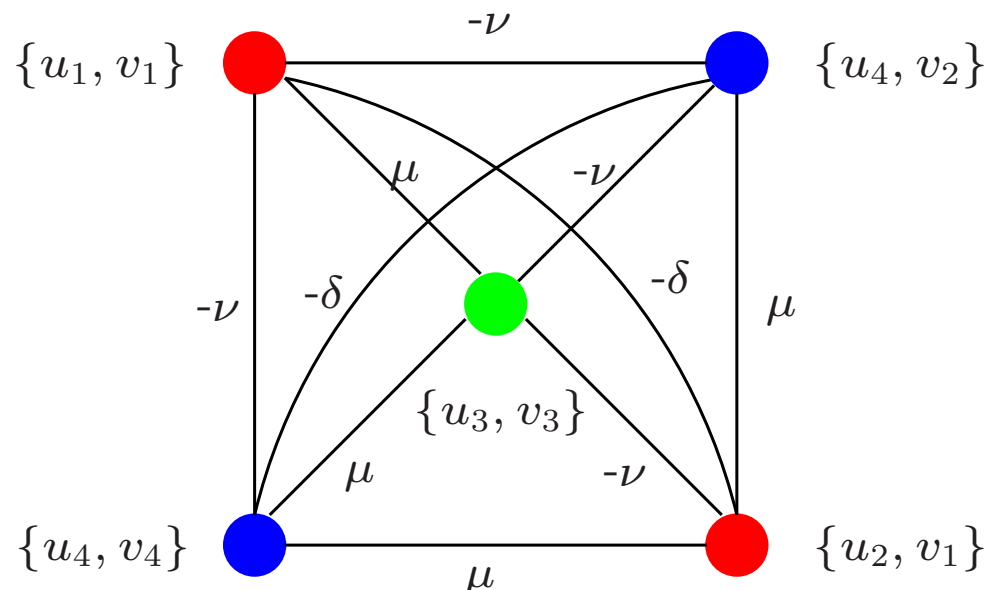
- Matches are rewarded for conservation of interactions
- Duplications are rewarded/penalized for functional conservation/differentiation after split
- Mismatches are penalized for divergence.

- Problem: Find all subnet pairs with significant alignment score

- A graph equivalent to BLAST

Weighted Alignment Graph

- $G(V, E)$: V consists of all pairs of homologous nodes $\mathbf{v} = \{u \in U, v \in V\}$
- An edge $\mathbf{v}\mathbf{v}' = \{uv\}\{u'v'\}$ in E is a
 - **match edge** if $uu' \in E$ and $vv' \in V$, with weight $w(\mathbf{v}\mathbf{v}') = \mu(uv, u'v')$
 - **mismatch edge** if $uu' \in E$ and $vv' \notin V$ or vice versa, with weight $w(\mathbf{v}\mathbf{v}') = -\nu(uv, u'v')$
 - **duplication edge** if $S(u, u') > 0$ or $S(v, v') > 0$, with weight $w(\mathbf{v}\mathbf{v}') = \delta(u, u')$ or $w(\mathbf{v}\mathbf{v}') = \delta(v, v')$



Maximum Weight Induced Subgraph Problem

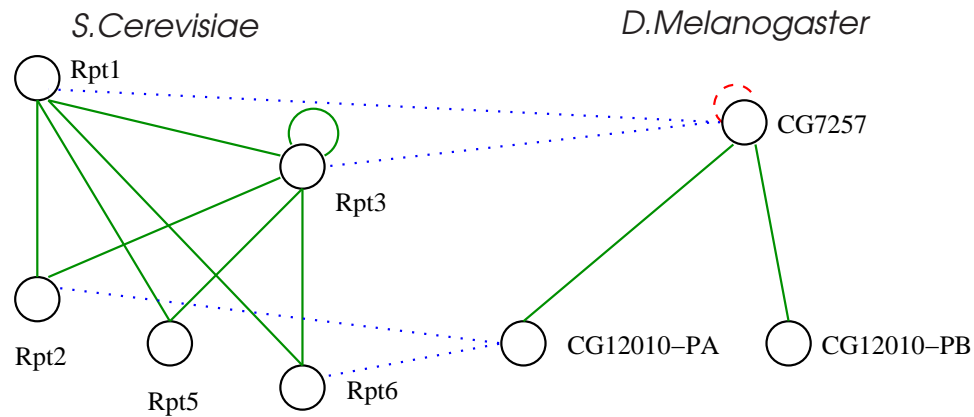
- **Definition:** (MAWISH)
 - Given graph $\mathcal{G}(\mathcal{V}, \mathcal{E})$ and a constant ϵ , find $\tilde{\mathcal{V}} \subseteq \mathcal{V}$ such that $\sum_{\mathbf{v}, \mathbf{u} \in \tilde{\mathcal{V}}} w(\mathbf{vu}) \geq \epsilon$.
 - NP-complete by reduction from Maximum-Clique
- **Theorem:** (MAWISH \equiv Pairwise alignment)
 - If $\tilde{\mathcal{V}}$ is a solution for the MAWISH problem on $\mathcal{G}(\mathcal{V}, \mathcal{E})$, then $P = \{\tilde{U}, \tilde{V}\}$ induces an alignment $\mathcal{A}(P)$ with $\sigma(\mathcal{A}) \geq \epsilon$, where $\tilde{\mathcal{V}} = \tilde{U} \times \tilde{V}$.
- **Solution:** Local graph expansion
 - Greedy graph growing + iterative refinement
 - Linear-time heuristic
- Source code available at
<http://www.cs.purdue.edu/pdsl/>

Alignment of Yeast and Fruit Fly PPI Networks

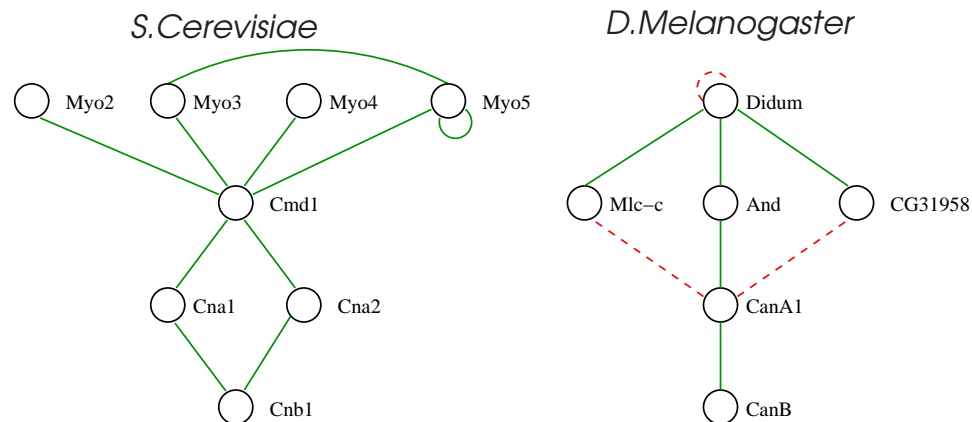
Rank	Score	<i>z</i> -score	# Proteins	# Matches	# Mismatches	# Dups.
1	15.97	6.6	18 (16, 5)	28	6	(4, 0)
	protein amino acid phosphorylation (69%) JAK-STAT cascade (40%)					
2	13.93	3.7	13 (8, 7)	25	7	(3, 1)
	endocytosis (50%) / calcium-mediated signaling (50%)					
5	8.22	13.5	9 (5, 3)	19	11	(1, 0)
	invasive growth (sensu <i>Saccharomyces</i>) (100%) oxygen and reactive oxygen species metabolism (33%)					
6	8.05	7.6	8 (5, 3)	12	2	(0, 1)
	ubiquitin-dependent protein catabolism (100%) mitosis (67%)					
21	4.36	6.2	9 (5, 4)	18	13	(0, 5)
	cytokinesis (100%, 50%)					
30	3.76	39.6	6 (3, 5)	5	1	(0, 6)
	DNA replication initiation (100%, 80%)					

Subnets Conserved in Yeast and Fruit Fly

Proteasome regulatory particle subnet



Calcium-dependent stress-activated signaling pathway



Runtime Characteristics

Dataset	Number of Patterns	Runtime (secs.)	2 Cores	4 Cores
Yeast/Fruit Fly	8	0.16	0.12	0.10
	13	1.80	1.02	0.68
	20	2.93	1.61	0.94

All times on a 2.66 MHz i7 Processor.

Analytical Assessment of Statistical Significance

- What is the **significance** of a **dense** component in a network?
- What is the **significance** of a **conserved** component in multiple networks?
- Existing techniques
 - Mostly computational (e.g., Monte-Carlo simulations)
 - Compute probability that **the** pattern exists rather than **a** pattern with **the property** (e.g., size, density) exists
 - **Overestimation of significance**

Random Graph Models

- Interaction networks generally exhibit **power-law** property (or exponential, geometric, etc.)
- Analysis simplified through **independence** assumption (Itzkovitz et al., *Physical Review*, 2003)
- Independence assumption may cause problems for networks with **arbitrary degree distribution**
- $P(uv \in E) = d_u d_v / |E|$, where d_u is expected degree of u , but generally $d_{\max}^2 > |E|$ for PPI networks
- Analytical techniques based on simplified models (Koyutürk, Grama, Szpankowski, *RECOMB*, 2006)
 - **Rigorous analysis** on $G(n, p)$ model
 - Extension to piecewise $G(n, p)$ to **capture network characteristics** more accurately

Significance of Dense Subgraphs

- A subnet of r proteins is said to be ρ -dense if $F(r) \geq \rho r^2$, where $F(r)$ is the number of interactions between these r proteins
- What is the expected size of the largest ρ -dense subgraph in a random graph?
 - Any ρ -dense subgraph with larger size is statistically significant!
- $G(n, p)$ model
 - n proteins, each interaction occurs with probability p
 - Simple enough to facilitate rigorous analysis
 - If we let $p = d_{\max}/n$, largest ρ -dense subgraph in $G(n, p)$ stochastically dominates that in a graph with arbitrary degree distribution
- Piecewise $G(n, p)$ model
 - Few proteins with many interacting partners, many proteins with few interacting partners
 - Captures the basic characteristics of PPI networks
 - Analysis of $G(n, p)$ model immediately generalized to this model

Largest Dense Subgraph

- **Theorem:** If G is a random graph with n nodes, where every edge exists with probability p , then

$$\lim_{n \rightarrow \infty} \frac{R_\rho}{\log n} = \frac{1}{\kappa(p, \rho)} \quad (pr.), \quad (1)$$

where

$$\kappa(p, \rho) = \rho \log \frac{\rho}{p} + (1 - \rho) \log \frac{1 - \rho}{1 - p}. \quad (2)$$

More precisely,

$$P(R_\rho \geq r_0) \leq O\left(\frac{\log n}{n^{1/\kappa(p, \rho)}}\right), \quad (3)$$

where

$$r_0 = \frac{\log n - \log \log n + \log \kappa(p, \rho)}{\kappa(p, \rho)} \quad (4)$$

for large n .

Piecewise $G(n, p)$ model

- The size of largest dense subgraph is still proportional to $\log n / \kappa$ with a constant factor depending on **number of hubs**
- **Model:**

$$P(uv \in E(G)) = \begin{cases} p_h & \text{if } u, v \in V_h \\ p_l & \text{if } u, v \in V_l \\ p_b & \text{if } u \in V_h, v \in V_l \text{ or } u \in V_l, v \in V_h \end{cases}$$

- **Result:**
Let $n_h = |V_h|$. If $n_h = O(1)$, then $P(R_n(\rho) \geq r_1) \leq O\left(\frac{\log n}{n^{1/\kappa(p_l, \rho)}}\right)$,
where

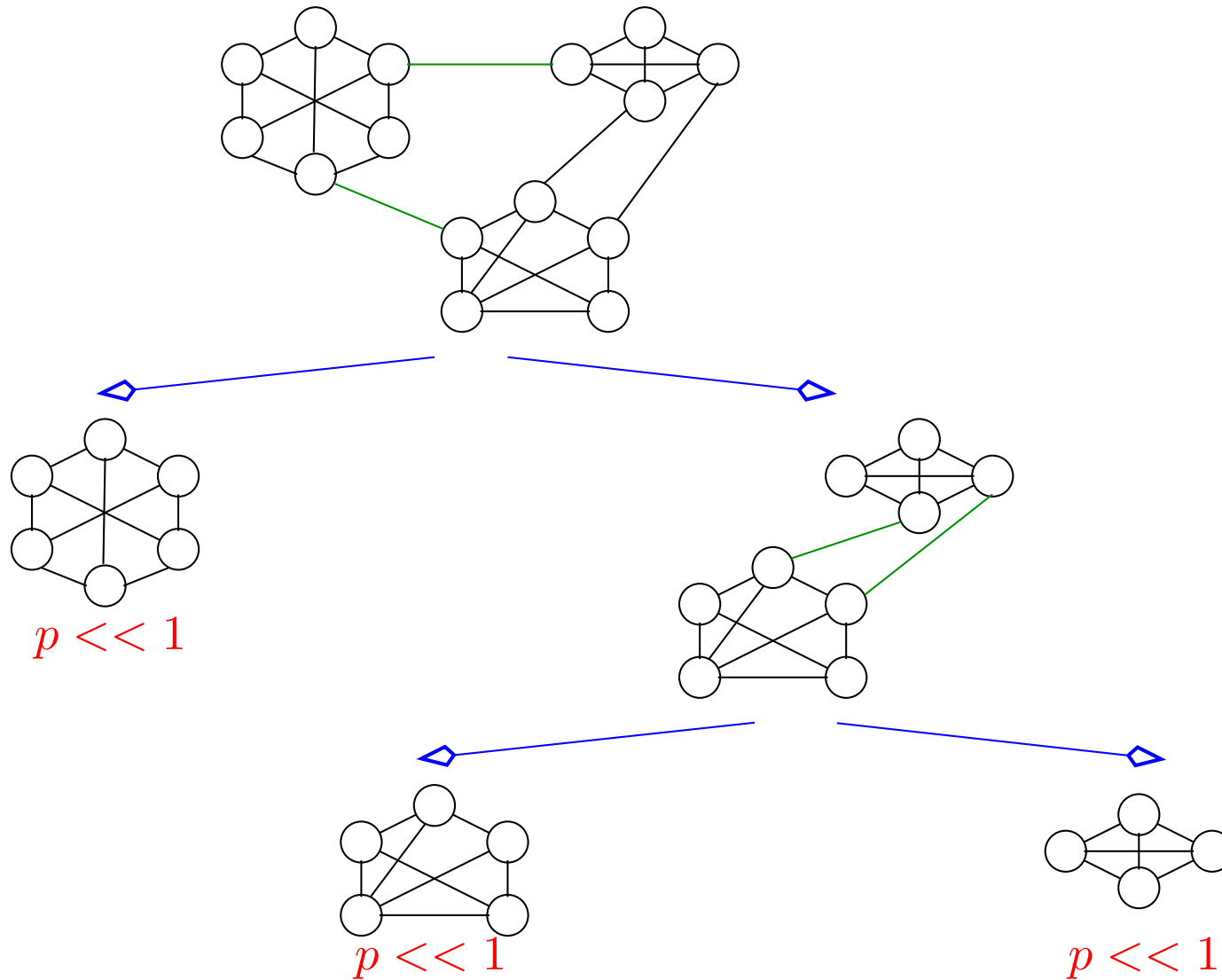
$$r_1 = \frac{\log n - \log \log n + 2n_h \log B + \log \kappa(p_l, \rho) - \log e + 1}{\kappa(p_l, \rho)}$$

and $B = \frac{p_b q_l}{p_l} + q_b$, where $q_b = 1 - p_b$ and $q_l = 1 - p_l$.

Algorithms Based on Statistical Significance

- Identification of **topological modules**
- Use **statistical significance** as a **stopping criterion** for graph clustering heuristics
- HCS Algorithm (*Hartuv & Shamir, Inf. Proc. Let., 2000*)
 - Find a minimum-cut bipartitioning of the network
 - If any of the parts is **dense enough**, record it as a dense cluster of proteins
 - Else, further partition them recursively
- **SIDES**: Use **statistical significance** to determine whether a subgraph is sufficiently dense
 - For given number of proteins and interactions between them, we can determine whether those proteins induce a significantly dense subnet

SIDES Algorithm



SIDES is available at <http://www.cs.purdue.edu/pds1>

Performance of SIdES

- Biological relevance of identified clusters is assessed with respect to **Gene Ontology (GO)**
 - Estimate the statistical significance of the **enrichment** of each GO term in the cluster
- **Quality** of the clusters with respect to GO annotations
 - Assume cluster C containing n_C genes is associated with term T that is attached to n_T genes and n_{CT} of genes in C are attached to T
 - **specificity** = $100 \times n_{CT}/n_C$
 - **sensitivity** = $100 \times n_{CT}/n_T$

	SIdES			MCode		
	Min.	Max.	Avg.	Min.	Max.	Avg.
Specificity (%)	43.0	100.0	91.2	0.0	100.0	77.8
Sensitivity (%)	2.0	100.0	55.8	0.0	100.0	47.6

Comparison of SIdES with MCode (Bader & Hogue, *BMC Bioinformatics*, 2003)

Runtime Characteristics

Dataset	Number of Clusters	Runtime (secs.)	2 Cores	4 Cores
Yeast PPI	11	4.80	2.64	1.60
	18	7.32	3.70	1.99
	26	10.19	5.61	2.90

All times on a 2.66 MHz i7 Processor.

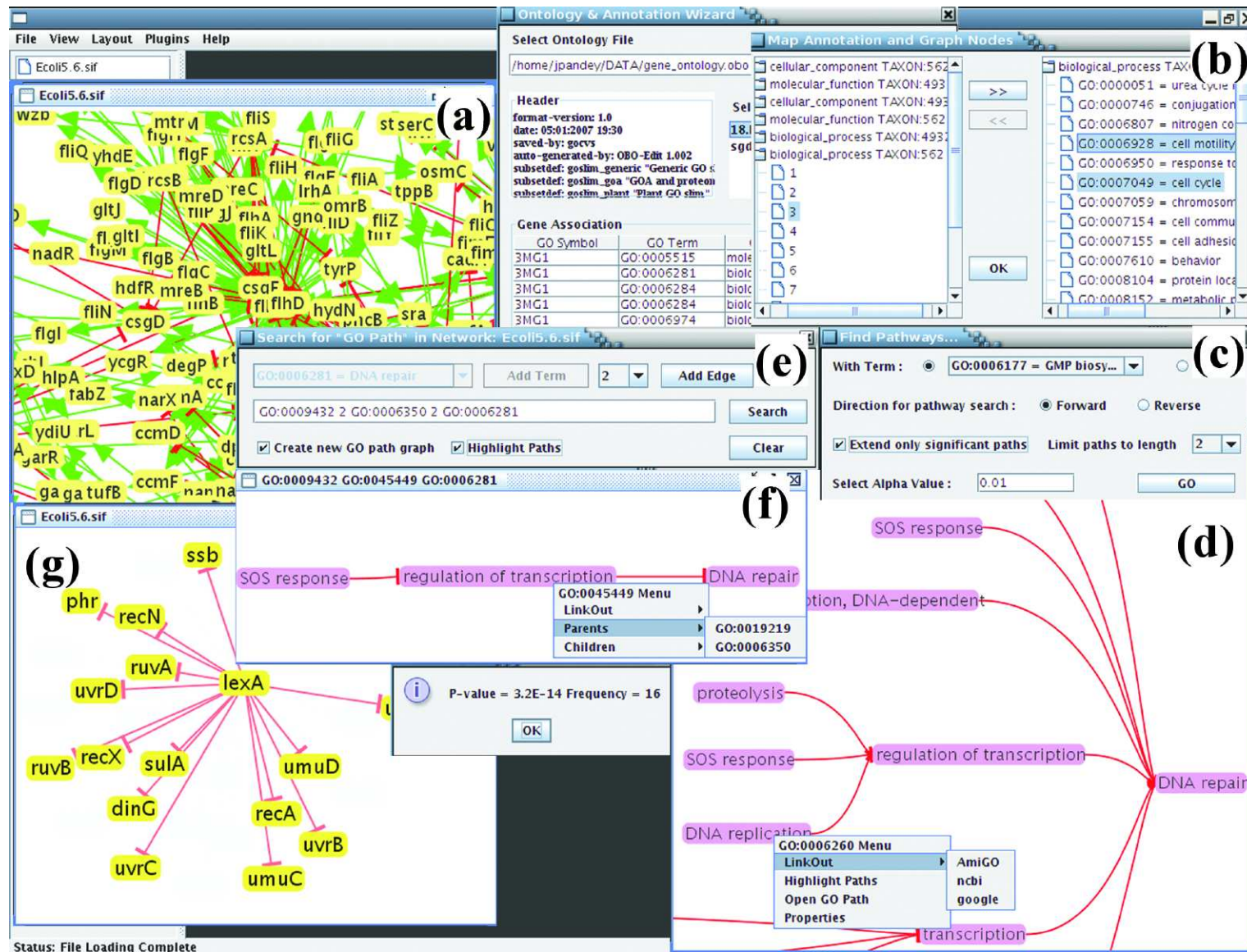
Functional Annotation of Pathways

- Identifying Significant Pathways
- Annotations and Metrics
- Application to Protein/Domain Interaction Networks
- Implementation and Results

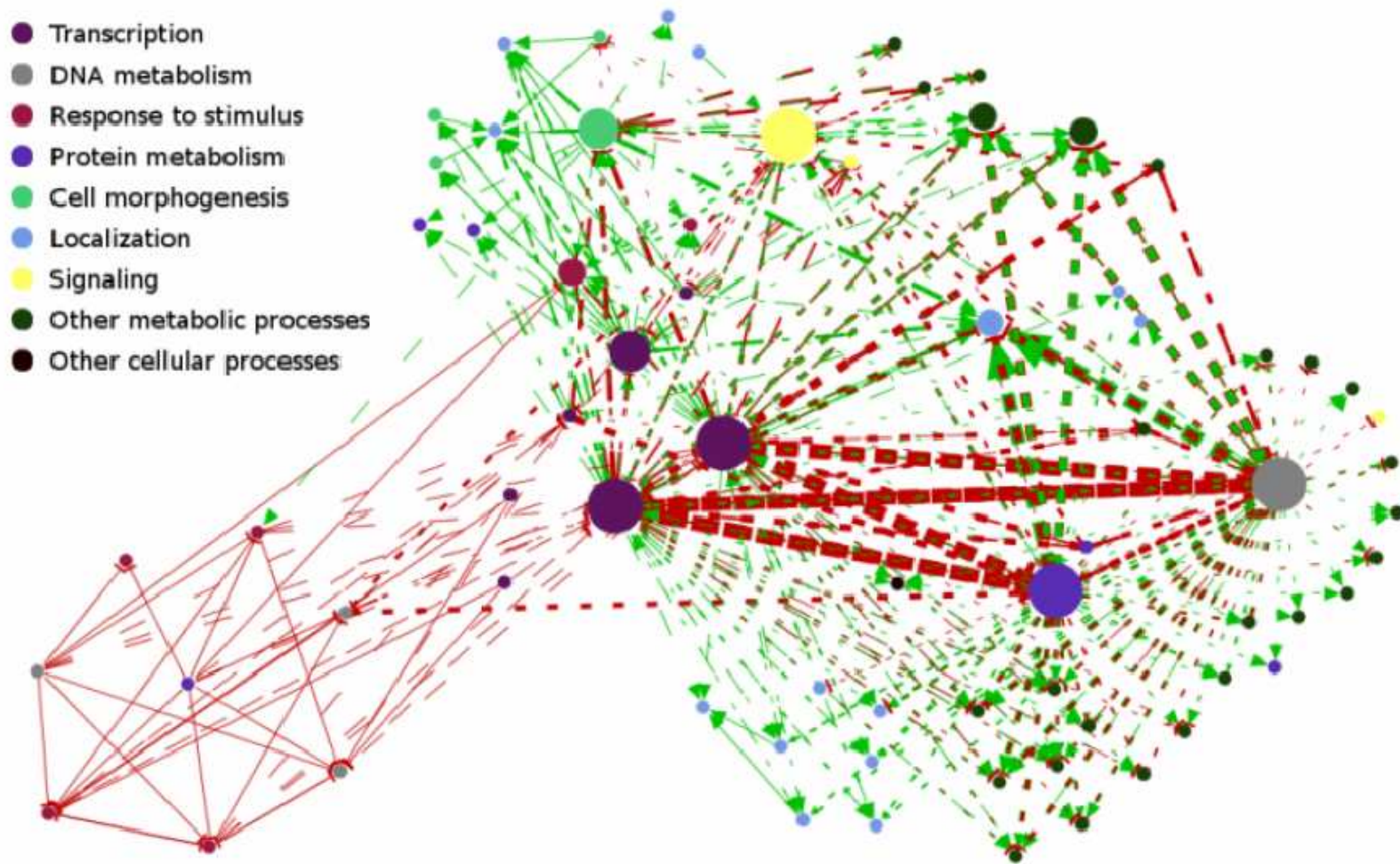
Node Annotation

- Node annotation is in the form of an ontology.
- For instance, Gene Ontology provides a library of molecular annotations (we refer to each annotation class as a functional attribute).

Narada Functionality

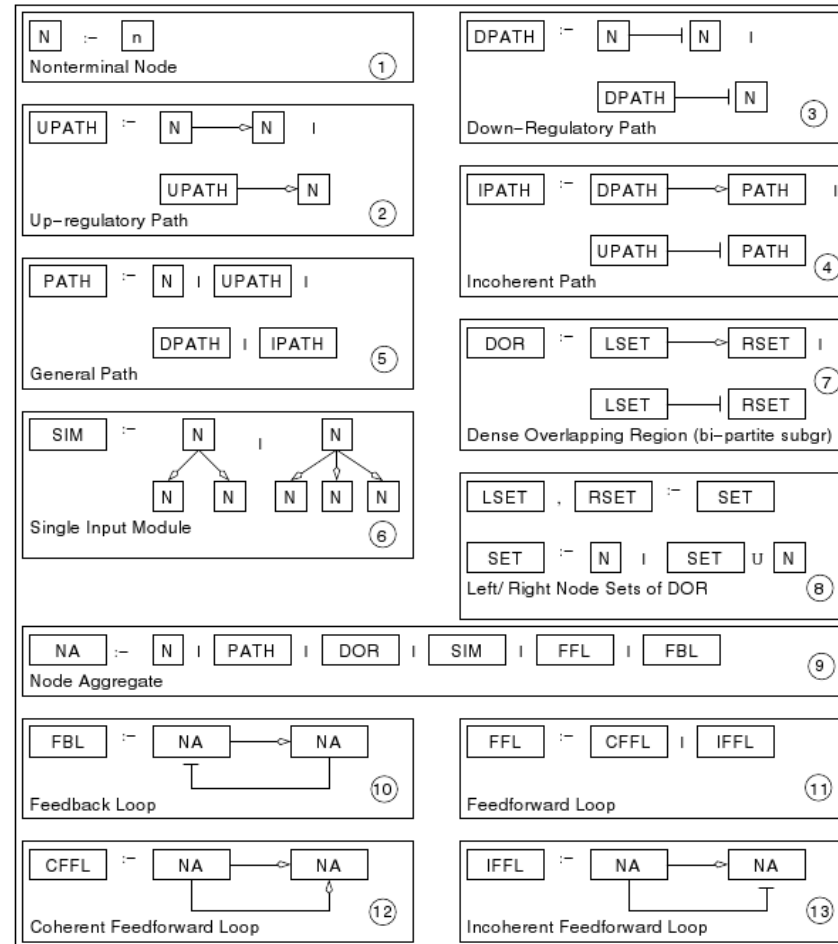


Narada Network Annotation

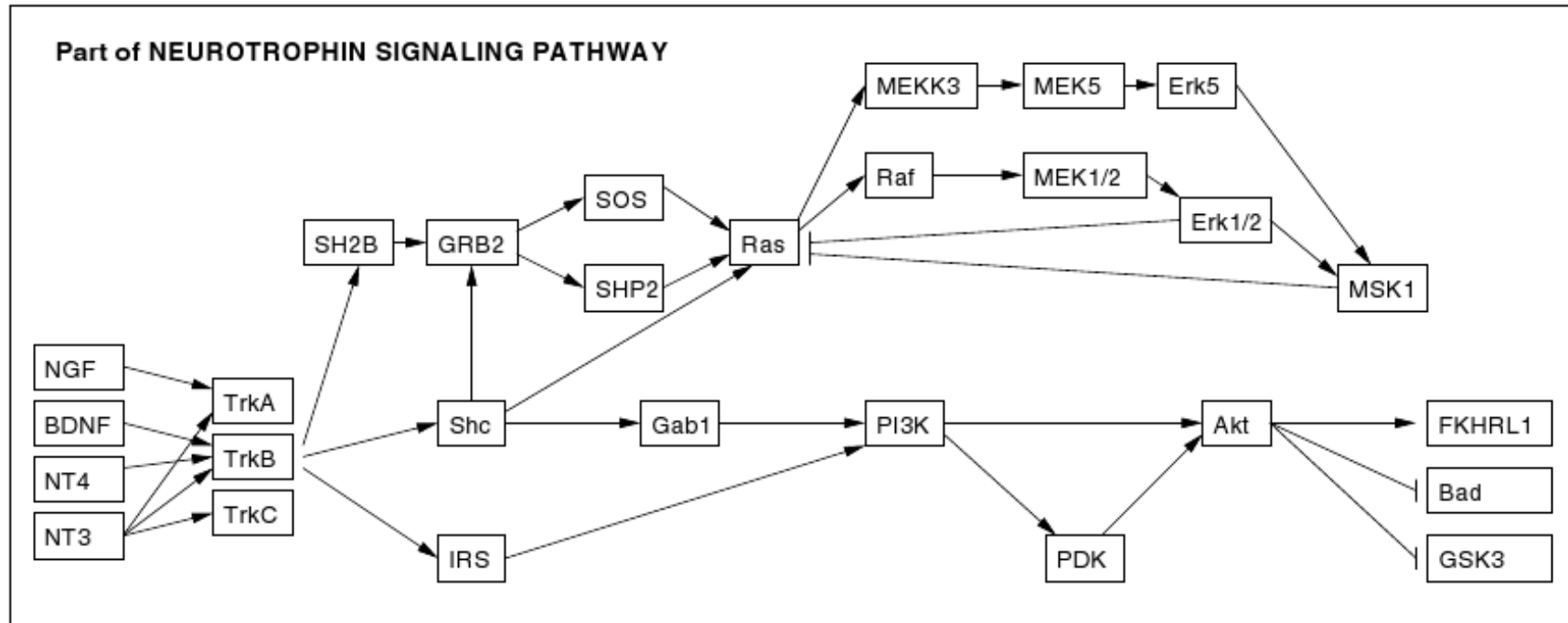


* Parallel implementation ongoing.

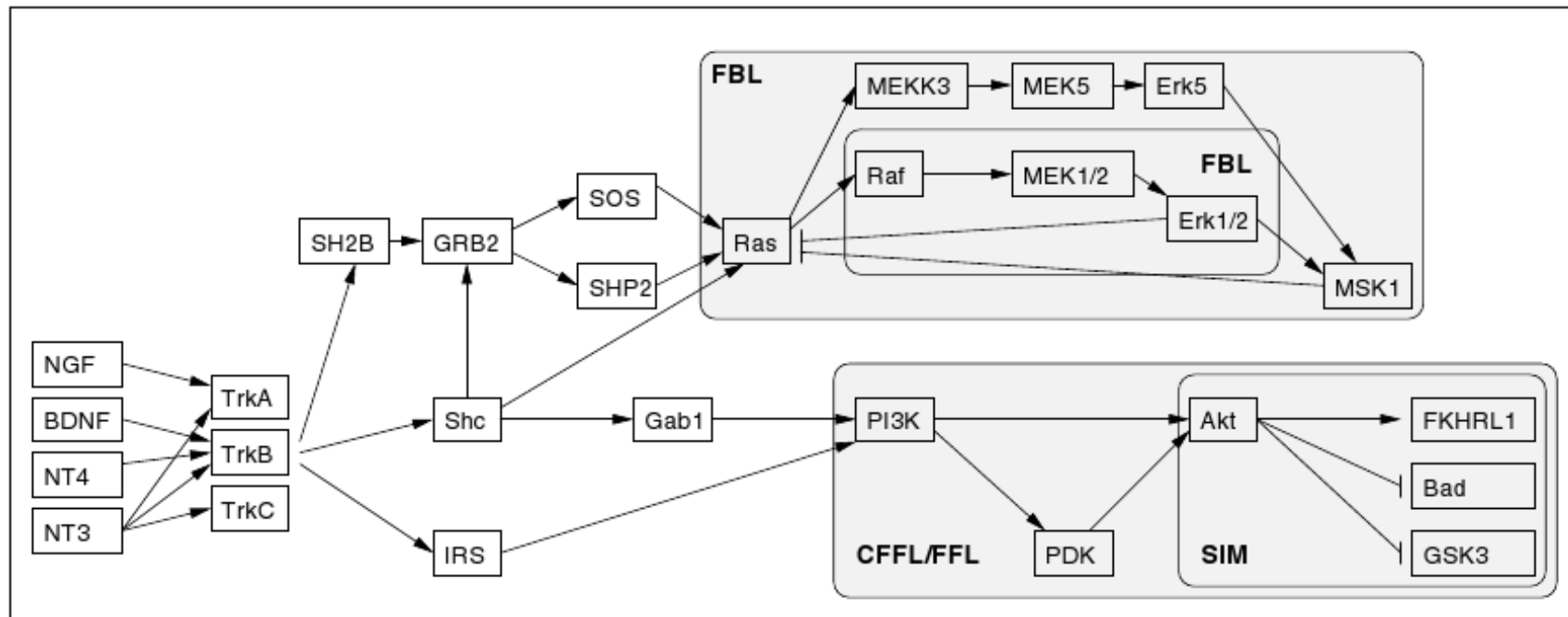
Graph Grammars and Parsing



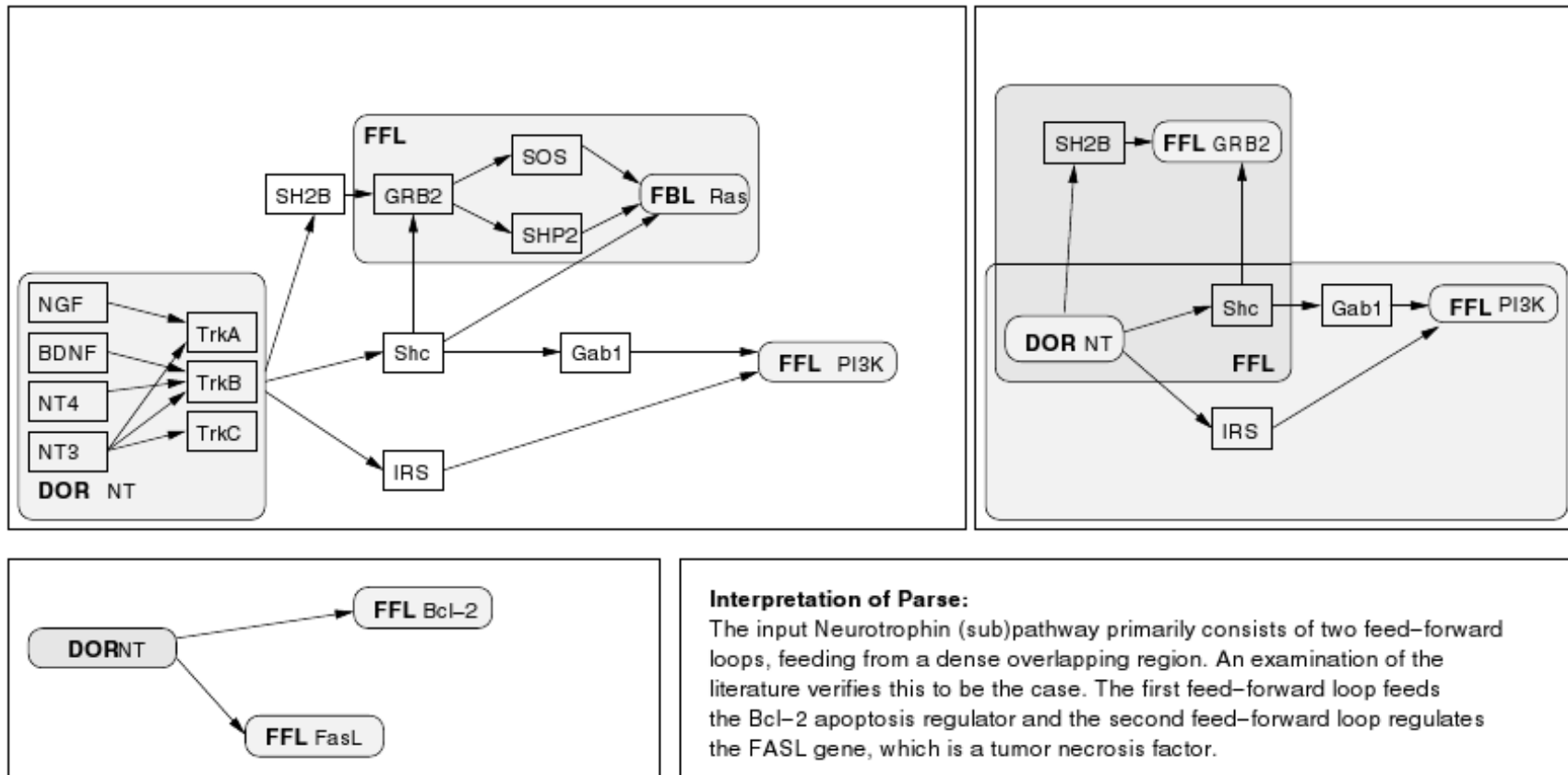
Graph Grammars and Parsing



Graph Grammars and Parsing



Graph Grammars and Parsing



Graph Grammars and Parsing: Status

- Serial parser complete.
- Parallel parser currently in implementation.
- Grammar inference currently under implementation.

Science of Information

“An NSF STC focused on post-Shannon Information theory.”