# Surface Triangulation and Voronoi Regions 

CS334<br>Spring 2022

Daniel G. Aliaga<br>Department of Computer Science<br>Purdue University

[Slides with help from Michael Kazhdan @ JHU, Ioannis Stamos @ CUNY, and
Profs. Shmuel Wimer and Andy Mirzaian]

## Motivation

- Time of flight
- Structured light
- Stereo images
- Shape from shading
- Etc.
http://graphics.stanford.edu/projects/mich/



## Motivation

## Surface reconstruction

Geometry
processing

etc.


## Marching Cubes

If the function is sampled on a regular voxel grid, we can independently triangulate each voxel.


## Marching Cubes



## Ball-pivoting



## Bernardini et al., IBM



Fixed-radius ball "rolling" over points selects subset of alpha-shape.

## Pivoting in 2D


(a)

(b)

(c)

Circle of radius $\rho$ pivots from point to point, connecting them with edges.
When sampling density is low, some of the edges will not be created, leaving holes. (c) When the curvature of the manifold is larger than $1 / \rho$, some of the points will not be reached by the pivoting ball, and features will be missed.

## Ball Pivoting Algorithm



## Ball Pivoting Algorithm



## Implicit Representation

Another option is representing a 3D model by an implicit function for:

- Reconstruction
- Fluid Dynamics
- 3D Texturing


Kazhdan 2005


## Implicit Function Fitting

Given point samples:

- Define a function with value zero at the points.
- Extract the zero isosurface.


Sample points


## Triangulation Complexity (in general)

- Theorem: (Gary et. al. 1978) A simple n-vertex polygon can be triangulated in $\mathrm{O}(\mathrm{nlogn})$ time and $O(n)$ storage
- The problem has been studied extensively between 1978 and 1991, when in 1991 Chazelle presented an $\underline{O(n)}$ time complexity algorithm.


## Delaunay Triangulation

- Another very popular algorithm...
- But first, Voronoi Diagrams...
- Relevant Conversation:
- Captain Kirk: "Spock! Which tricorder tower (i.e., cell phone) should I be using?"
- Commander Spock: "Logically, the closest one, Jim."
- How do you do that?


## Where to place cell phone towers? or Which cell phone tower should I use?



## Cell phone towers

$$
P=\left\{p_{1}, p_{2}, \ldots, p_{n}\right\} \text { a set of } n \text { points in the plane. }
$$



## Voronoi Diagram

$P=\left\{p_{1}, p_{2}, \ldots, p_{n}\right\}$ a set of $n$ points in the plane.

## Voronoi Diagram:



Voronoi(P): \# regions = n, \# edges $\leq 3 n-6$, \# vertices $\leq 2 n-5$.

## Delaunay Triangulation = Dual of the Voronoi Diagram



DT(P): \# vertices = n, \# edges $\leq 3 n-6, ~ \# ~ t r i a n g l e s ~ \leq 2 n-5 . ~$

## Delaunay Triangulation



Delaunay triangles have the "empty circle" property.

## Voronoi Diagram and Delaunay Triangulation



## VD Properties

- Each Voronoi region $\mathrm{V}\left(\mathrm{p}_{\mathrm{i}}\right)$ is a convex polygon (possibly unbounded).
- $V\left(p_{i}\right)$ is unbounded $\Leftrightarrow p_{i}$ is on the boundary of $\mathrm{CH}(P)$.
- Consider a Voronoi vertex $\mathrm{v}=\mathrm{V}\left(\mathrm{p}_{\mathrm{i}}\right) \cap \mathrm{V}\left(\mathrm{p}_{\mathrm{j}}\right) \cap \mathrm{V}\left(\mathrm{p}_{\mathrm{k}}\right)$.

Let $C(v)=$ the circle centered at $v$ passing through $p_{i}, p_{j}, p_{k}$.

- $C(v)$ is circumcircle of Delaunay Triangle ( $p_{i}, p_{j}, p_{k}$ ).
- $C(v)$ is an empty circle, i.e., its interior contains no other sites of $P$.


## Computing Delaunay Triangulation

- Many algorithms: O(nlogn)
- Lets use flipping:
- Recall: A Delaunay Triangulation is a set of triangles T in which each edge of T possesses at least one empty circumcircle.
- Empty: A circumcircle is said to be empty if it contains no nodes of the set V


## What is a flip?



A non-Delaunay edge flipped

## Flip Algorithm

- ??


## Flip Algorithm

1. Let V be the set of input vertices.
2. $\mathrm{T}=$ Any Triangulation of V .
3. Repeat until all edges of $T$ are Delaunay edges.
a. Find a non-delaunay edge that is flippable b. Flip

Naïve Complexity: $\mathrm{O}\left(\mathrm{n}^{2}\right)$

## Locally Delaunay $\rightarrow$ Globally Delaunay

- If $T$ is a triangulation with all its edges locally Delaunay, then T is the Delaunay triangulation.
- Proof by contradiction:
- Let all edges of $T$ be locally Delaunay but an edge of $T$ is not Delaunay, so flip it...


## Flipping

- Other flipping ideas?


## Randomized Incremental Flipping

- Complexity can be O(nlogn)


## Fortune’s Algorithm

- "A sweepline algorithm for Voronoi "Algorithms", 1987, O(nlogn) https://www.youtube.com/watch?v=k2P9yWSMaXE Pseudocode:

```
add a site event in the event queue for each site
while the event queue is not empty
pop the top event
if the event is a site event
    insert a new arc in the beachline
    check for new circle events
else
    create a vertex in the diagram
    remove the shrunk arc from the beachline
    delete invalidated events
    check for new circle events
```


## Wave Propagation View

 waves.

Let Time be the $3^{\text {rd }}$ Dimension


All sites have identical opaque cones.

Let Time be the $3^{\text {rd }}$ Dimension


All sites have identical opaque cones.
cone $(p) \cap$ cone $(q)=$ vertical hyperbola $h(p, q)$.
vertical projection of $h(p, q)$ on the xy base plane is $P B(p, q)$.

## Voronoi Diagrams

- http://alexbeutel.com/webgl/voronoi.html


## Voronoi Diagram

- http://www.raymondhill.net/voronoi/rhill-voronoi-demo5.html


## Examples Triangulations



## And Beyond...

- Not "relaxation" but more general:
- Recall: Reaction Diffusion
- https://pmneila.github.io/jsexp/grayscott/
- Textures:


