

Recursive Algorithms Implemented in Python

Reading

- Zelle 13.2
- Recursive algorithms (pseudocode, Algorithms slides)

Recursion

- A problem solving paradigm
- An approach for designing algorithms
- Given a recursive algorithm, there is always an equivalent non-recursive algorithm
 - Recursive algorithm often simpler

Recursion problem solving paradigm

- You don't solve the problem directly
- Split the problem until it becomes trivial
- Compute solution to problem by combining solutions of sub-problems

3 elements of recursive algorithm

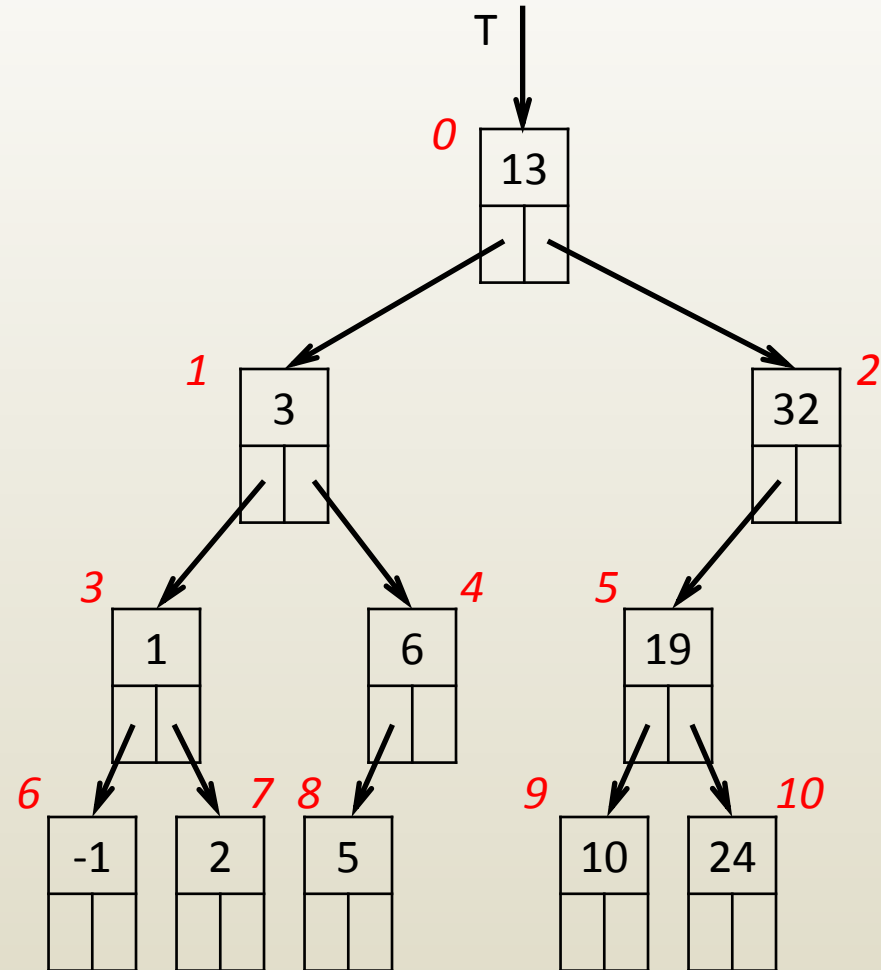
- Termination condition
 - At some point recursion has to stop
 - For example, don't go beyond leafs
 - Leafs don't have children, referring to children leafs causes algorithm to crash
- Recursive call
 - Algorithm calls itself on subsets of the input data
 - One or more recursive calls
 - For binary tree we had two recursive calls, one for each child
- Work done before, between, and after recursive calls

Examples

- We have seen several recursive algorithms
 - Binary tree traversal
 - Counting number of nodes in binary trees
 - Evaluation of arithmetic expression stored in a binary tree
 - Printing out arithmetic expression stored in a binary tree
 - Computing factorial of n
 - Finding the minimum element of an array of numbers
 - Binary search
- Now let's implement these and other recursive algorithms in Python

Binary trees in Python

- An array of triples (i.e. an array of arrays with 3 elements, a 2-D array of nx3 in size)
- One triple per node
 - Data, index of left child, index of right child
- First triple corresponds to root
- An index of -1 corresponds to null (i.e. no such child)
- Example for tree on the right
 - `[[13, 1, 2], [3, 3, 4], [32, 5, -1], [1, 6, 7], [6, 8, -1], [19, 9, 10], [-1, -1, -1], [2, -1, -1], [5, -1, -1], [10, -1, -1], [24, -1, -1]]`
 - Black numbers are data
 - Red numbers are indices



iClicker Question

- Given a binary tree T encoded using an array of triples, how does one test whether a node with index $currNode$ is a leaf?
 - A. $T[currNode][0] == -1$ and $T[currNode][1] == -1$
 - B. $T[currNode][1] == -1$ and $T[currNode][2] == -1$
 - C. $T[currNode][1] == -1$ or $T[currNode][2] == -1$
 - D. Either B or C
 - E. None of the above

Counting nodes in binary tree

// PSEUDOCODE

Input:

T // link to root of binary tree

Output:

// count of nodes

CountBTR(T)

if T == NULL

return 0

endif

return **CountBTR(T->left) + 1 +**
CountBTR (T->right)

endCountBTR

Python

currNode is index of current node

def CountBTR(T, currNode):

if currNode == -1:

return 0

left = T[currNode][1]

right = T[currNode][2]

return **CountBTR(T, left) + 1 +**
CountBTR(T, right)

Arithmetic expression evaluation

Input:

T // link to root of arithm. expr. tree

Output:

// value of arithmetic expression

EvalAEBTR(T)

if T-> left == NULL

return T->val

endif

switch T->symbol

case '+': **return** EvalAEBTR(T->left) +
 EvalAEBTR(T->right)

case '-': **return** EvalAEBTR(T->left) -
 EvalAEBTR(T->right)

case '*': **return** EvalAEBTR(T->left) *
 EvalAEBTR(T->right)

case '/': **return** EvalAEBTR(T->left) /
 EvalAEBTR(T->right)

endswitch

endEvalAEBTR

def EvalAEBTR(T, currNode):

 left = T[currNode][1]

 right = T[currNode][2]

 if left == -1:

return T[currNode][0]

if T[currNode][0] == '+':

return EvalAEBTR(T, left) +
 EvalAEBTR(T, right)

elif T[currNode][0] == '-':

return EvalAEBTR(T, left) -
 EvalAEBTR(T, right)

elif T[currNode][0] == '*':

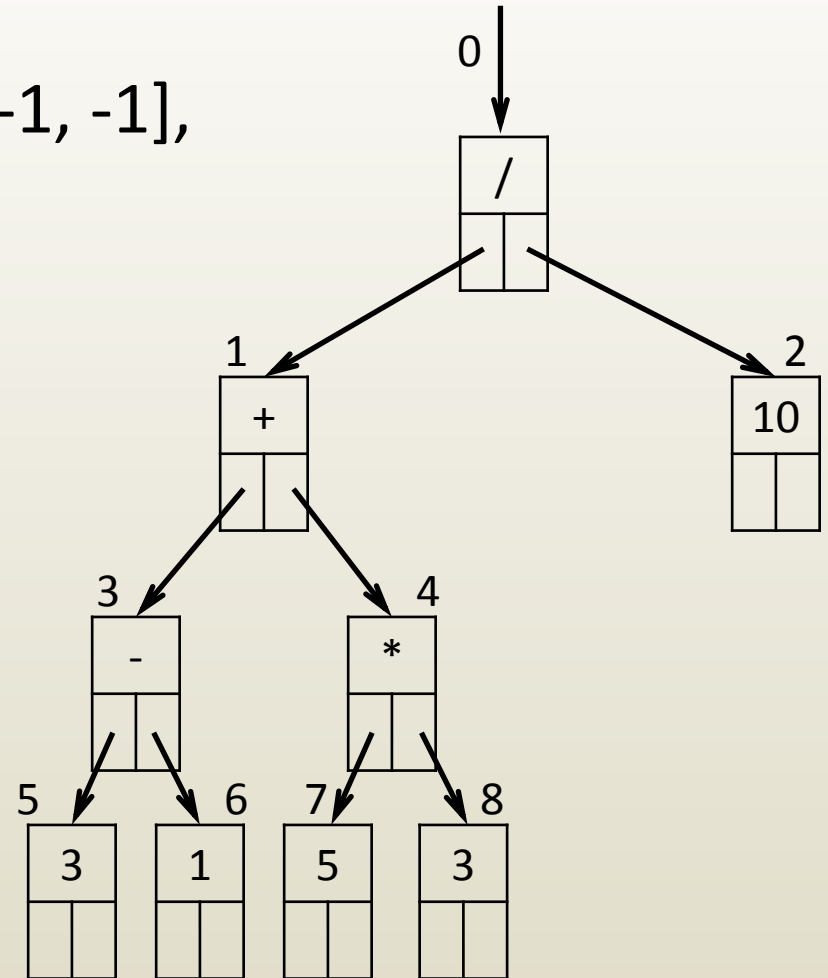
return EvalAEBTR(T, left) *
 EvalAEBTR(T, right)

elif T[currNode][0] == '/':

return EvalAEBTR(T, left) /
 EvalAEBTR(T, right)

Example

- $T = [['/ ', 1, 2], ['+ ', 3, 4], [10, -1, -1], ['- ', 5, 6], ['* ', 7, 8], [3, -1, -1], [1, -1, -1], [5, -1, -1], [3, -1, -1]]$



Printing out arithmetic expression

Input:

T // link to root of arithmetic expression
binary tree

Output:

// expression printed out with
parentheses

PrintAEBTR(T)

```
    if T->left != NULL
        print "("
        PrintAEBTR(T->left)
    endif
    print T->string
    if T->right != NULL
        PrintAEBTR(T->right)
    endif
endPrintAEBTR
```

Example: (((3-1)+(5*3))/10) for tree on
previous slide

def PrintAEBTR(T, currNode):

```
    if currNode == -1:
        return
    left = T[currNode][1]
    right = T[currNode][2]
    if left != -1:
        print('(' , end="")
        PrintAEBTR(T, left)
    print(T[currNode][0], end="")
    if right != -1:
        PrintAEBTR(T, right)
    print(')', end="")
```

Drawing Binary Tree

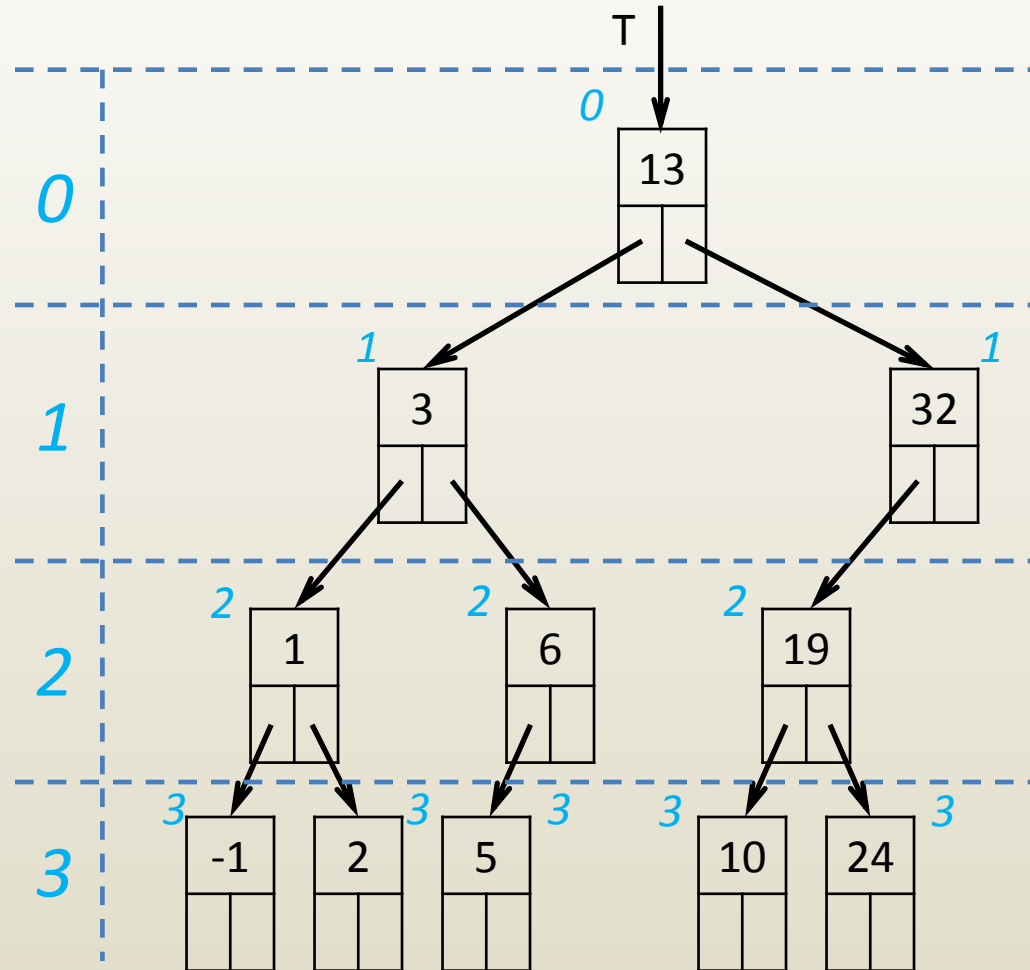
- Needed ingredients
 - How to draw
 - a node—*as a circle*
 - an edge between parent and child—*as a line segment*
 - text data—*using text object*
 - Where to draw
 - Node vertical coordinate given by depth in tree
 - Node horizontal coordinate given by in-order traversal index

2 step approach

- Step 1: set coordinates of each node
 - Recursive
 - Coordinates appended to node triple, which is now defined by 5 numbers: [data, leftIndex, rightIndex, hCoord, vCoord]
- Step 2: actually draw the nodes (and edges)
 - Recursive
 - Uses coordinates stored during step 1

Setting node coordinates

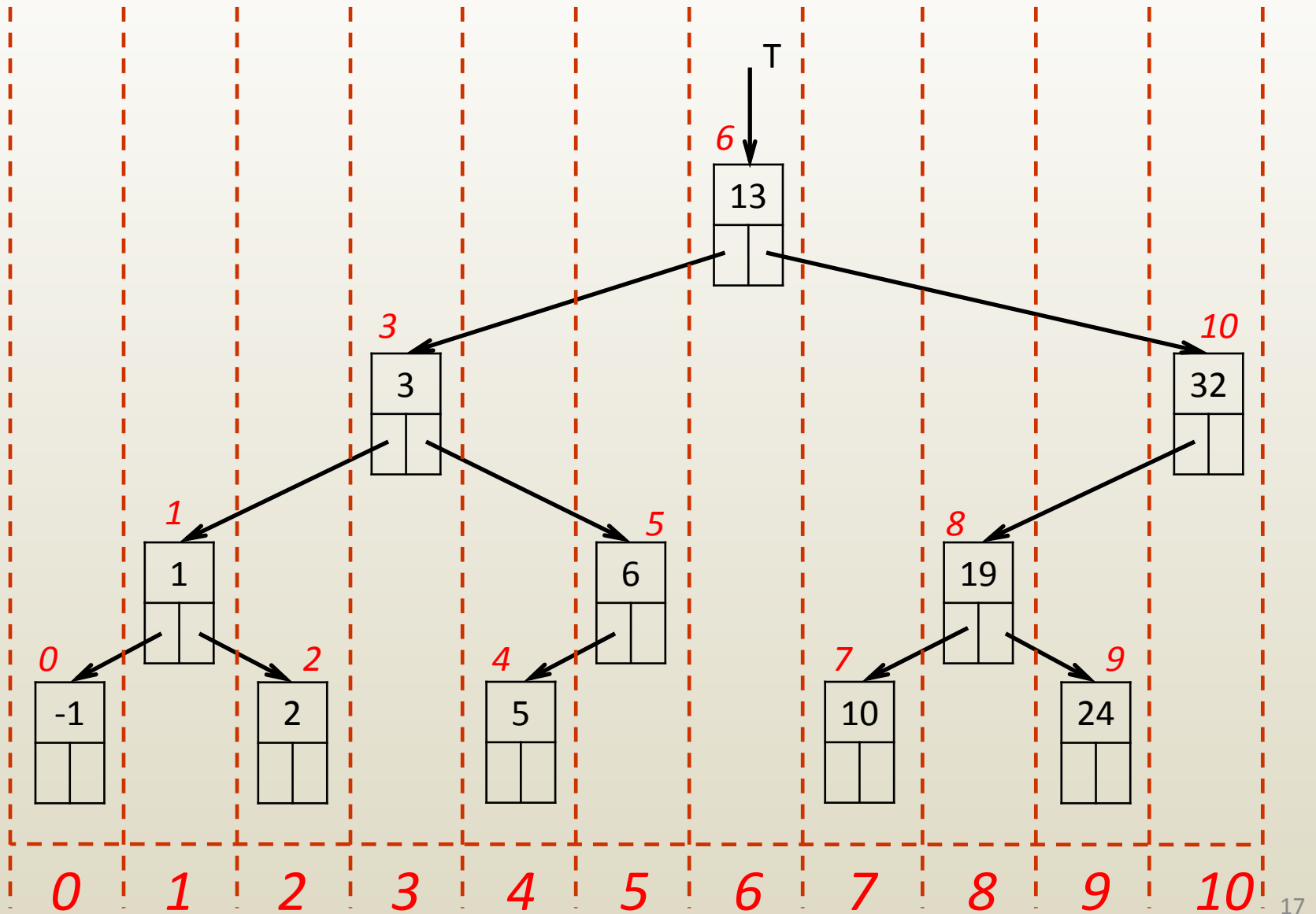
- Vertical coordinate “vCoord”
 - Equals depth in tree
 - Same for all nodes on same level



Setting horizontal node coordinates

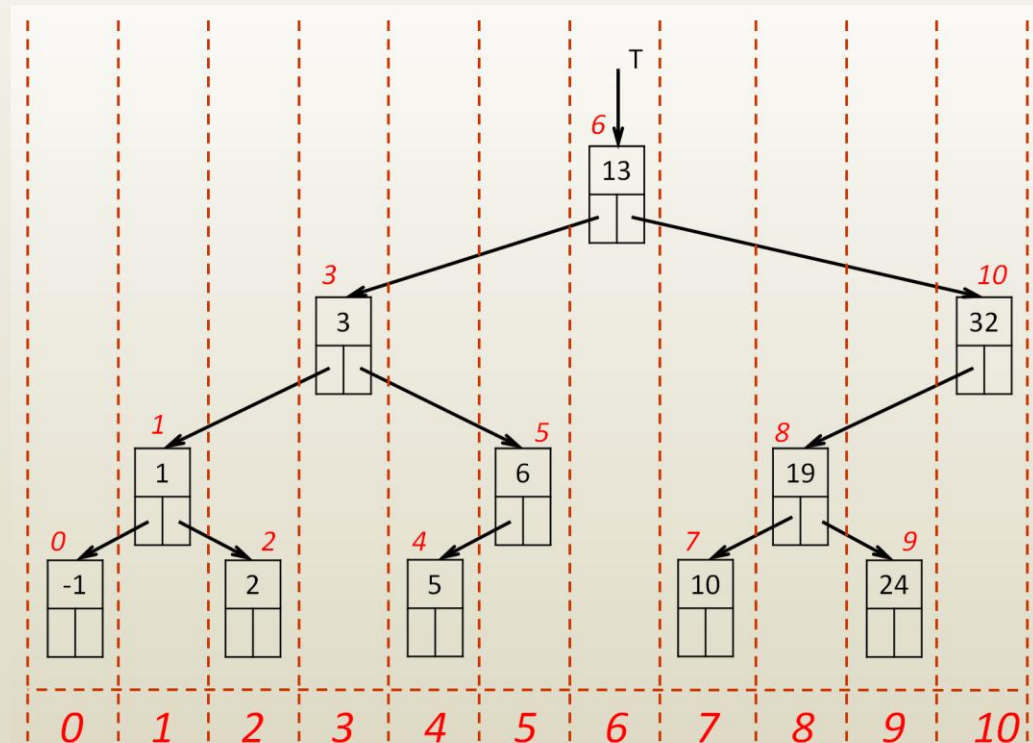
- Each node has a different coordinate
- Each node belongs to a different column
- If there are n nodes, we need n columns
- A node has a greater coordinate than any node in its left sub-tree
- Left-most node has coordinate 0

Setting horizontal node coordinates

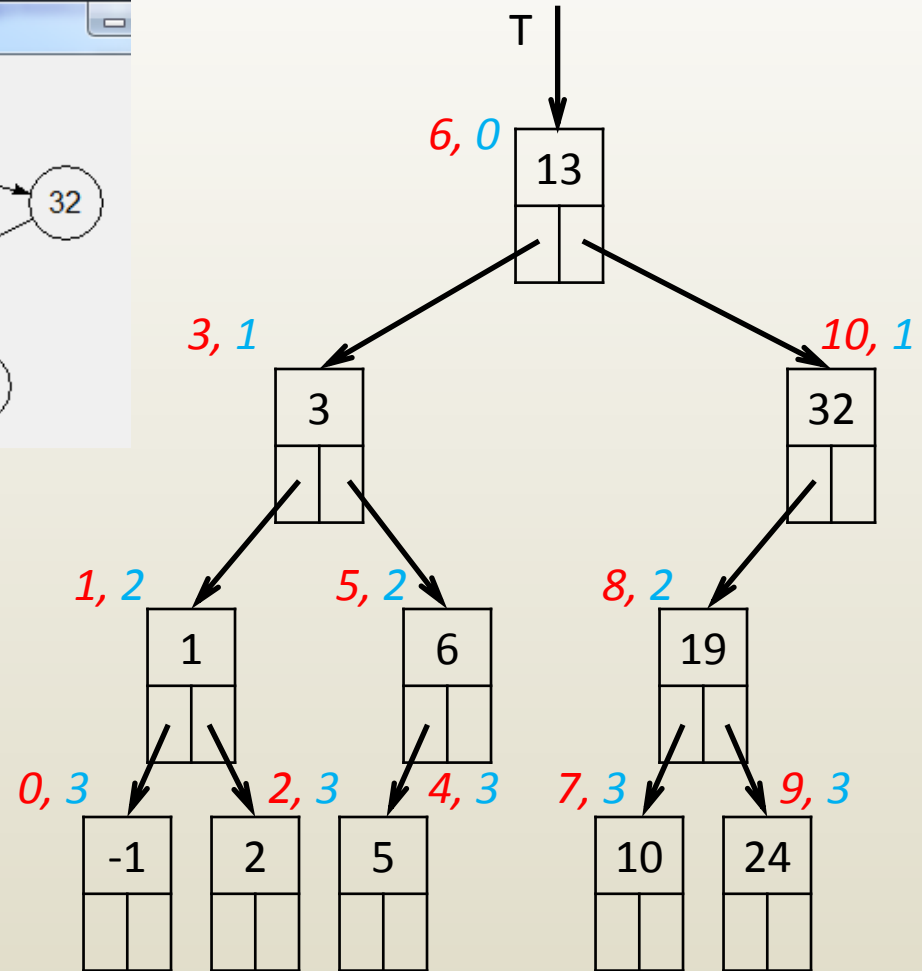
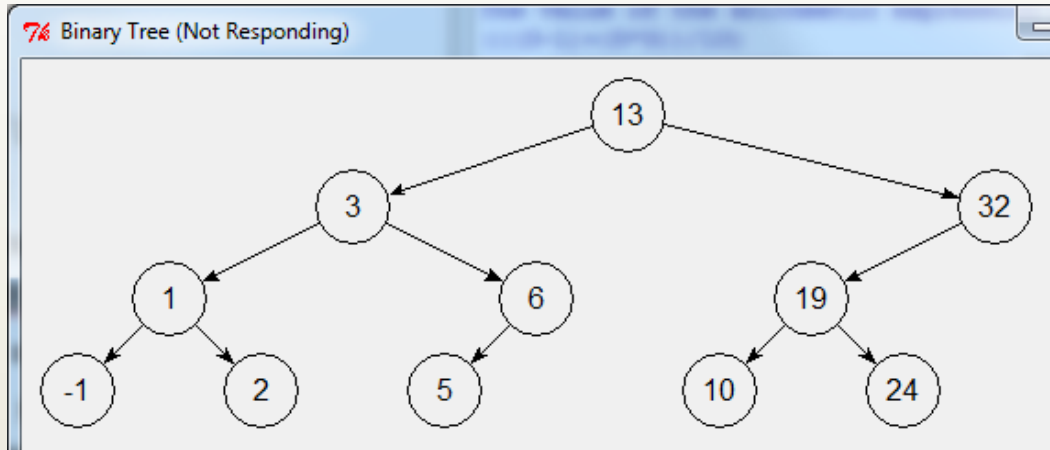


Setting horizontal node coordinates

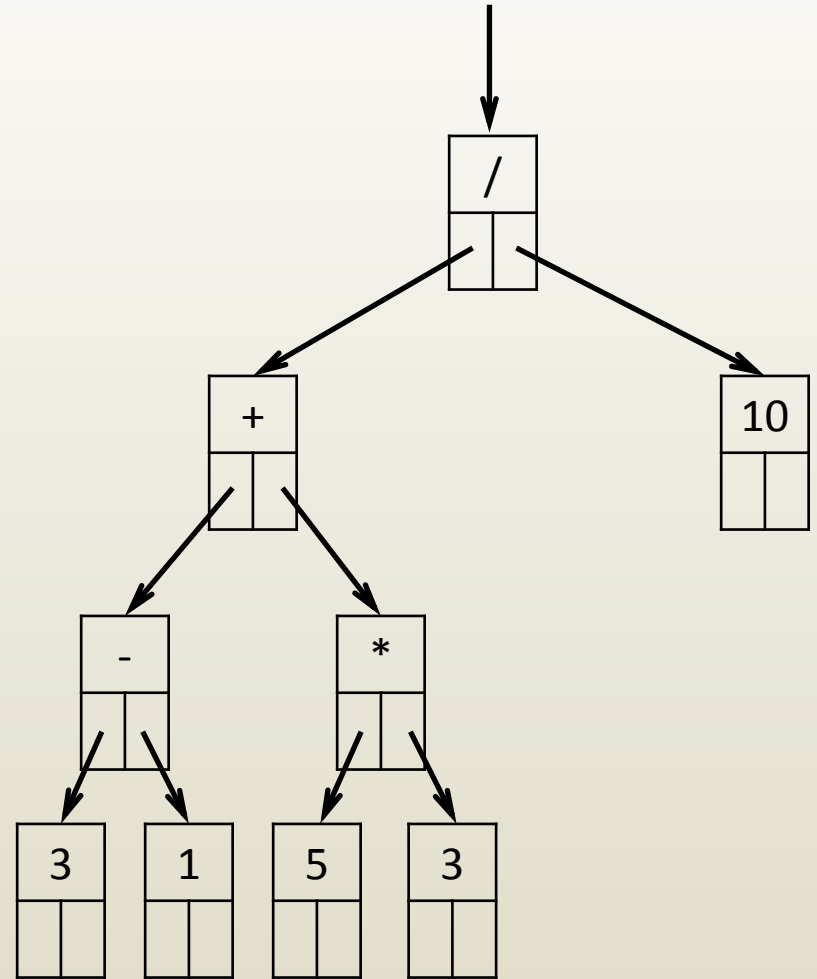
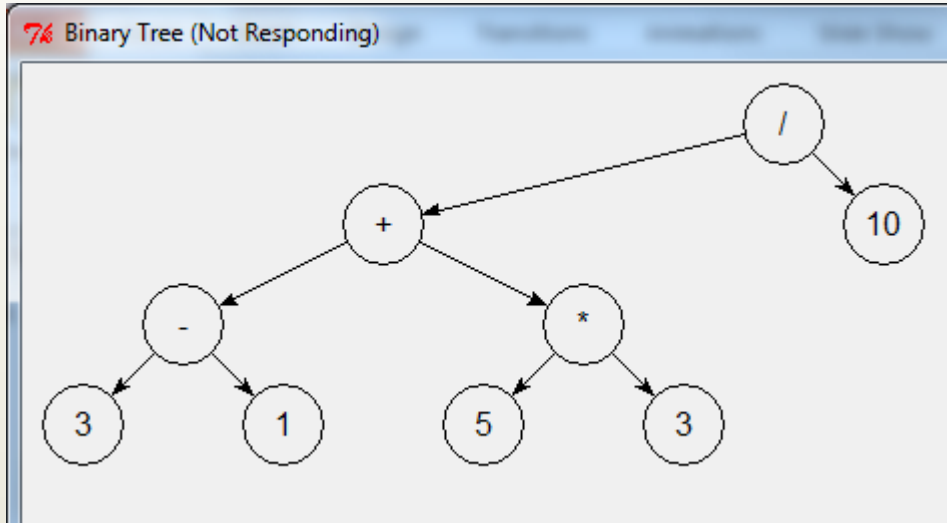
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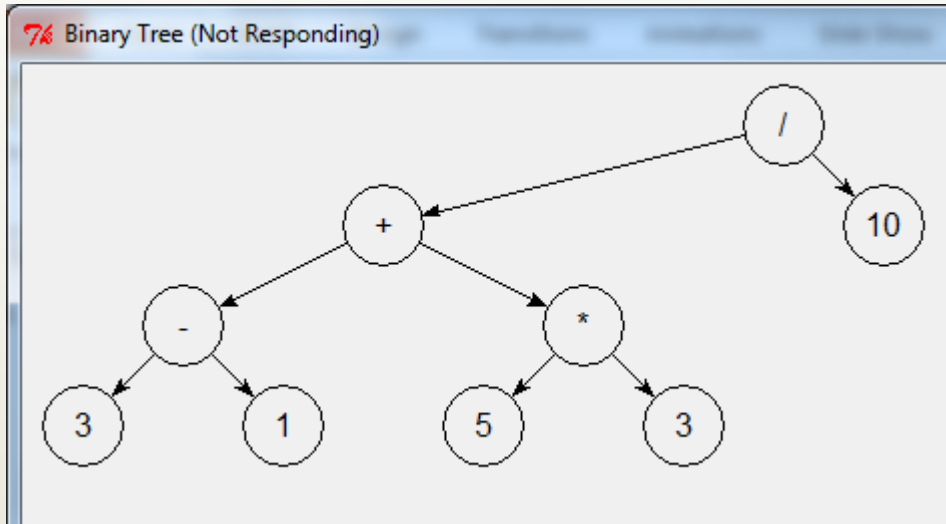
Example 1



Example 2



iClicker question



Our tree drawing algorithm doesn't place the root half way btw. its children

- A. The implementation is not correct.
- B. The implementation is correct.
- C. We could enhance the algorithm to avoid this.
- D. B and C
- E. A and C

Computing factorial

Input:

n // factorial function
argument

Output:

// n!

Factorial(n)

```
if n == 1
    return 1
endif
return n * Factorial(n-1)
```

endMinimum

def Fact(n):

```
if n == 1:
    return 1
return n * Fact(n-1)
```

Finding minimum in array

Input:

A // array of integers
n // number of elements in array (array size)
 i_0 // consider elements from i_0 onwards

Output:

Min // value of element with smallest value

MinR(A, n, i_0) // recursive version

```
if  $i_0 == n-1$  // last element
    return A[ $i_0$ ]
endif
tmp = MinR(A, n,  $i_0+1$ )
if A[ $i_0$ ] < tmp
    return A[ $i_0$ ]
else
    return tmp
endMinR
```

def MinR(A, i_0):

```
if  $i_0 == \text{len}(A)-1$ :
    return A[ $i_0$ ]
tmp = MinR(A,  $i_0+1$ )
if A[ $i_0$ ] < tmp:
    return A[ $i_0$ ]
else:
    return tmp
```

Binary search

- Finds whether a number appears in a sorted array in logarithmic time
 - $\log n$, where n is the number of elements in the array

```
def BinSearchR(A, l, r, a):  
    if l == r:  
        return A[l] == a  
    m = (int) ((l+r)/2)  
    if A[m] >= a:  
        return BinSearchR(A, l, m, a)  
    else:  
        return BinSearchR(A, m+1, r, a)
```