Irregular Data Structures

Linked lists, trees, and graphs

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Motivation

- Irregular data structures needed to overcome disadvantages of arrays
 - Easy expansion and contraction to keep up with dynamic data size
 - Modeling of irregular data, with complex "neighboring" relationship

Cost

- Irregular data structures
 - Increased complexity
 - Decreased efficiency
 - Structure stored explicitly, not all storage used to store data
 - No direct access to all data

- A 1-D sequence data structure
- Not an array
- Each data element is linked to the next
- Link: memory address pointing to a data element
- Link list node: data element
 + link

- Example
 - credit card transaction amounts in dollars, sorted
 - Links stored explicitly
 - E.g. 32 bit / link
 - Actual address irrelevant here
 - Link shown with arrow
 - Link to first element has to be known (shown in red)
 - Link of last element is null



- Add a new transaction in the amount of \$8.12
 - Start at first node (using known red arrow link)
 - Is amount (13.40) smaller than \$8.12?
 - No, use node link to go to next node
 - Is amount (12.50) smaller than \$8.12?
 - No, use node link to go to next node
 - Is amount (7.45) smaller than \$8.12?



- Add a new transaction in the amount of \$8.12
 - Yes, insert new node
 - Make new node
 - Set new node amount to \$8.12



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 - Yes, insert new node
 - Make new node
 - Set new node amount to \$8.12
 - Set new node link to next node



- Add a new transaction in the amount of \$8.12
 - Yes, insert new node
 - Make new node
 - Set new node amount to \$8.12
 - Set new node link to next node
 - Set previous node link to new node



• Add a new transaction in the amount of \$8.12



• Delete transaction \$12.50

– Move to node storing \$12.50 transaction



- Delete transaction \$12.50
 - Move to node storing \$12.50 transaction
 - Set link of previous node to next node



- Delete transaction \$12.50
 - Move to node storing \$12.50 transaction
 - Set link of previous node to next node
 - Delete current node



• Delete transaction \$12.50



- Delete transaction \$13.40
 - Special case
 - Set red link equal to link of first node
 - Delete first node



- Delete transaction \$13.40
 - Special case
 - Set red link equal to link of first node
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- Advantages
 - List grows and shrinks as needed, w/o having to modify entire list
 - Insertion & deletion imply local changes
- Disadvantages
 - You cannot find third transaction directly
 - Have to traverse list
 - Storing link implies overhead



iClicker question

When inserting a transaction with value 1.00 in the linked list below, which of the following statements is true:

- A. The transaction cannot be inserted since there is no transaction of smaller value.
- B. The next node link of the new node will be NULL.
- C. The insertion point is found when the next node link of the current node is found to be NULL.
- D. A, B, and C are true.
- E. B and C are true.



Binary tree

- Definition
 - A hierarchical data structure
 - A (parent) node links to 0,1, or 2 (children) nodes
 - The starting node is called root; the root is not the child of any node
 - Nodes with 0 children are called leafs
 - Non-leaf nodes are called internal



- Operators at internal nodes
- Operands at leafs



- Operators at internal nodes
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- Arithmetic expression can be recovered by traversing tree
- Traversal: visiting all nodes
- Traversal rules
 - Start at root
 - For every node
 - go left until dead end
 - then go right until dead end
 - then go back up
- Printout rules
 - Write "(" before going left
 - Write current node symbol before going right
 - Write ")" after having gone right



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iClicker question

- Which traversal called COUNT counts the number of leafs in a binary tree.
- A. If leaf, return 1. If not leaf, return COUNT(left child) + COUNT(right child)
- B. If leaf, return 1. If not leaf return0. COUNT(left child).COUNT(right child).
- C. If leaf, return operand. If not leaf, return COUNT(left child) + COUNT(right child)
- D. None of the above.
- E. All of the above.

Graphs

- Graphs
 - Nodes (also called vertices) connected by links, called edges
 - Nodes have a variable number of incident edges
 - Great flexibility
- Example: airline routes

- "List of edges"
 - Array of nodes & array of edges
 - Edges pair of node indices
 - Origin node first

2

JFK

0

ORD

1

IND

Destination node second

3

RDU

4

DFW

5

ATL

6

• "List of edges"

2

JFK

0

ORD

1

IND

- Used only for sparse graphs
 (i.e. a small number of edges)
- Difficult to find whether there is an edge between two nodes (requires traversal of edge list)

3

RDU

4

DFW

5

ATL

6

MIA

- "Adjacency lists"
 - One array for each node
 - Array stores adjacent nodes

Adjacency list

for node 6

0

3

1

5

49

ORD

0,1

2,0

1,0

IND

JFK

- "Adjacency lists"
 - Finding an edge only requires traversing the starting node's adjacency list

- "Adjacency matrix"
 - A 2-D matrix
 - Row corresponds to start node
 - Column corresponds to end node
 - 0 if no edge, 1 if edge

	0	1	2	3	4	5	6	
0	0	1	0	0	1	0	0	
1	1	0	0	0	1	1	0	
2	1	0	0	0	0	0	0	
3	0	0	1	0	0	1	0	
4	0	1	0	0	0	1	0	
5	0	0	1	1	1	0	1	
6	0	0	0	1	0	1	0	

- "Adjacency matrix"
 - An edge is found in constant time
 - Is there an edge between ATL and ORD?
 - A[5][0] is 0 so the answer is no
 - Storage quadratic in number of nodes
 - Inefficient for sparse graphs

	0	1	2	3	4	5	6	
0	0	1	0	0	1	0	0	
1	1	0	0	0	1	1	0	
2	1	0	0	0	0	0	0	
3	0	0	1	0	0	1	0	
4	0	1	0	0	0	1	0	
5	0	0	1	1	1	0	1	
6	0	0	0	1	0	1	0	

Directed graphs

- So far we talked about directed graphs
 - an edge started from one node and ended at another
 - the edge could only be traversed in one direction

Undirected graphs

• In an undirected graph edges can be traversed in either direction.

