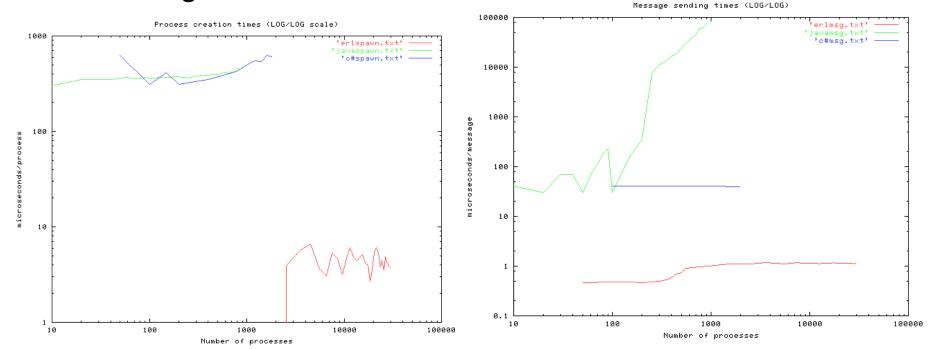
Principles of Concurrency and Parallelism

Lecture 5: Erlang

CS390C: Principles of Concurrency and Parallelism

Example

- Suppose we have N threads (or processes, tasks) that form a ring
- Each thread communicates with its neighbor forwarding a message
- How big can we make the ring?
- How long does it take to send a message?



CS390C: Principles of Concurrency and Parallelism

Philosophy

- Independent processes
 - suitable for executing on distributed machines
- No sharing
 - (Deep) copy data sent on messages
 - no cross-machine pointers
 - no locks, data races, synchronization issues, ...
- All processes have a unique name
- Asynchronous sends, synchronous receives
 - Eventual delivery
 - But if A sends messages m1 and m2 to B, m2 will never arrive before m1

Key features

- Functional
 - single assignment (every variable assigned to at most once)
- Lightweight first-class processes
- Pattern-matching
- Small collection of datatypes
 - lists, tuples, pairs
- Dynamic typing
- Realtime concurrent garbage collection

Examples

```
-module(math).
-export([fac/1]).
fac(N) when N > 0 \rightarrow N * fac(N-1);
            -> 1.
fac(0)
> math:fac(25).
15511210043330985984000000
lookup(Key, {Key, Val, _, _}) ->
  {ok, Val};
lookup(Key, {Key1, Val, S, B}) when Key<Key1 ->
  lookup(Key, S);
lookup(Key, {Key1,Val,S,B}) ->
  lookup(Key, B);
lookup(Key, nil) ->
  not found.
```

Examples

```
append([H|T], L) -> [H|append(T, L)];
append([], L) -> L.
sort([Pivot|T]) ->
  sort([X|X < T, X < Pivot]) ++
  [Pivot] ++
  sort([X | X < T, X > = Pivot]);
sort([]) -> [].
> Adder = fun(N) -> fun(X) -> X + N end end.
#Fun
> G = Adder(10).
#Fun
> G(5).
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```

Concurrency

Concurrency

```
-module(counter).
-export([start/0,loop/1]).

start() ->
   spawn(counter, loop, [0]).

loop(Val) ->
   receive
        increment -> loop(Val + 1)
end.
```

Issues:

- Cannot directly access counter value.
- Messaging protocol is explicit (via message increment)

Refinement

```
-module(counter).
-export([start/0,loop/1,increment/1,value/1,stop/1]).
%% First the interface functions.
start() ->
  spawn(counter, loop, [0]).
increment(Counter) -> Counter ! increment.
value(Counter) ->
  Counter ! {self(),value},
  receive
     {Counter, Value} -> Value
end.
stop(Counter) -> Counter ! stop.
loop(Val) ->
    receive
        increment -> loop(Val + 1);
        {From, value} -> From ! {self(), Val}, loop(Val);
        stop -> true;
        Other -> loop(Val)
    end.
```

Concurrency

```
-module(M).
-export([start/0, ping/1, pong/0]).
ping(0) \rightarrow
   pong! finished,
   io:format("ping finished~n", []);
ping(N) \rightarrow
   pong ! {ping, self()},
   receive pong ->
      io:format("Ping received pong~n", [])
   end,
ping(N - 1).
pong() \rightarrow
    receive
      finished -> io:format("Pong finished~n", []);
      {ping, Ping PID} ->
         io:format("Pong received ping~n", []),
         Ping PID ! pong,
         pong()
       end.
start() -> register(pong, spawn(M, pong, [])),
            spawn(M, ping, [3]).
```

Distributed Programming

Can generalize previous example to a distributed environment

```
-module(M).
-export([start/0, ping/1, pong/0]).
ping(0,Pong node) ->
                                                    On one host: erl -sname ping
   {pong, Pong Node} ! finished,
   io:format("ping finished~n", []);
                                                    On another: erl -sname pong
ping(N) \rightarrow
   {pong, Pong Node} ! {ping, self()},
                                                    On one node:
  receive pong ->
      io:format("Ping received pong~n", [])
                                                       M:start pong().
   end,
                                                    On another:
ping(N - 1, Pong Node).
                                                       M:start ping(pong@<host>).
ponq() \rightarrow
    receive
      finished -> io:format("Pong finished~n", []);
      {ping, Ping PID} ->
         io:format("Pong received ping~n", []),
         Ping PID ! pong,
         pong()
       end.
start pong() -> register(pong, spawn(M, pong, [])),
start ping(Pong Node) -> spawn(M, ping, [3, Pong Node]).
                                                                            11
```

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Monitoring

Client/Server

```
server(Fun, Data) ->
  receive
    {new fun, Fun1} ->
      server(Fun1, Data);
    {rpc, From, ReplyAs, Q} ->
      {Reply, Data1} = Fun(Q, Data),
      From ! {ReplyAs, Reply},
      server(Fun, Data1)
  end.
 rpc(A, B) \rightarrow
   Tag = new ref(),
   A ! {rpc, self(), Tag, B},
   receive
      {Tag, Val} -> Val
   end
```

Concurrency Patterns

```
RPC call
Unicast
                                                Call (RPC)
                                                  A ! {self(), B},
  A ! B
                                                  receive
                                                     {A, Reply} \rightarrow
                                                       Reply
                                                  end
Event Handling
                                             Callback
  receive A -> A end
                                                   receive
                                                     {From, A} ->
                                                       From ! F(A)
```

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end

Concurrency Patterns

Callback within RPC

```
A ! {Tag, X}, g(A, Tag).

g(A, Tag) ->
  receive
    {Tag, Val} -> Val;
        {A, X} ->
        A ! F(X),
        go(A, Tag)
  end.
```

Timeouts

Readings

Concurrent Programming in Erlang (Part 1)

http://www.erlang.org/download/erlang-book-partl.pdf

See also: www.erlang.org