

Networks: how Information theory met the space and time

Philippe Jacquet

INRIA

Ecole Polytechnique

France

Plan of the talk

- History of networking and telecommunication
- Physics, mathematics, computer science
- Internet Routing complexity
- Mobile ad hoc networking complexity
- Wireless networking: Shannon's law.
- Information and space-time
- Wireless capacity and space-time

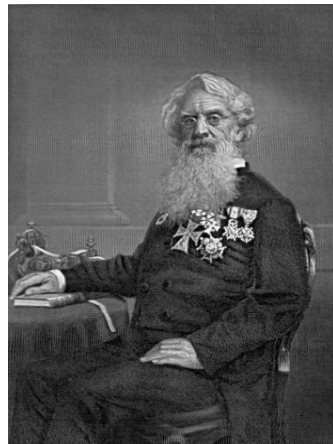
History of networking and telecommunication

- 1900: Marconi on Eiffel tower: first wireless telecommunication
- 1948: Shannon and Information Theory, transistor: first telecommunication massive optimization
- 1986: Birth of Internet, World Wide Web: First massive data, multimedia, multi-user network.
- 2000: the internet got wireless

Rise of Telecommunications (1800-1900)



C. Chappe 1793



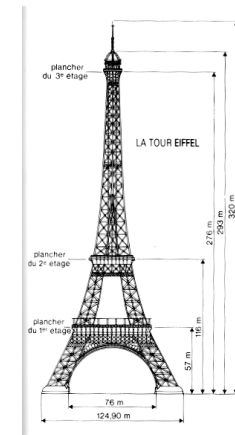
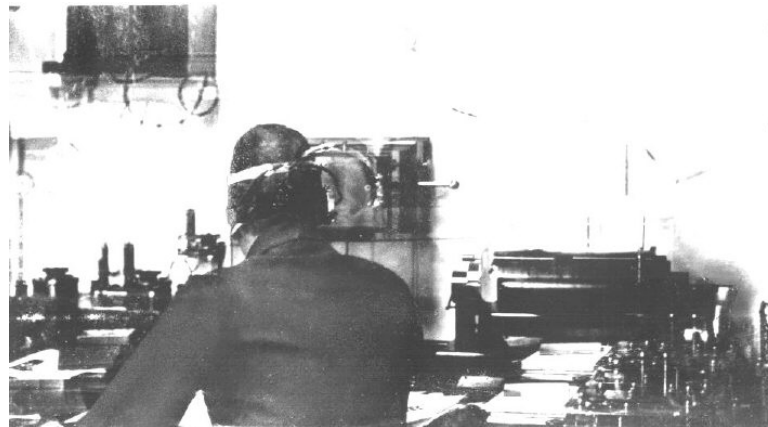
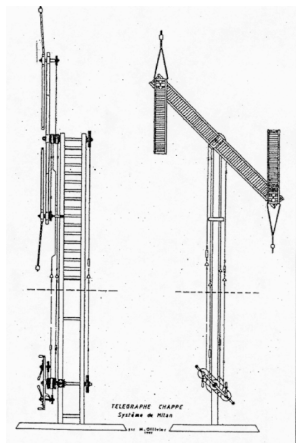
S. Morse 1837



T. Edison 1878

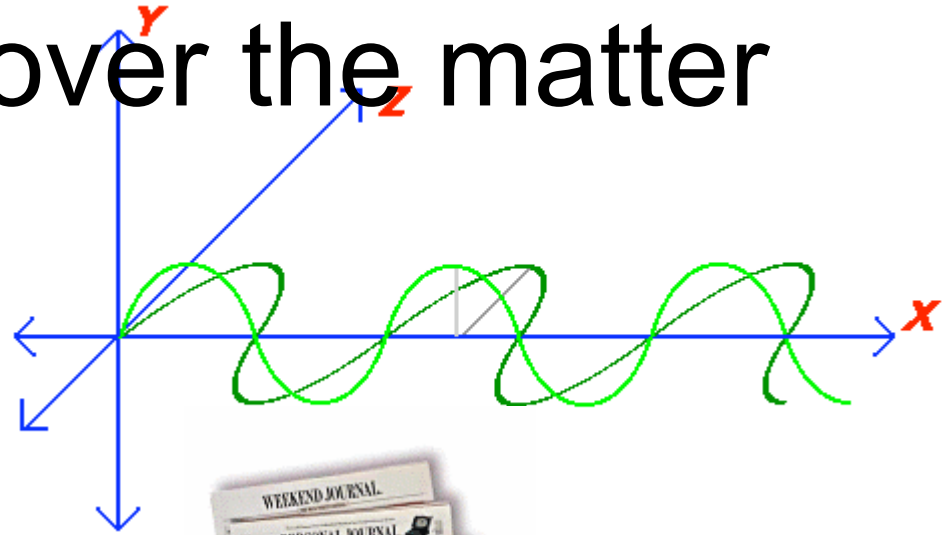


G. Marconi 1899



From pigeon to light speed: the triumph over the matter

- Electromagnetic
- Journal paper
- Pigeon post



From lamp to transistor: code and signal processing revolution.

The triumph over the numbers

- Telecommunications cross the border from physics to mathematics



enigma 1941

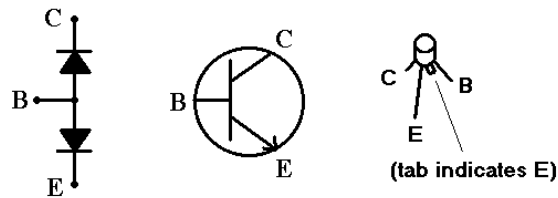
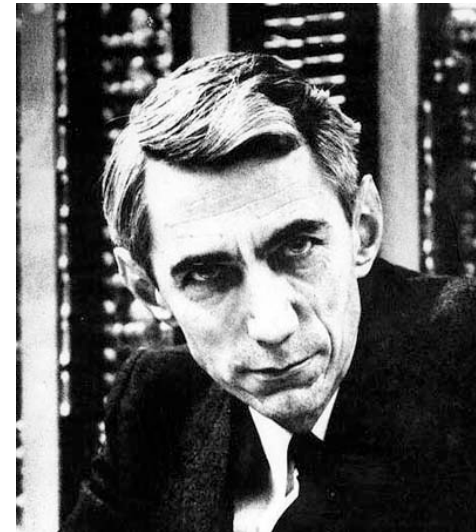


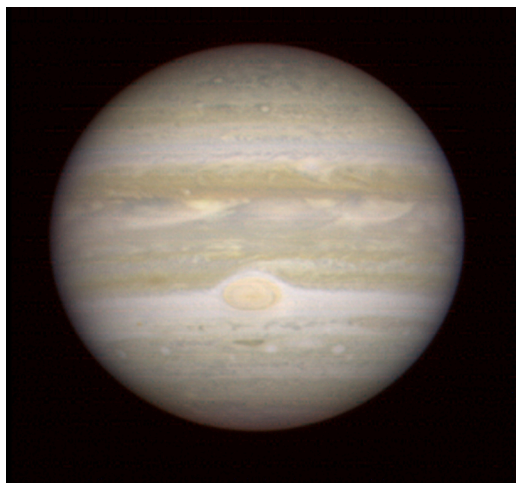
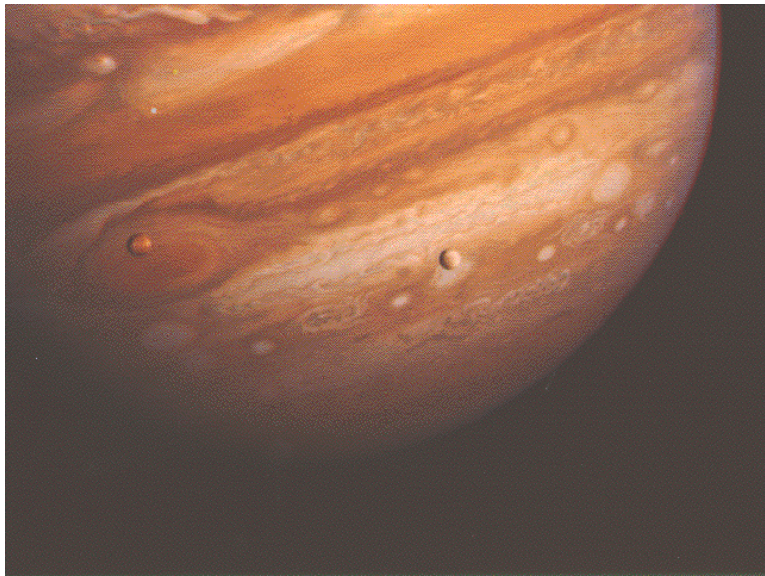
Fig. 1 The NPN Transistor

transistor 1947



C. Shannon 1948

The wonders of the digital revolution



C. Berrou 2000

Internet: The triumph over the complexity

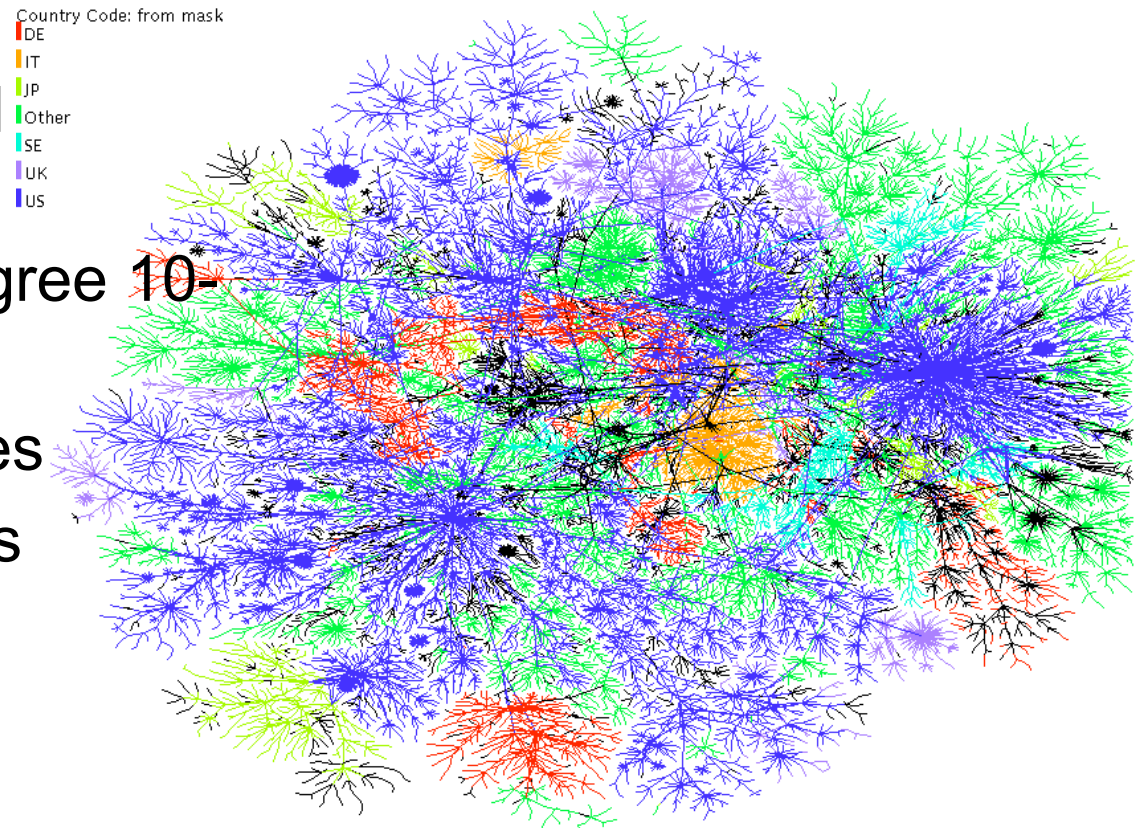
- Telecommunications cross the border of computer science
 - Cold war 1970: the strategic network identified as the weak point
- Answer: ARPANET then INTERNET
 - First, connect missiles sites
 - Then, universities
 - Every user
- The protocol (IP) binds heterogeneous networks (fiber, telephone, *etc*)
- Detects and avoids damaged parts in real time.



Facts and figure about internet

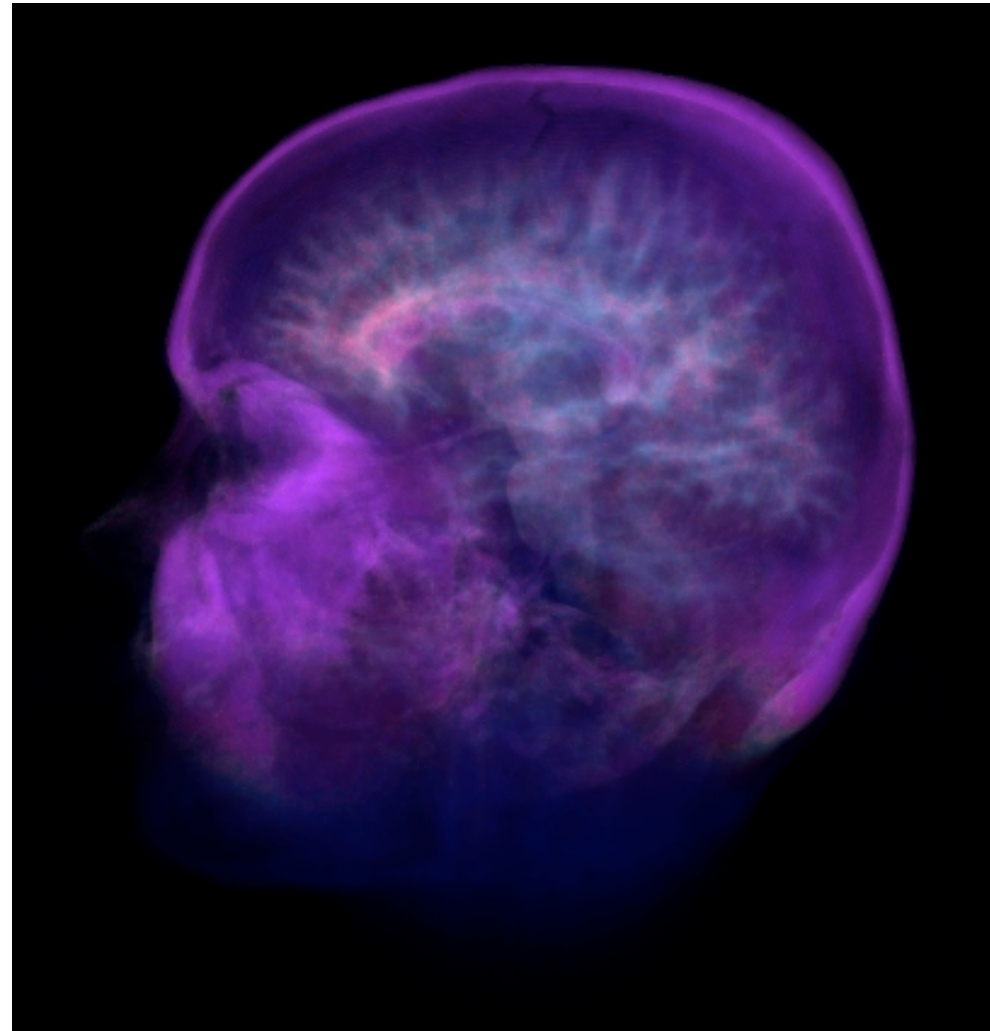
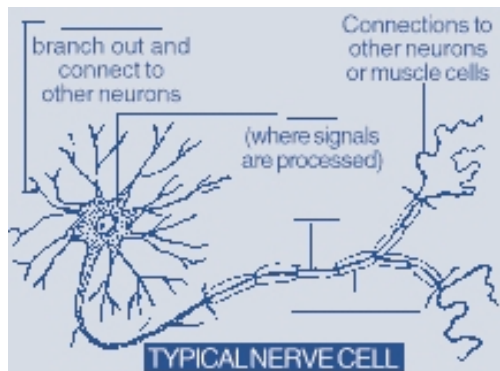
- internet:

- $6 \cdot 10^8$ connected machines
- Connectivity degree 10-100
- $3 \cdot 10^{10}$ web pages
- $2 \cdot 10^{15}$ online bits
- speed sub c



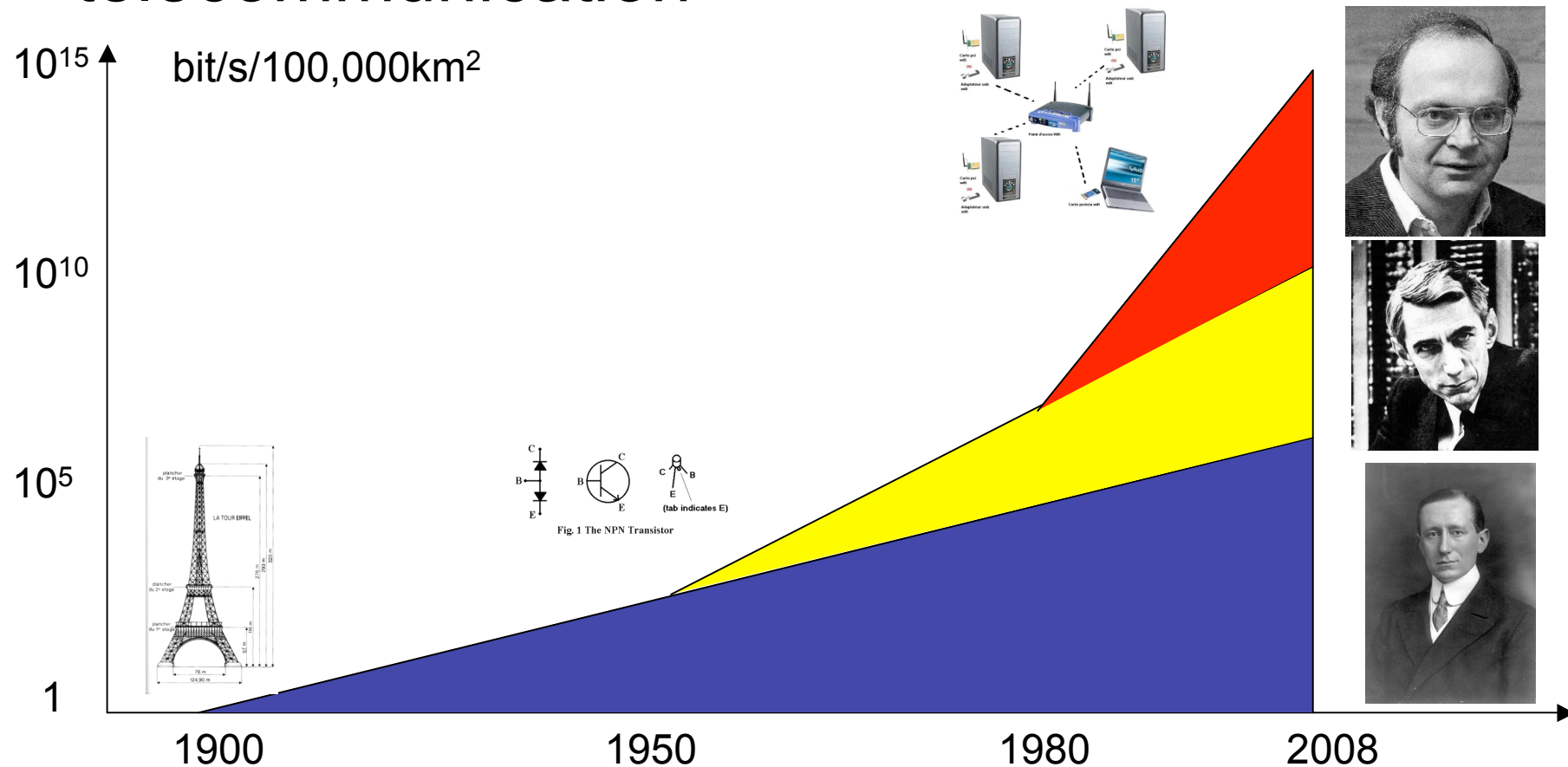
Internet is the most complicated artefact

- Far from human brain:
 - 10^{11} neurons
 - Connectivity degree : 10,000 -100,000
 - Speed: 100-400 m/s (scale with world internet)



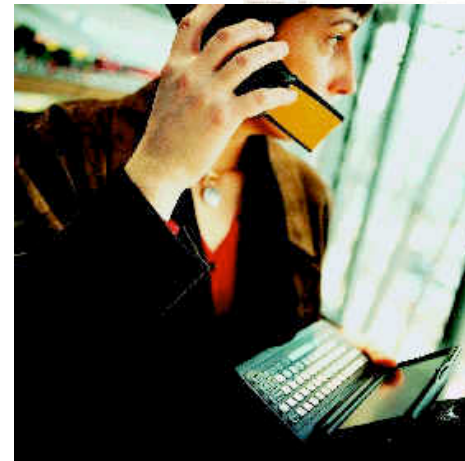
Histogram

- Physics, math and computer sciences in telecommunication



wireless performance from Marconi to Wifi

- Traffic density
 - 1900:
 - 10 bit/s/100.000 km² ,
 - 1000 watt
 - 2008:
 - 10,000,000 bit/s/ha,
 - 0.01 watt



A factor 10^{14}
100,000,000,
000,000

Comparison: road traffic



Road traffic increase 1900-2008: $< 10^5 = 100,000?$

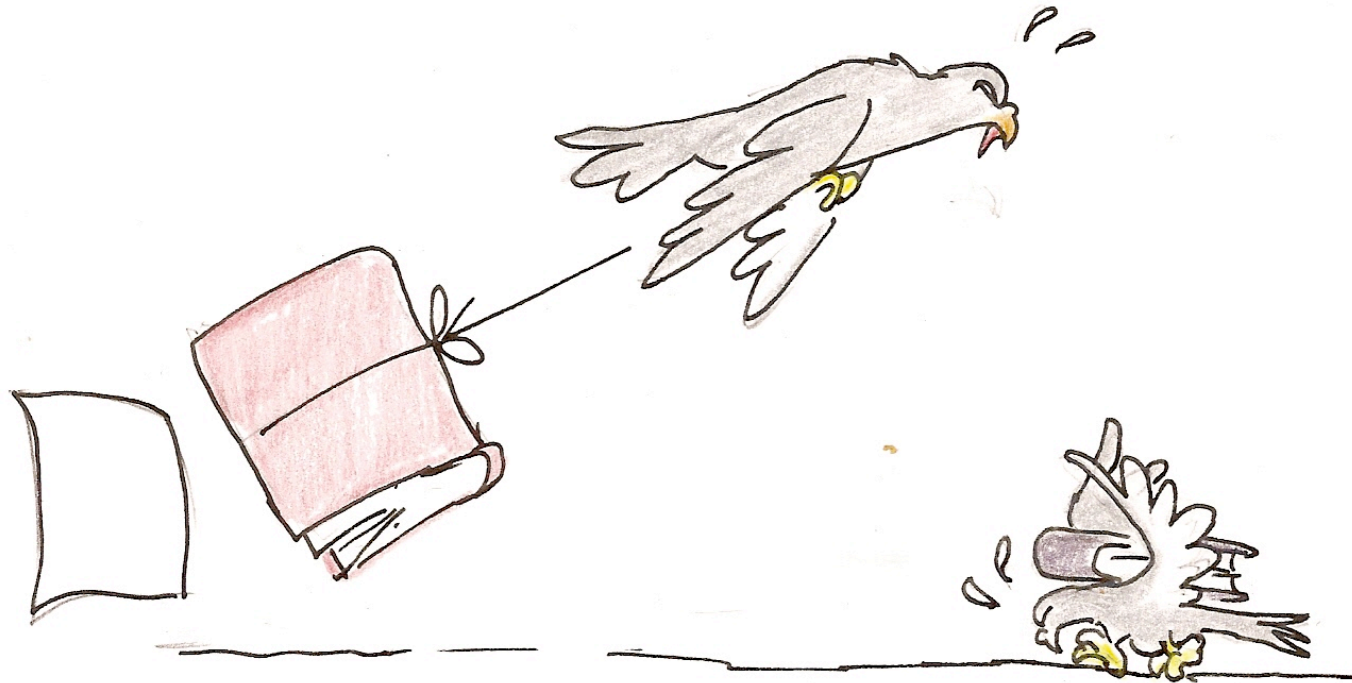
Physics, mathematics, computer science

- The telecommunications without
 - The physics



Physics, mathematics, computer science

- The telecommunications without
 - The mathematics



Physics, mathematics, computer science

- The telecommunications without
 - The computer science

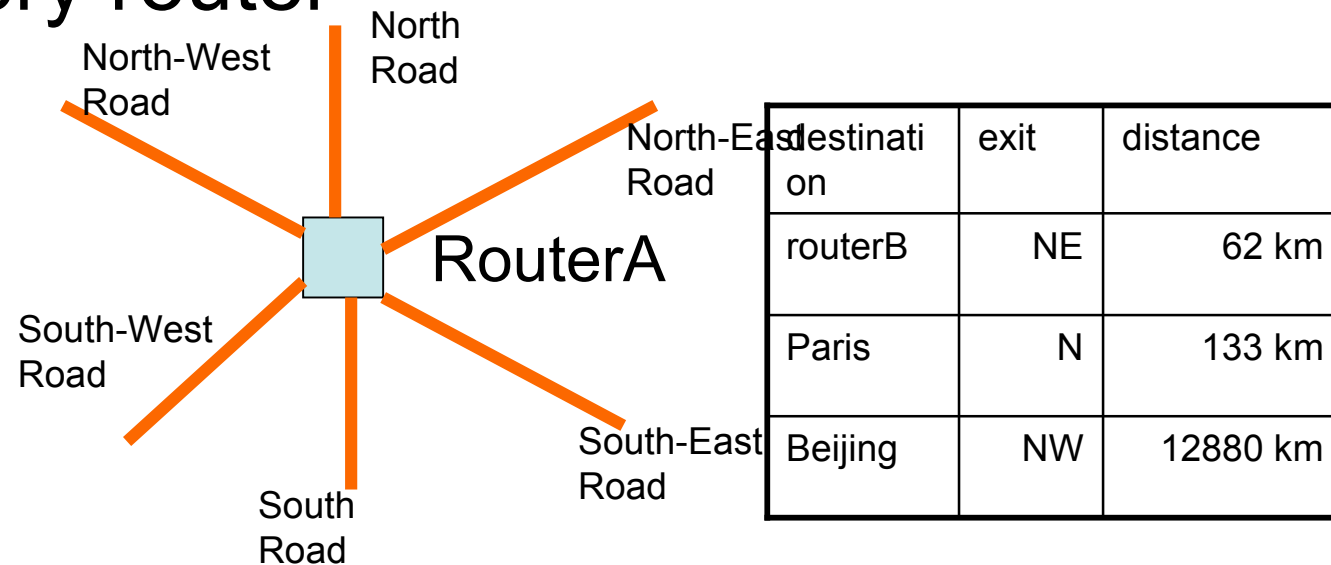


Physics-math-computer science

- Networks are together physics, math and computer science
- Computer science is the science which makes a complex system a simple object
 - Gee, how hard sometimes to reach this simplicity...

Example: routing protocol

- Routing tables is like orientation maps in every router

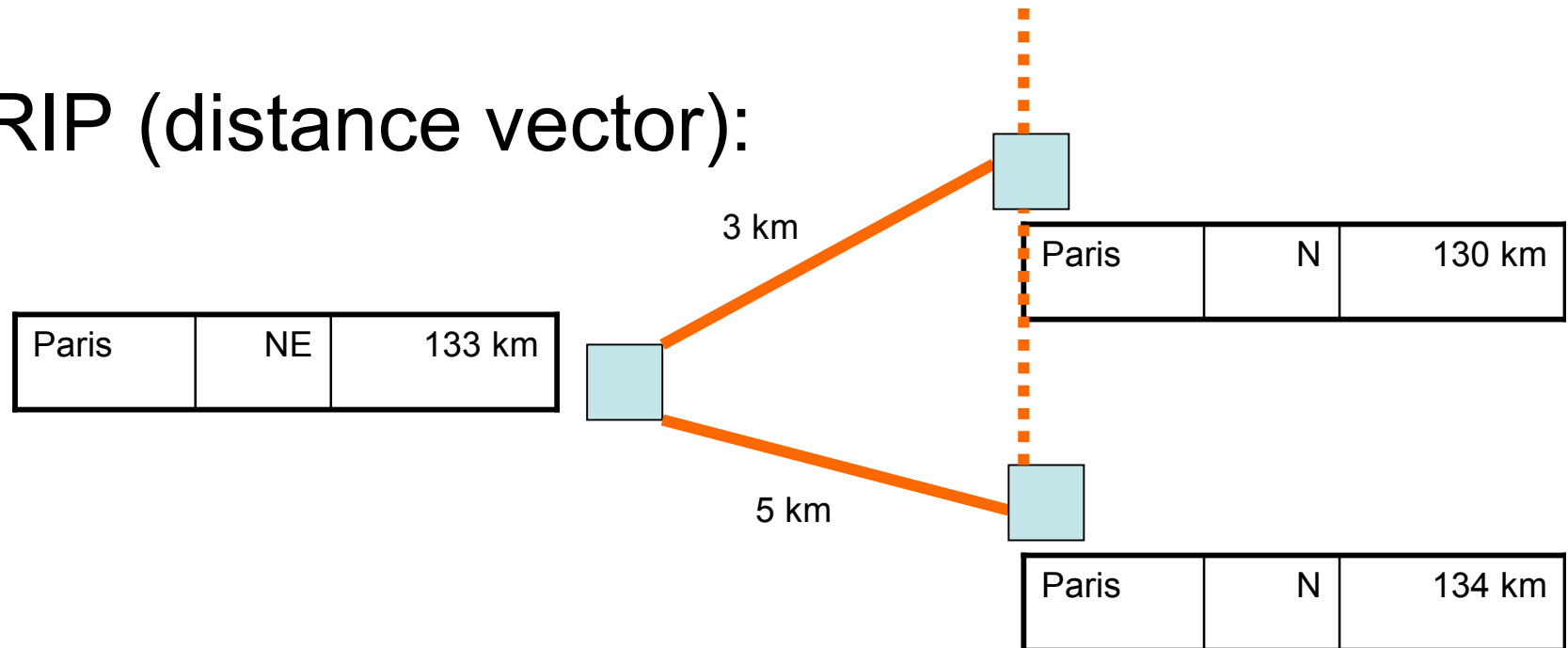


Example: routing protocol

- Two protocols;
 - RIP: run to neighbor protocol
 - Deliver the routing table to local neighbor
 - BGP: run the Tour de France protocol
 - Deliver the local routing table to whole network

Example: routing protocol

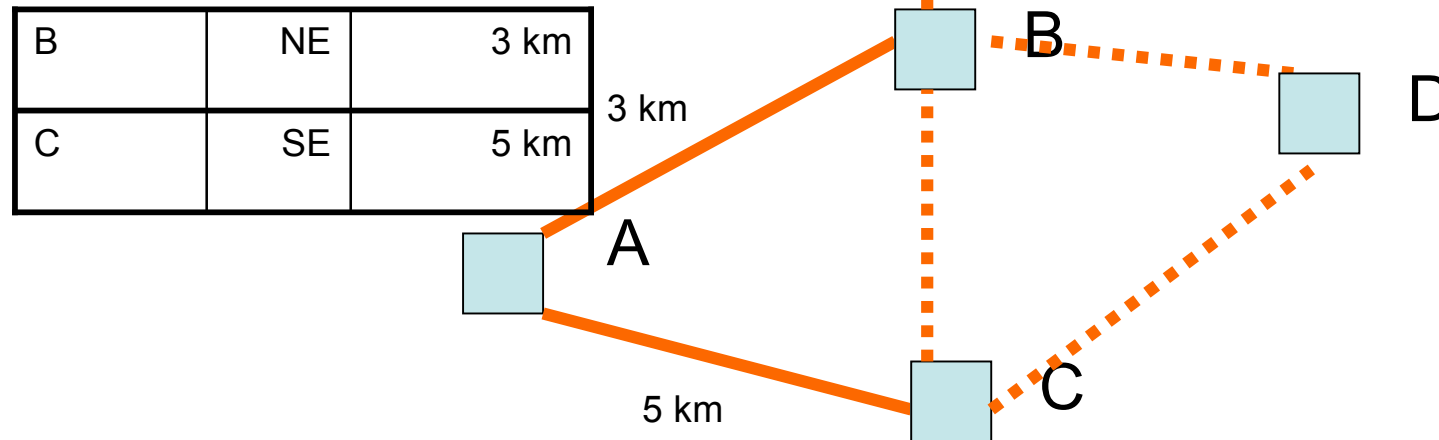
- RIP (distance vector):



- Complexity: NL (per refresh period)
- Convergence time: network diameter


Example: routing protocol

- BGP (link state):



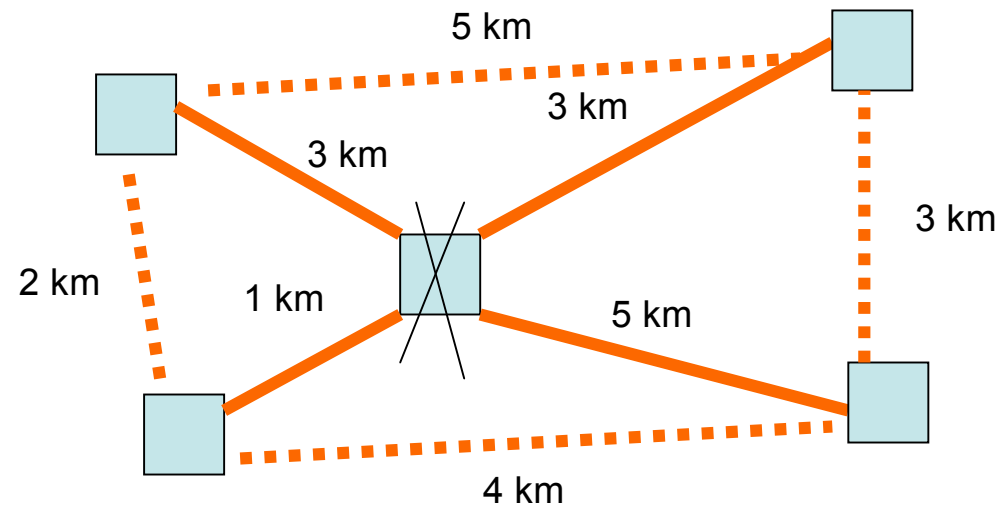
- Routing table computed on local link database
- Complexity L^2 (per refresh period)
- Convergence: network diameter

Example: routing protocol

- Which is best:
 - Deliver whole table to local?
 - Deliver local table to all? 
- Divergence time makes the difference!
 - Network Diameter with BGP (symmetric)
 - At least Diameter \times L with RIP (asymmetric)

RIP failure

- Count to infinity (1983 ARPANET incident)



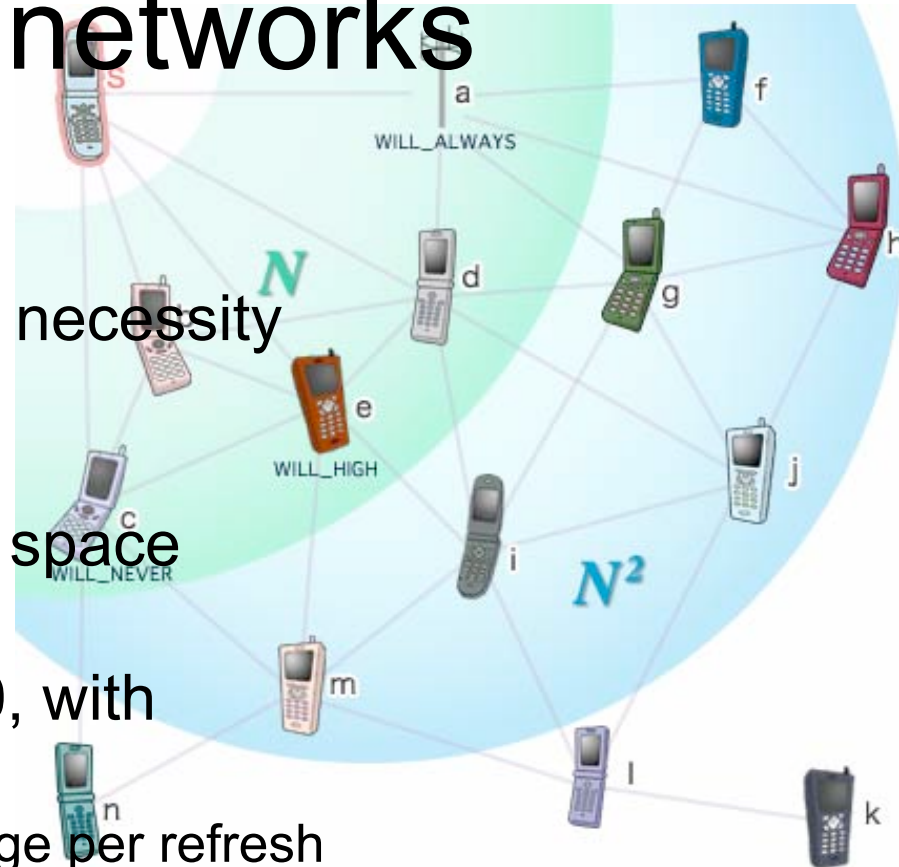
Diameter max limited to 15

$$\infty \times L < 15 \times L$$

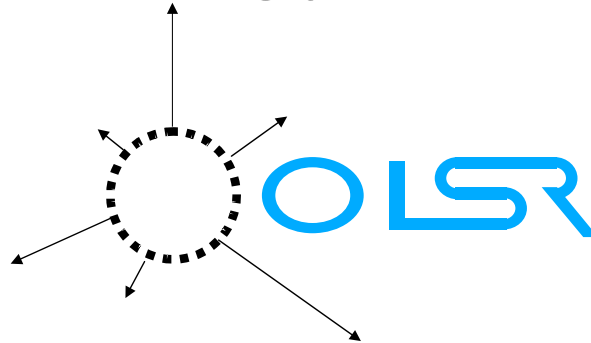
Wireless networks

Mobile *ad hoc* networks

- Mobility makes link failure a necessity
 - Refresh period 1 second
 - Automatic self-healing
- Local neighborhood is local space
 - Unlimited neighborhood size
- Stadium network: $N=10,000$, with average degree 1,000
 - BGP needs 10^{14} links exchange per refresh time
 - BGP fails on Wifi networks with 20 users at walking speed.
 - Heavy density kills link state management.



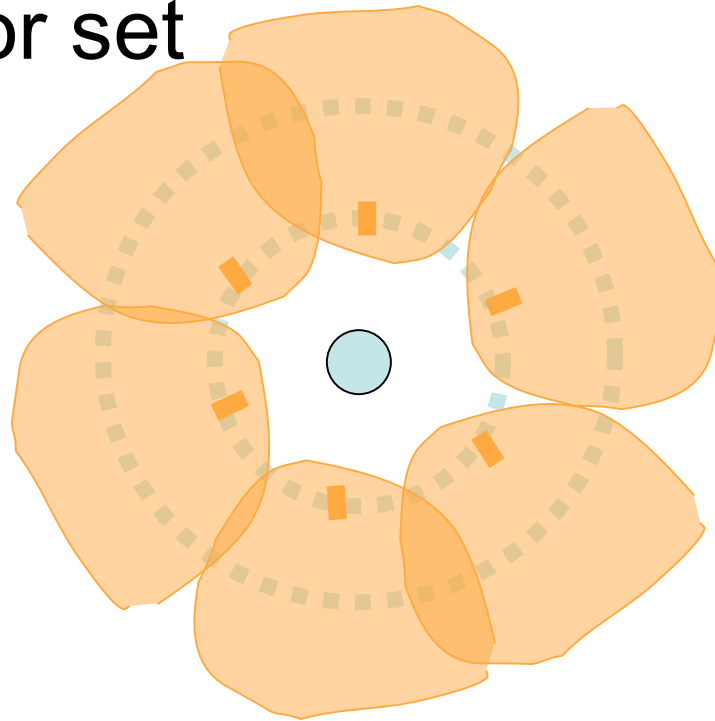
Wireless topology compression



- Optimized Link State Routing protocol
 - Advertise local table subset to whole network
 - Local neighborhood subset is the MultiPoint Relay (MPR) set of the node.
 - Every node receives a compressed topology information.

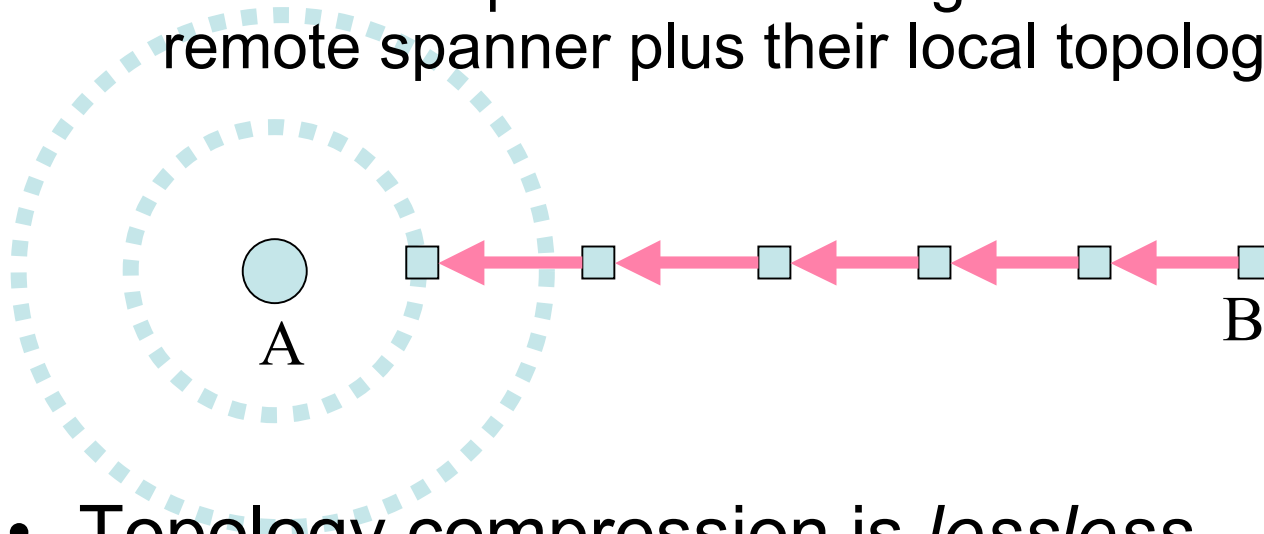
Wireless topology compression

- the MPR set covers the two-hop neighbor set



Wireless topology compression

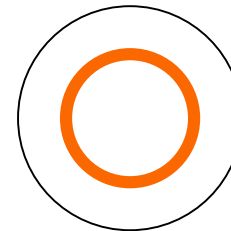
- MPR sets forms a *remote spanner*
 - Nodes compute their routing table on remote spanner plus their local topology.



- Topology compression is *lossless*
 - Optimal routes on remote spanner also optimal without compression.

Wireless topology compression

- In Erdos-Renyi random graphs
 - Random selection gives compression rate $\tau_r = \frac{N^3}{L^2} \log N$
- In unit disk graph model
 - Greedy selection gives compression rate $\tau_r = 3 \left(\frac{\pi N}{L} \right)^{\frac{2}{3}}$
- Stadium network
 - Compression rate 10^{-2}



Dissemination compression

- Only MPR retransmit local routing table
 - Another compression of rate $\approx \tau_r \frac{N}{L}$
- Stadium network
 - Overall compression 10^{-7}

Wireless protocol compression

- 10^7 ratio is like
 - your car traveling at light speed



Information theory in networks

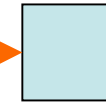
- Originally for point to point coding

A



X: Emitted code

B



Y: Received code

Channel capacity: $I(X, Y) = \text{entropy of } Y - \text{channel entropy}$

$$I(X, Y) = h(Y) - h(Y | X)$$

Example: wireless transmission

- Emitted Symbol: X an integer in $[1, S]$
- Received Symbol: $Y = X + \beta$
 - Noise β an integer in $[1, B]$

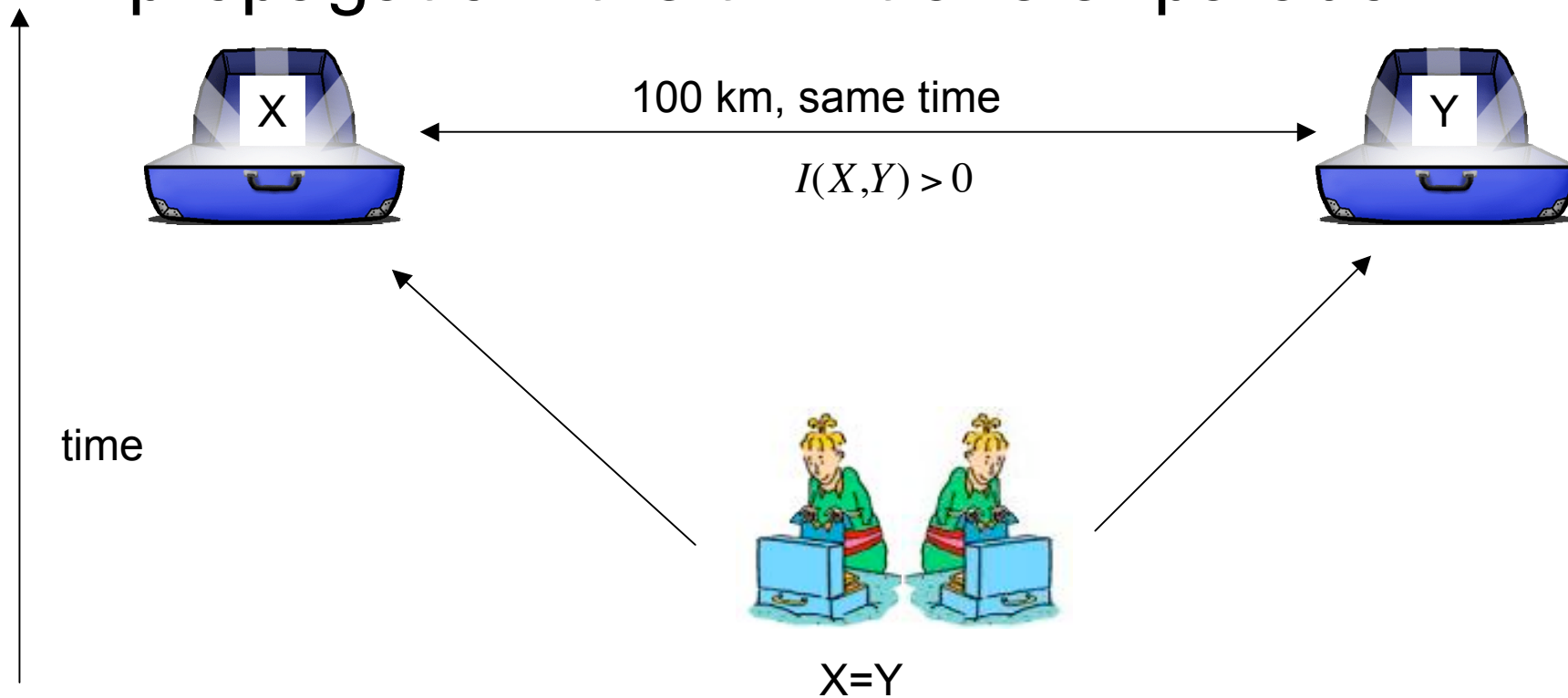
– Capacity per symbol $I(X, Y) = \log\left(1 + \frac{S}{B}\right)$

$$h(Y) = \log(S + B)$$

$$h(Y | X) = \log B$$

Limitation of information theory within space and time

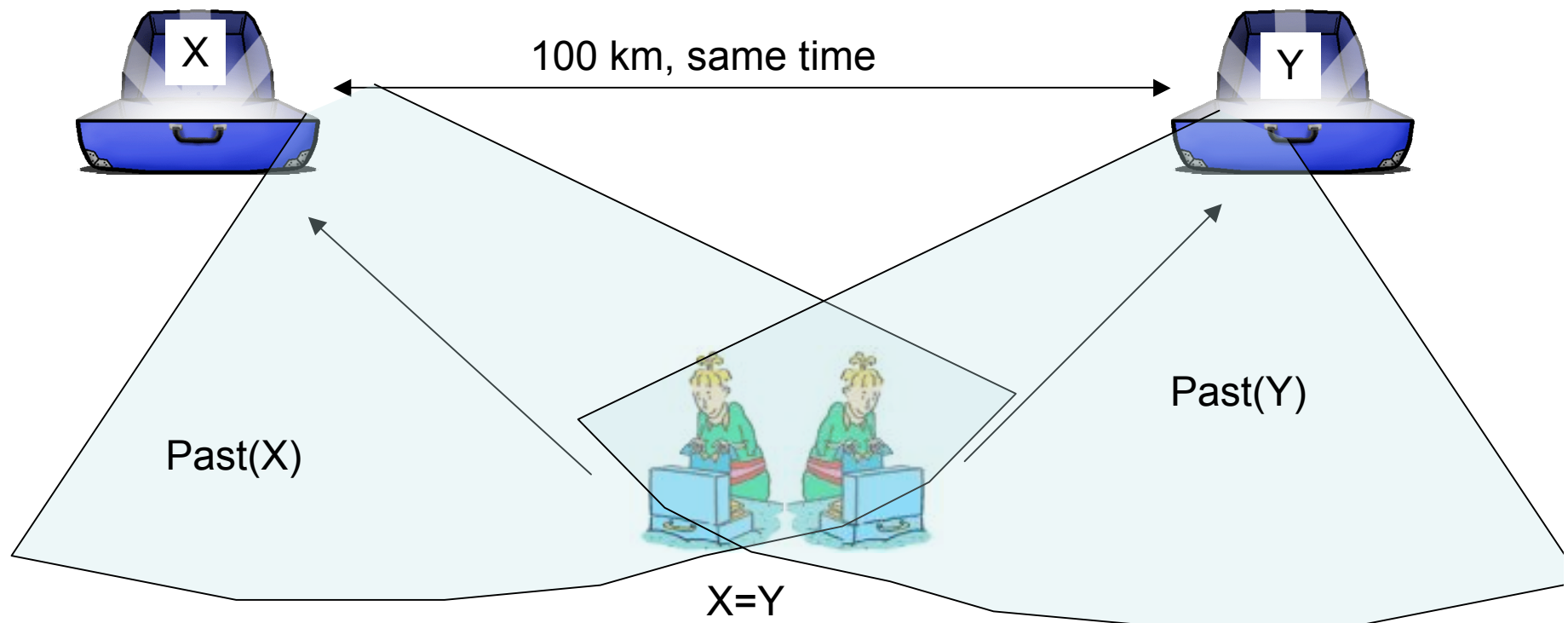
- Fake superluminal information propagation: the twin traveler paradox



Causality in information theory

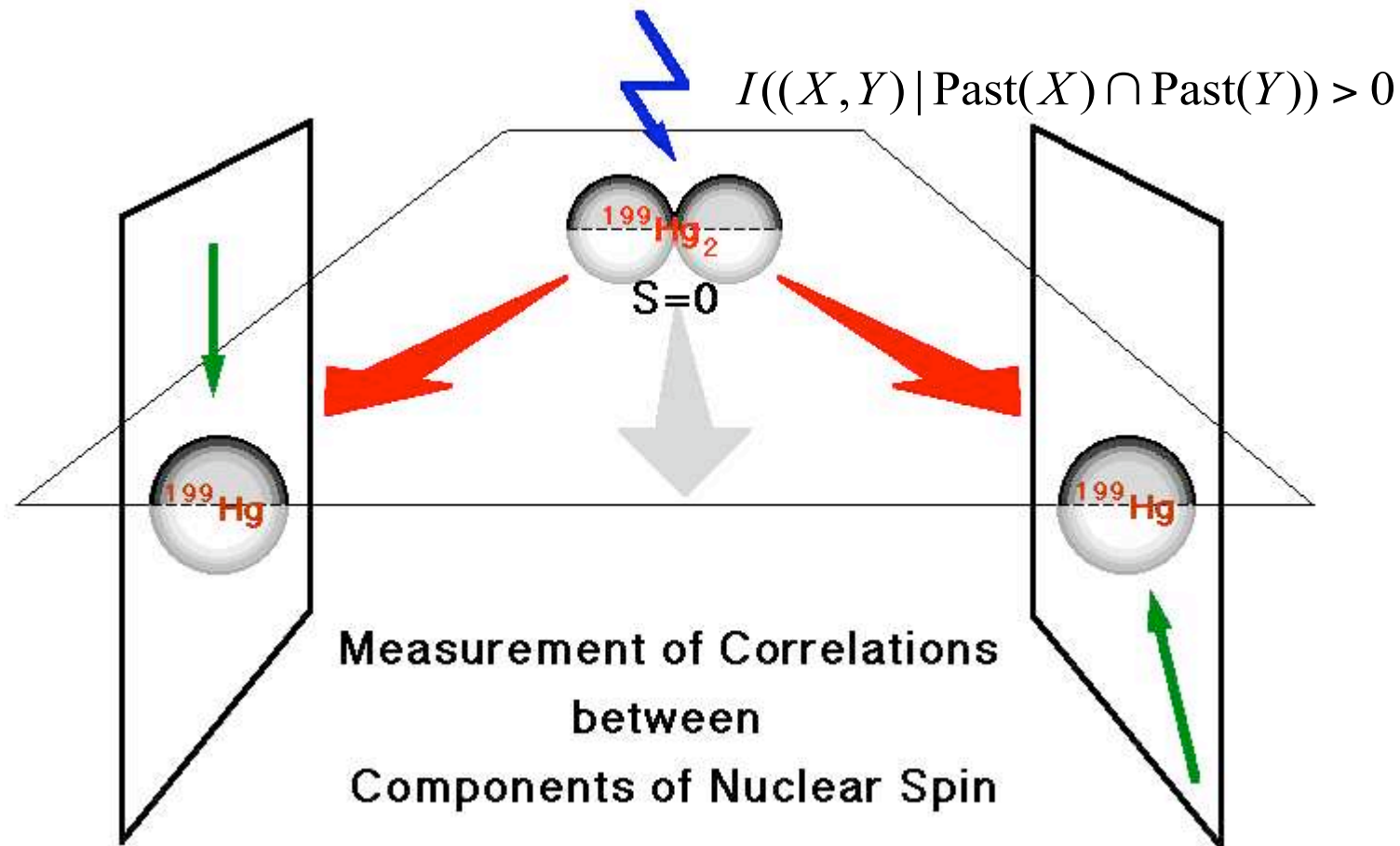
- Twin traveler paradox resolution?
 - For a common fixed past X et Y should be independent.

$$I((X, Y) | \text{Past}(X) \cap \text{Past}(Y)) = 0$$



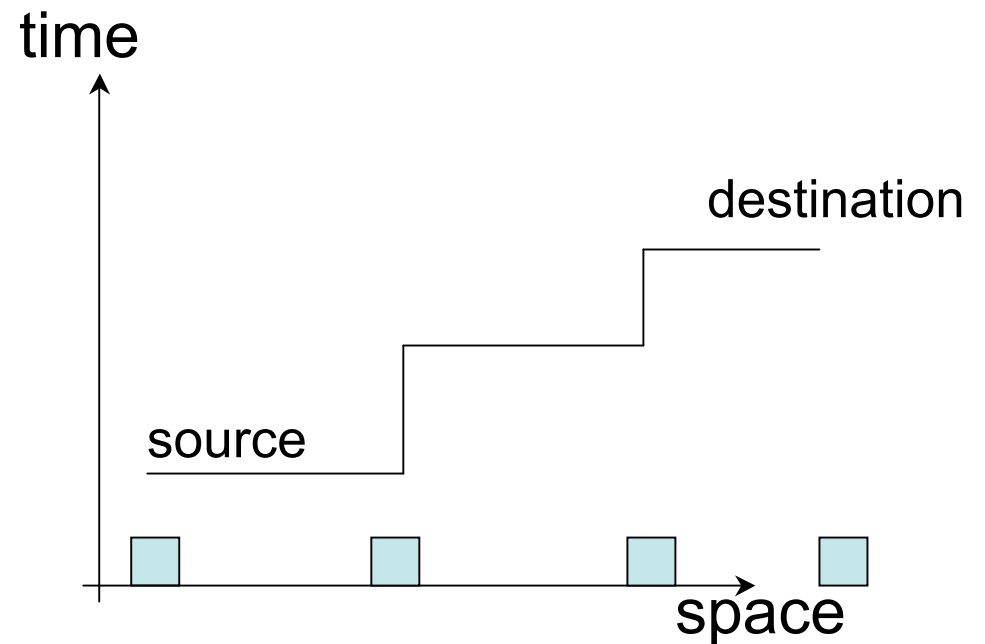
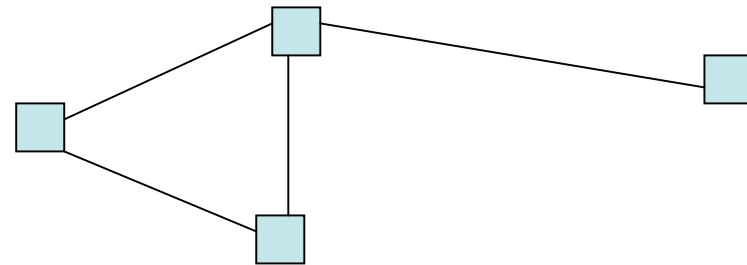
Bell theorem

- Information is non causal:



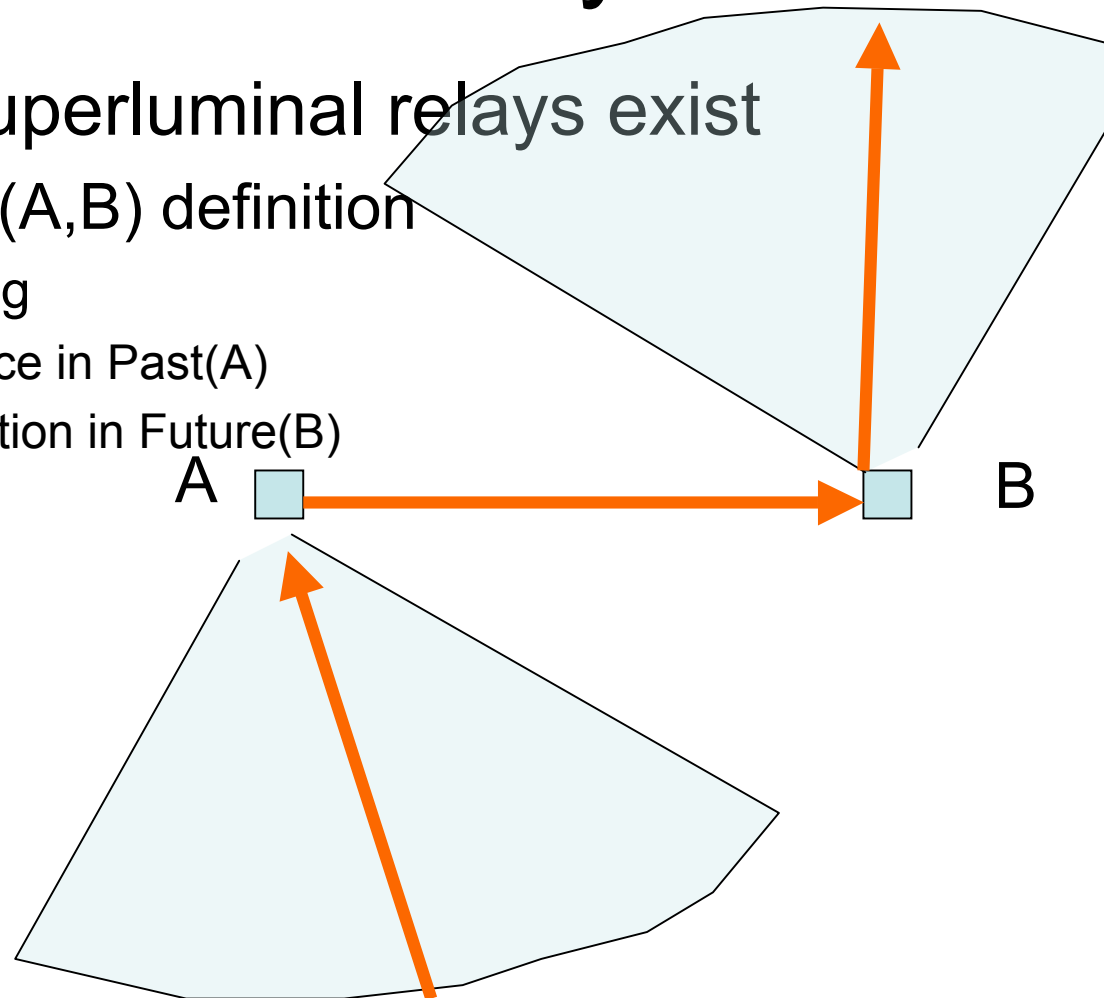
Network: information and space-time

- A Network is
 - Set of objects in physical space
 - relay information
 - from arbitrary source
 - to arbitrary destinations.
 - Concept of router
 - Concept of information propagation path

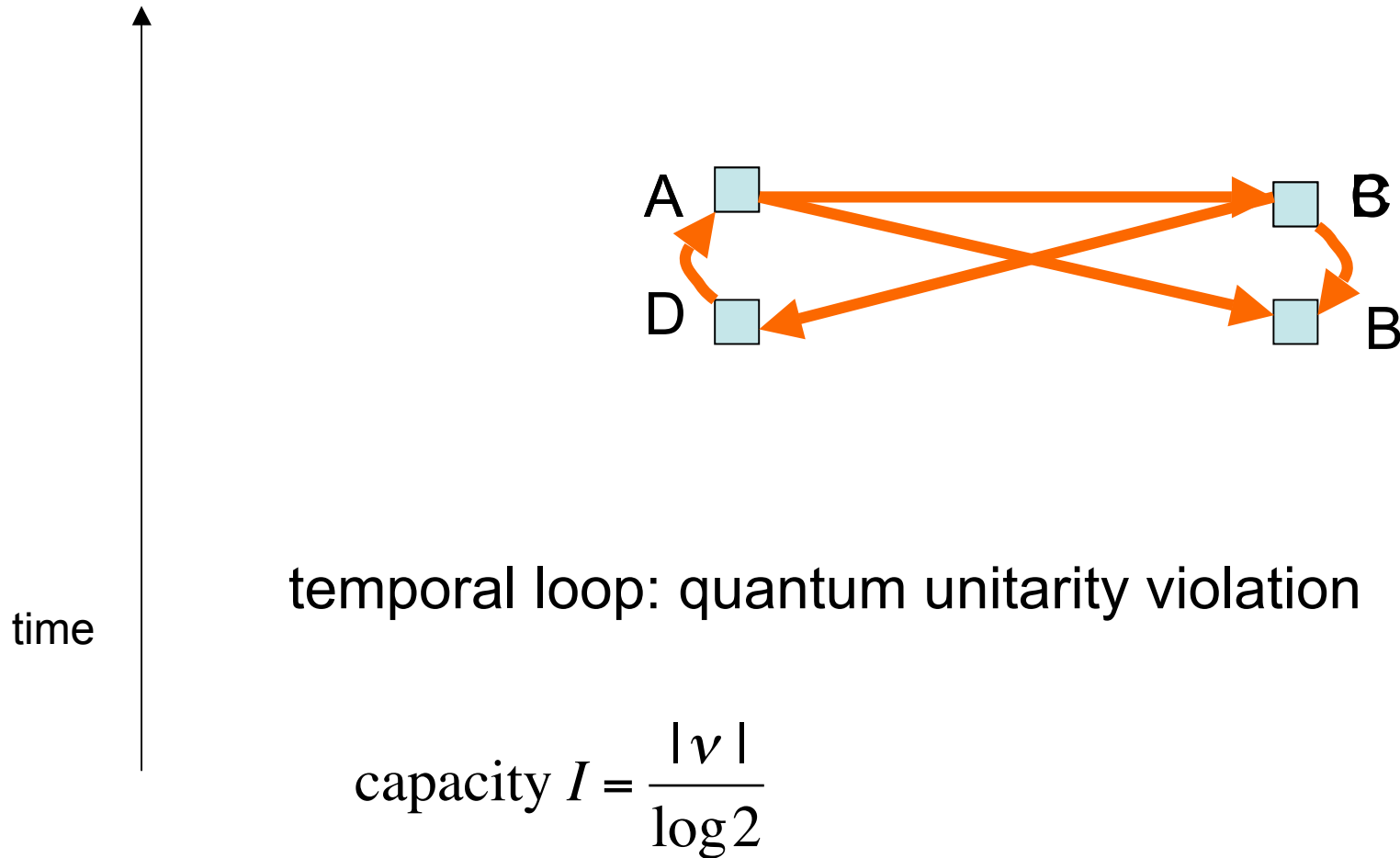


Space time relay versus information causality

- Can network of superluminal relays exist
 - Space time relay (A,B) definition
 - Can relay anything
 - From any source in Past(A)
 - To any destination in Future(B)



Space time relay versus information causality



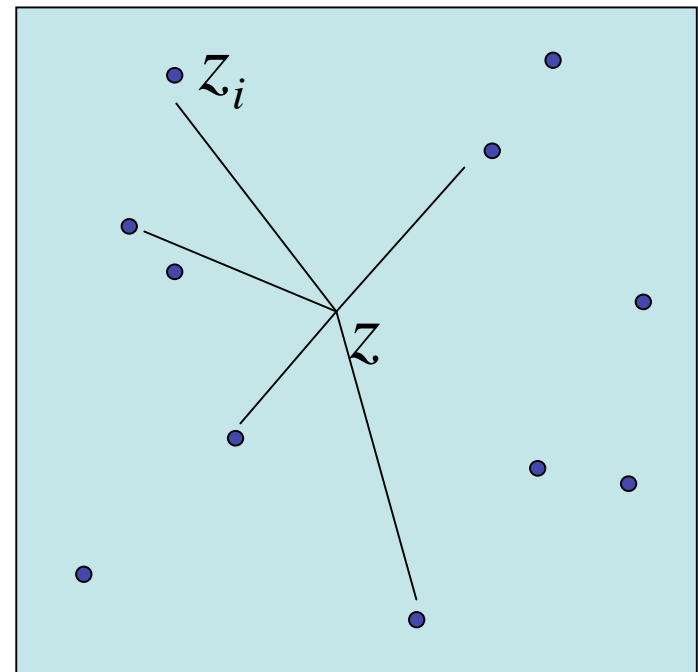
A wireless model

- Emitters are distributed as Poisson process in the plan

density λ

– Signals sum

$$S = \sum_i |z - z_i|^{-\alpha}$$

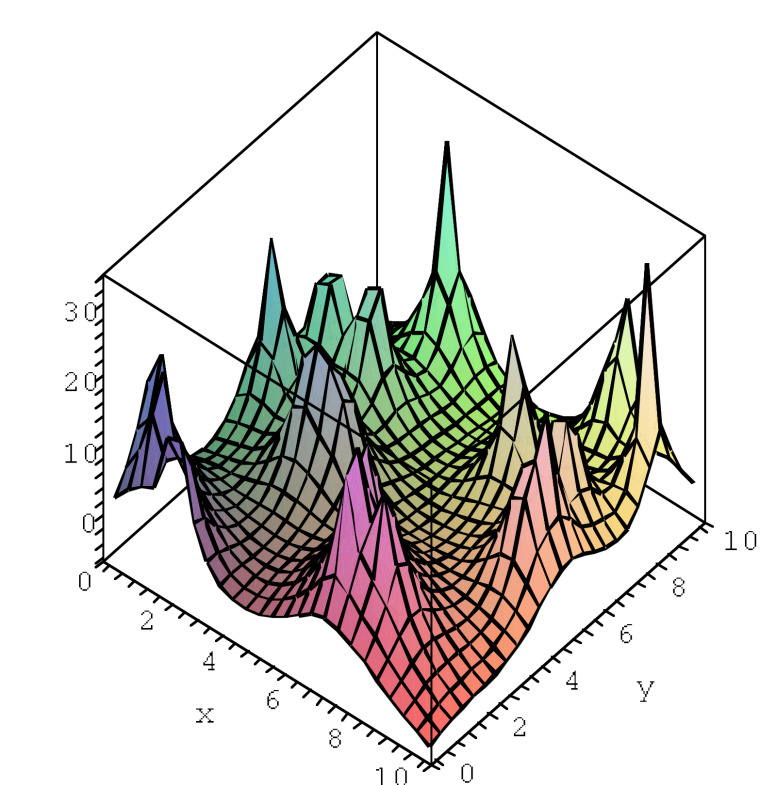


A wireless model

- Signal distribution

$$E(e^{-\theta S}) = \exp(-\lambda \pi \Gamma(1 - \gamma) \theta^\gamma)$$

$$\gamma = \frac{2}{\alpha}$$



Wireless information theory

- Wireless capacity,
 - emitters send independent information

$$I_0 = E \left(\sum_i \log_2 \left(1 + \frac{|z - z_i|^\alpha}{\sum_{j \neq i} |z - z_j|^\alpha} \right) \right)$$

- Computable and invariant even with random fading

$$I_0 = \frac{\gamma}{\log 2}$$

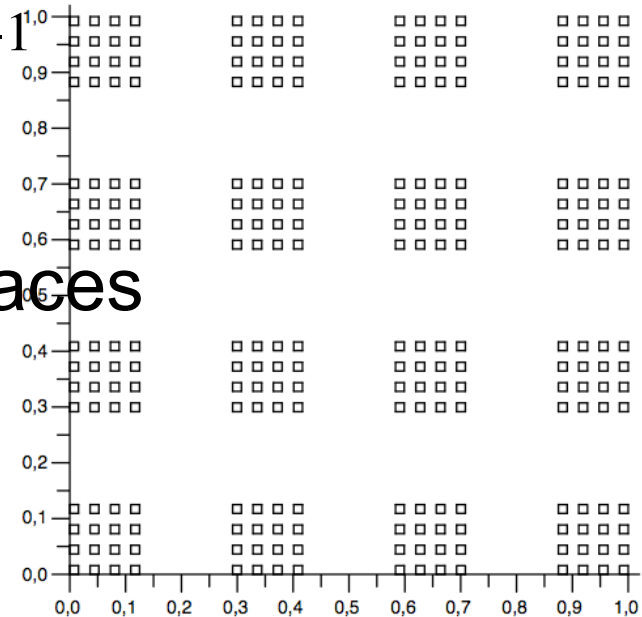
Wireless Shannon

- Invariant with dimensions

$$I_0 = \frac{\alpha}{D} (\log 2)^{-1}$$

– Works also with fractal spaces

- $D=4/3$



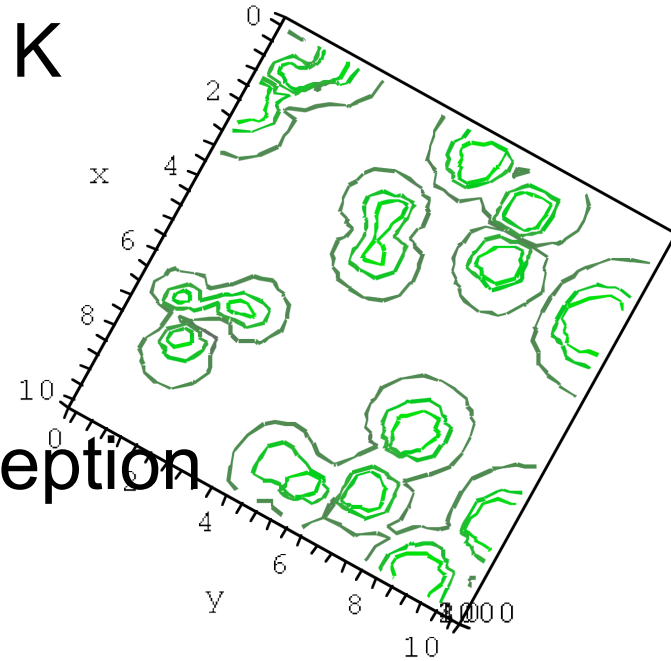
Wireless Space capacity

- With signal over noise ratio K requirement

$$I(K) = \frac{\sin(\gamma\pi)}{\gamma\pi} K^{-\gamma}$$

- Average area of correct reception

$$\sigma(K) = \frac{I(K)}{\lambda}$$



$$\sigma(10) \approx 0.037066$$

Wireless Space capacity

- Reception probability vs distance

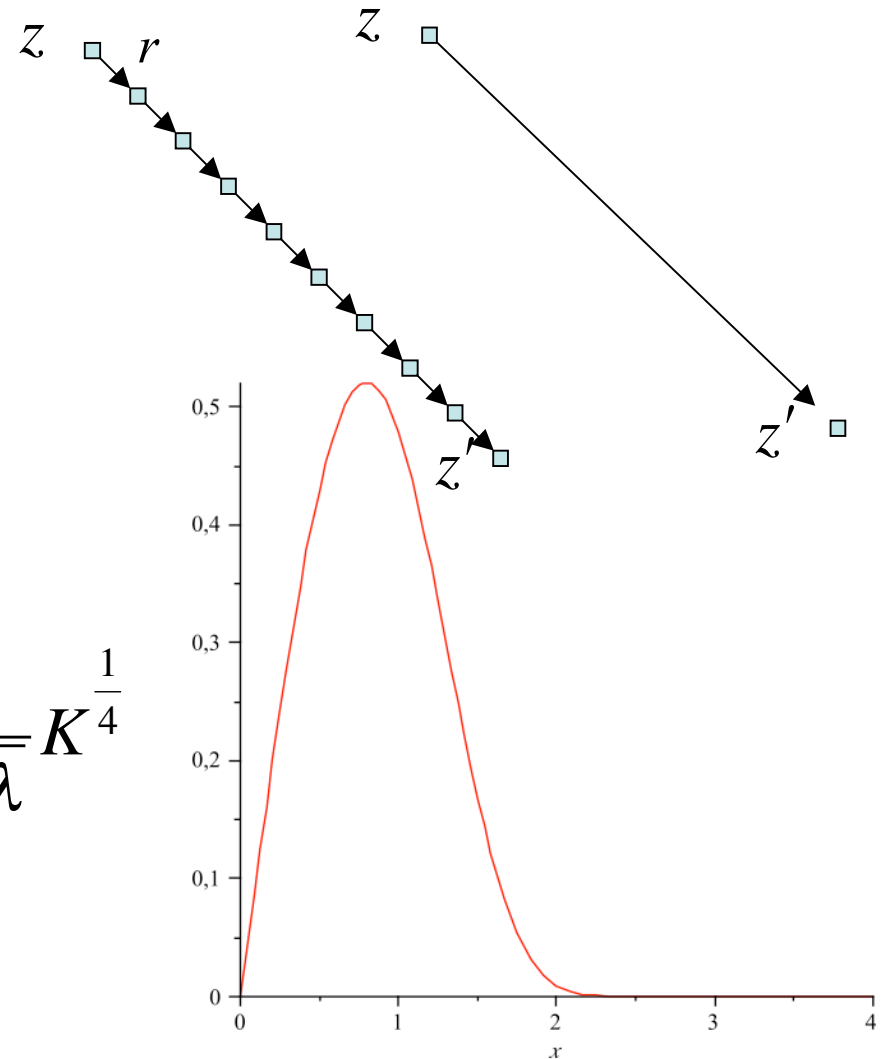
$$p(r, \lambda, K) = p(r\sqrt{\lambda}K^{-\frac{1}{4}}, 1, 1)$$

$$p(r, 1, 1) = \sum_n (-1)^n \frac{\sin(\pi n \gamma)}{\pi} \frac{\Gamma(n \gamma)}{n!} r^{2n}$$

$$p(r, 1, 1) = 1 - \operatorname{erf}\left(\frac{r^2}{2}\right) \text{ when } \alpha = 4$$

- Optimal routing radius

$$r_m = \arg \max_{r>0} \{rp(r, \lambda, K)\} = \frac{r_1}{\sqrt{\lambda}} K^{\frac{1}{4}}$$



Wireless Space capacity

- Average number of retransmissions

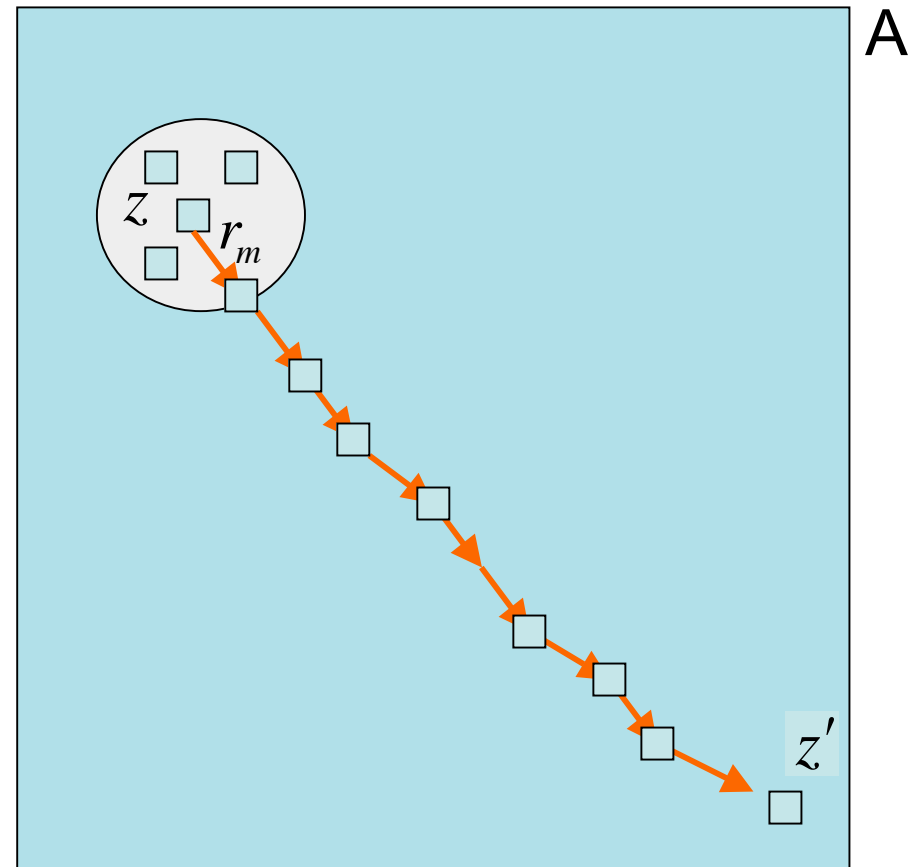
$$\frac{|z - z'|}{r_m P(r_m)}$$

- Net traffic density

$$\rho = \lambda \frac{r_m P(r_m)}{E(|z - z'|)}$$

- True if neighborhood is dense enough

$$\pi r_m^2 \frac{N}{A} > \log N$$

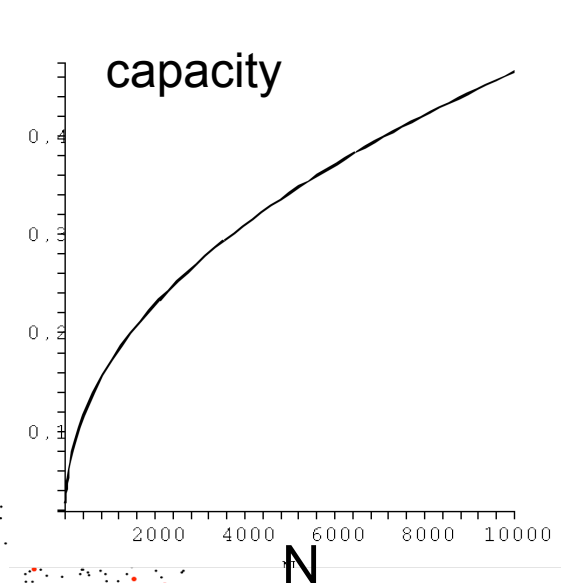
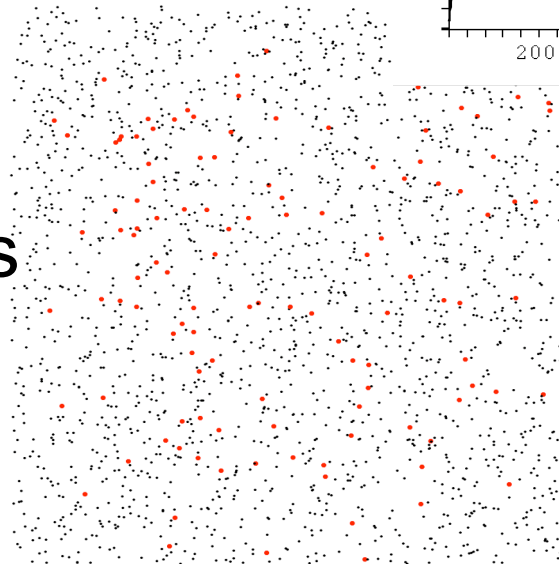


Space capacity result (Gupta-Kumar 2000)

- The capacity increases with the density

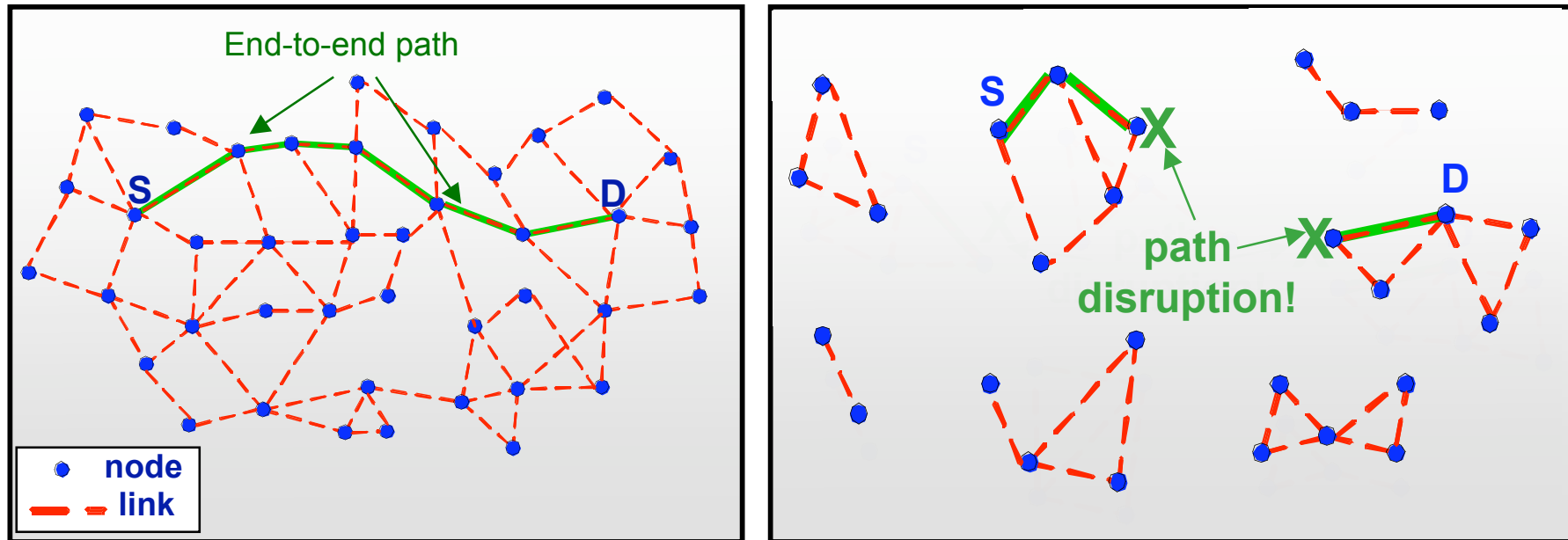
$$A\rho = \sqrt{\pi}r_1^2 p_1 \frac{\sqrt{A}}{E(|z - z'|)} \sqrt{\frac{N}{\log N}} = O\left(\sqrt{\frac{N}{\log N}}\right)$$

- Massively dense wireless networks



Time capacity paradox

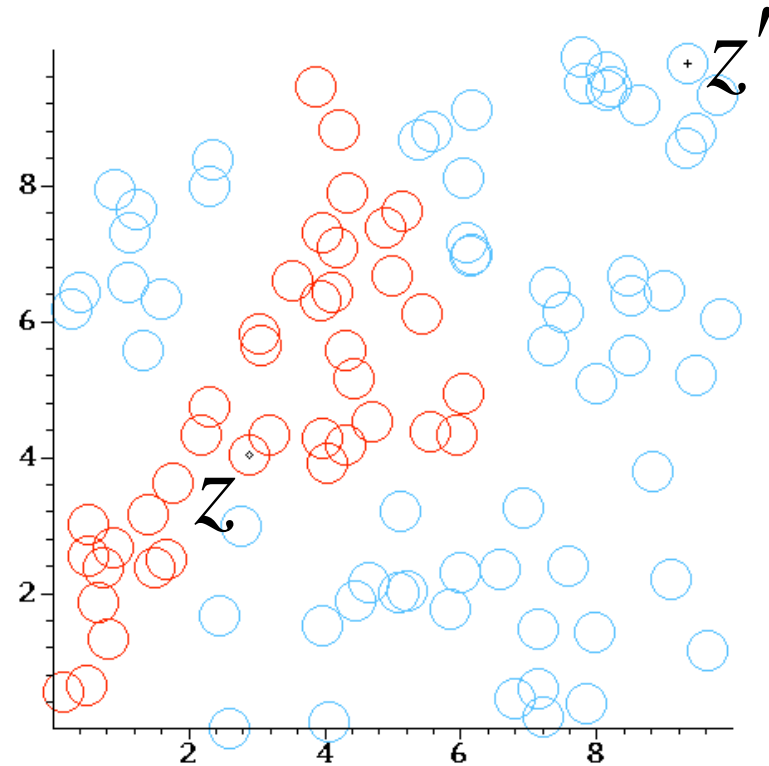
- Mobility can create capacity in sparse networks



- Delay Tolerant Networks

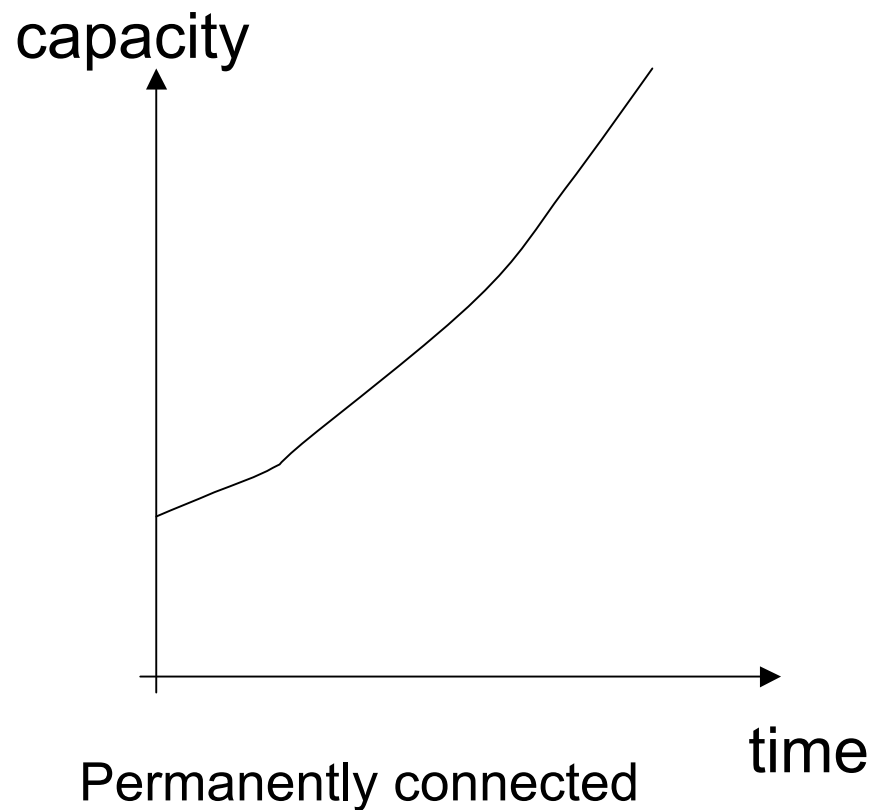
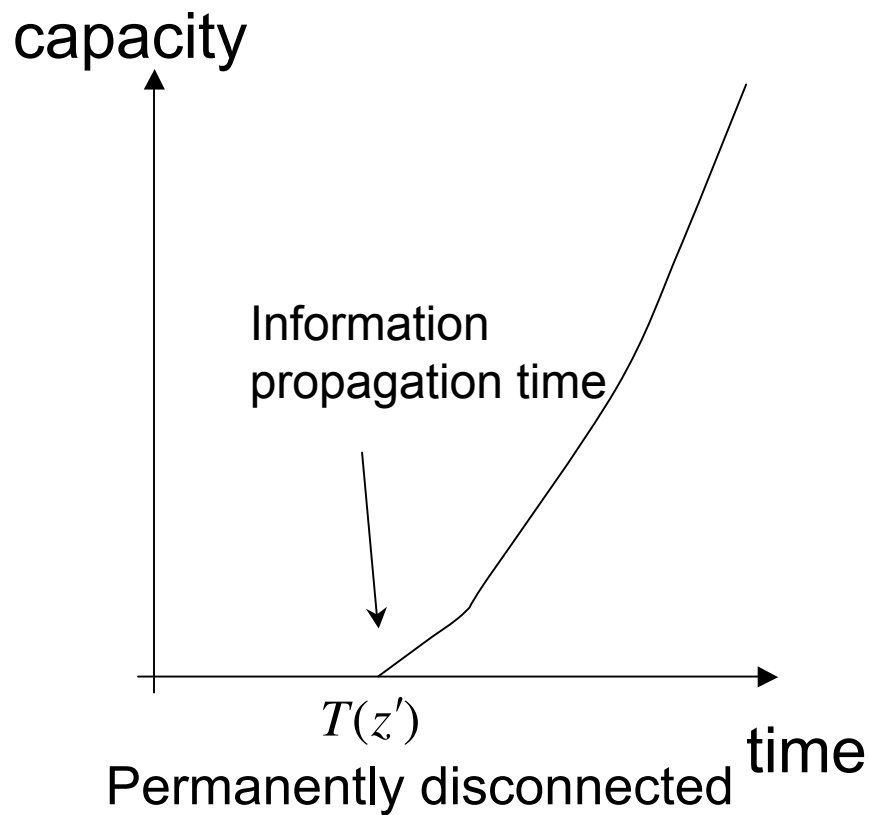
Information propagation speed

- Unit disk graph model
- Random walk mobility model



Time capacity paradox

- Mobility creates capacity



Information propagation speed

- Upper bound of information propagation speed

– Any quantity c such that

$$\lim_{z' \rightarrow \infty} P(T(z') < \frac{|z - z'|}{c}) = 0$$

– Is the smallest ratio $\frac{\theta}{\rho}$ in the kernel of

$$D(\rho, \theta) = \sqrt{(\tau + \theta)^2 - \rho^2 v^2} - \tau - \frac{2\pi v s I_0(\rho)}{1 - \pi v \frac{2}{\rho} I_1(\rho)}$$

$I_k()$ are modified Bessel functions

Information propagation speed

