## Homework 1. Basic Logic

1a. $p \vee(\overline{r \vee q}), \overline{r \vee q}=\neg(r \vee q)=\neg r \wedge \neg q$.

Table 1: 1a

| p | q | r | $r \vee q$ | $\neg(r \vee q)$ | $p \vee(\overline{r \vee q})$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T | T | T | T | F | T |
| T | T | F | T | F | T |
| T | F | T | T | F | T |
| T | F | F | F | T | T |
| F | T | T | T | F | F |
| F | T | F | T | F | F |
| F | F | T | T | F | F |
| F | F | F | F | T | T |

1 b .

Table 2: 1b

| p | q | r | $\neg q$ | $p \wedge \neg q$ | $(p \wedge \neg q) \rightarrow r$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T | T | T | F | F | T |
| T | T | F | F | F | T |
| T | F | T | T | T | T |
| T | F | F | T | T | F |
| F | T | T | F | F | T |
| F | T | F | F | F | T |
| F | F | T | T | F | T |
| F | F | F | T | F | T |

2. 

$$
\begin{aligned}
& (p \wedge q) \rightarrow(p \vee q) \\
\equiv & \neg(p \wedge q) \vee(p \vee q) \\
\equiv & (\neg p \vee \neg q) \vee(p \vee q) \\
\equiv & (\neg p \vee p) \vee(\neg q \vee q) \\
\equiv & T \vee T \\
\equiv & T
\end{aligned}
$$

Table 7
De Morgan's laws
Associative laws
Negation laws
Domination laws

3a. $P(x, y): x+y=5$ where $x, y$ are positive integers. $\forall x \forall y P(x, y)$. This statement is false.
Choose $x=1, y=1$, then we have,
$1+1 \geq 5$
$2 \geq 5$
which is false.
3b. $\forall x \exists y P(x, y)$. This statement is true.

We can translate this as, for all $x$, there is some $y$, such that $x+y \geq 5$. Let $x=1$, which is the min possible value. If we choose $y \geq 4$, then this statement holds as it only needs to be true for some $x$.
4. Pushing our negation all the way through we have,

$$
\begin{aligned}
& \forall x \exists y P(x, y) \\
\equiv & \exists x \neg \exists y P(x, y) \\
\equiv & \exists x \forall y \neg P(x, y)
\end{aligned}
$$

Hence, it follows that just (c) is equivalent.

