# Implementation

Major Internet routing protocols:

- RIP (v1 and v2): intra-domain, Bellman-Ford
  - $\rightarrow$  also called "distance vector"
  - $\rightarrow$  metric: hop count
  - $\rightarrow$  UDP
  - → nearest neighbor advertisement
  - → popular in small intra-domain networks
- OSPF (v1 and v2): intra-domain, Dijkstra
  - $\rightarrow$  also called "link state"
  - $\rightarrow$  metric: average delay
  - $\rightarrow$  directly over IP: protocol number 89
  - $\rightarrow$  broadcasting via flooding
  - $\rightarrow$  popular in larger intra-domain networks

- IS-IS: intra-domain, Dijkstra
  - $\rightarrow$  "link state"
  - → directly over link layer (e.g., Ethernet)
  - $\rightarrow$  more recently: also available over IP
  - $\rightarrow {\rm flooding}$
  - $\rightarrow$  popular in larger intra-domain networks

# BGP (Border Gateway Protocol):

 $\longrightarrow$  inter-domain routing

# Autonomous System B Peering Border Routers

- —— "peering" between two domains
- → includes customer-provider relationship
- → Internet exchanges: multiple domains

# • CIDR addressing

- $\rightarrow$  i.e., a.b.c.d/x
- $\rightarrow$  Purdue: 128.10.0.0/16, 128.210.0.0/16, 204.52.32.0/20
- → check at www.iana.org (e.g., ARIN for US)
- Metric: policy
  - $\rightarrow$  e.g., shortest-path, trust, pricing
  - → meaning of "shortest": delay, router hop, AS hop
  - → mechanism: path vector routing
  - $\rightarrow$  BPG update message

# BGP route update:

→ BGP update message propagation

BGP update message:

$$ASNA_k \rightarrow \cdots \rightarrow ASNA_2 \rightarrow ASNA_1$$
; a.b.c.d/x

Meaning: ASN  $A_1$  (with CIDR address a.b.c.d/x) can be reached through indicated path

- → "path vector"
- $\longrightarrow$  called AS-PATH

# Some AS numbers:

- Purdue: 17
- BBN: 1
- UUNET: 701
- Level3: 3356
- Abilene (aka "Internet2"): 11537

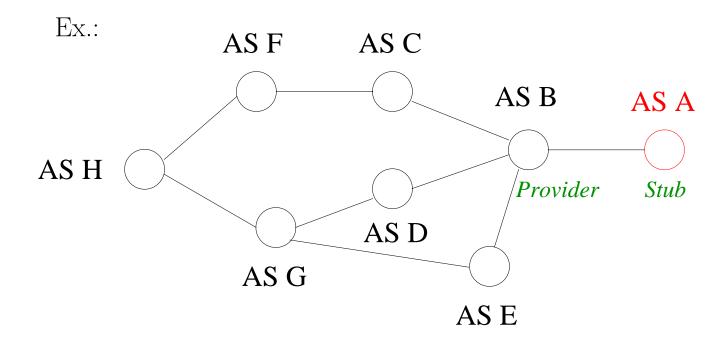
### Policy:

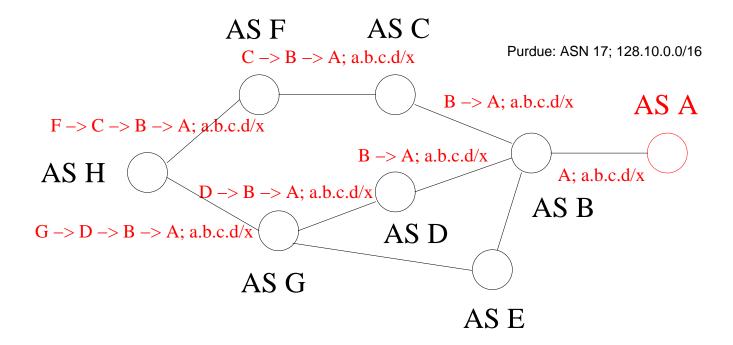
• if multiple AS-PATHs to target AS are known, choose one based on policy

- → e.g., shortest AS path length, cheapest, least worrisome
- advertise to neighbors target AS's reachability
  - $\rightarrow$  also subject to policy
  - $\rightarrow$  no obligation to advertise
  - $\rightarrow$  specifics depend on bilateral contract (SLA)

# SLA (service level agreement):

- $\longrightarrow$  bandwidth (e.g., 1 Gbps, OC-3, DS3
- $\longrightarrow$  delay (e.g., avrg. 25ms US), loss (e.g., 0.05%)
- → pricing (e.g., 1 Mbps: below \$100)
- $\longrightarrow$  availability (e.g., 99.999%)
- $\longrightarrow$  etc.





### Performance

Route update frequency:

- → routing table stability vs. responsiveness
- → rule: not too frequently
- $\longrightarrow$  30 seconds
- → stability wins
- → hard lesson learned from the past (sub-second)
- $\longrightarrow$  legacy: TTL

Other factors for route instability:

- → selfishness (e.g., fluttering)
- → BGP's vector path routing: inherently unstable
- → more common: slow convergence
- → target of denial-of-service (DoS) attack

### Route amplification:

- $\longrightarrow$  shortest AS path  $\neq$  shortest router path
- → e.g., may be several router hops longer
- → AS graph vs. router graph
- → inter- vs. intra-domain routing: separate subsystems
- → policy: company in Denmark

# Route asymmetry:

- → routes are not symmetric
- $\longrightarrow$  estimate: > 50%
- → mainly artifact of inter-domain policy routing
- → various performance implications
- → source traceback

# Black holes:

---- persistent unreachable destination prefixes

→ BGP routing problems

→ further aggrevated by DNS

---- purely application layer: end system problem