An efficient, incremental, automatic garbage collector Deutsch & Bobrow

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Outline

- Introduction
- Advantage/Disadvantage of Reference counting
- Methodology
- Conclusion

Introduction – Reference counting

Reference counting (RC) contains

A counter for each object

- The counter means how many other objects reference it.
- When the value of the counter becomes 0

The object is garbage and can be reclaimed.

Advantage of RC

- GC costs are distributed throughout the computation
- Unreferenced structures can be reclaimed immediately
- Do not need to know the roots of the program
 - It can reclaim some memory even if some parts of the system are unavailable

Disadvantage of RC

 Redefine all read / write operations in order to manipulate reference counts

- Undesirable, i.e. traverse a list

- Atomic operations to update reference counts
- Circular structures are not reclaimed
- Space overhead to store reference counts

Three ways to improve RC

• Deferral RC

Defer the identification of garbage objects to a reclamation phase

Coalescing

Eliminate unnecessary temporal adjustments of reference counts

• Buffering

Buffer the adjustment of reference counts for later processing

Deferred RC

- Observation:
 - The majority of objects have a RC of 1
- Solution:
 - Record transactions that affect the accessibility in the transaction file
 - Allocate a new cell
 - *Create* a pointer to a cell
 - **Destroy** a pointer to a cell
 - Store the reference count for objects with 2 or more references in a multireference table (MRT)

Deferred RC

- Observation:
 - The majority of pointer loads are to local and temporary variables
- Solution:
 - Do not count local variable references
 - Counts are no longer accurate
 - Need a zero count table (ZCT) for objects with reference counts of 0 but may be referenced from locals

Transaction file processing

```
for ( Transaction* t = FirstTransaction(); t; t = Next(t) ) {
```

```
if (IsAllocate(t)) {
```

AddToZCT(t);

```
} else if ( lsCreatePtr(t) ) {
```

```
if (IsInMRT(t))
```

```
else
```

```
SetMRTReferenceCount(t,2);
```

```
} else if ( IsDestroyPtr(t) ) {
```

```
if (IsInMRT(t)) {
```

```
if (GetMRTCount(t) == 2)
```

```
RemoveFromMRT(t);
```

else

```
DecrementMRTReference(t);
```

```
} else {
```

AddToZCT(t);

// Upon allocation, simply add to ZCT

```
if ( !RemoveFromZCT(t) ) { // Remove from ZCT if present
```

```
IncrementMRTReference(t); // Add reference to MRT or increment its ref count
```

```
// Add to ZCT if this was the last reference
```

Reclamation

- Reference count = 0 (in ZCT)
 - Not count yet referenced from local variables
 - "May" be live
- Any objects in ZCT are reclaimable if it is not reachable from the stack.
 - Variable reference table (VRT) contains all the pointers from the stack.
 - Reclaim (obj ∈ ZCT) if IsInVRT(obj) = true

Reclamation

- Upon freeing an object
 - Decrement the references counts of objects which it may point to
- Free a large data structure
 - May defer the decremented operations in order to disperse the disruptive effects

Coalescing

- Eliminate unnecessary temporal adjustments of reference counts
 - allocate create
 - destroy- create

Linearizing Garbage Collection

- Collect circular structures
- Generational garbage collector
 - In depth-first traversal order
 - Extended MRT
 - MRT must be updated before copying
 - (taking the stack into account)
 - Store the relocation address for each entry of MRT (?)
- It improves locality



- Old data can be released by reference counts
- Full linearization will be performed if the storage needs to be reclaimed
- MRT and ZCT are segmented by address

Further Observation



Reference

- www.cs.utexas.edu/users/mckinley/395Tmm/ talks/Mar-9-DRC.pdf
- Richard Jones, Antony Hosking, and Eliot Moss, The Garbage Collection Handbook