

# Boolean Circuits

**Definition 1. Syntax:** *A Boolean circuit is a (directed acyclic) graph  $C = (V, E)$  where the nodes (gates) in  $V = \{1, \dots, n\}$  can be either a boolean connective or a variable or  $\{true, false\}$  and edges are of the form  $(i, j)$  where  $i < j$ . Gates with no incoming edges (variables or true/false) are input gates. Output gates have no outgoing edges (typically one — numbered  $n$ ). The gates  $\wedge$  and  $\vee$  have indegree 2 and gate  $\neg$  has indegree 1.*

**Semantics:** *Specifies a truth value to each appropriate truth assignment of the variable gates. The value of the circuit is value of the output gate computed inductively "bottom up".*

# Computational Problems

CIRCUIT SAT: Given a circuit  $C$ , is there a truth assignment  $T$  appropriate to  $C$  such that  $T(C) = \text{true}$ .

Computationally equivalent to SAT,  $\in NP$ .

CIRCUIT VALUE: Given a circuit with no variable gates (i.e., only true/false inputs) is the value of the gate true.

$\in P$ .

# Size of Boolean Circuits

**Theorem 1. [Shannon]** *For any  $n \geq 2$  there is a  $n$ -ary boolean function  $f$  such that no boolean circuit with  $\frac{2^n}{2n}$  or fewer gates can compute it.*

**Proof:** We will give a counting argument.

Suppose not. Then for some  $n \geq 2$  all  $n$ -ary boolean functions can be computed with  $m = 2^n/2n$  or fewer gates.

The number of circuits with  $n$  variables and  $m$  or fewer gates is bounded by  $((n + 5).m^2)^m$ .

The number of different  $n$ -ary boolean functions is  $2^{2^n}$ .

$$((n + 5).m^2)^m < 2^{2^n} \text{ if } m = 2^n/2n.$$

That is, two *different* boolean functions are computed by the same circuit. Contradiction. QED

# Compactness Theorem

**Theorem 2.** *A set  $S$  of expressions is satisfiable iff every finite subset of  $S$  is satisfiable.*

# Resolution

A method to determine whether a Boolean expression  $\phi$  in CNF form is satisfiable.

0.  $\phi^* = \phi$ .

1. Repeat until possible:

Let  $(x \vee C)$  and  $(\neg x \vee D)$  be two clauses in  $\phi^*$ .

Add the clause  $(C \vee D)$  (*resolvent*) to  $\phi^*$ .

2. If  $\phi^*$  contains the empty clause then  $\phi$  is not satisfiable else satisfiable.