

CS490D:
Introduction to Data Mining
Chris Clifton

January 16, 2004
Data Warehousing



Data Warehousing and OLAP
Technology for Data Mining

- **What is a data warehouse?**
- A multi-dimensional data model
- Data warehouse architecture
- Data warehouse implementation
- Further development of data cube technology
- From data warehousing to data mining



What is Data Warehouse?

- Defined in many different ways, but not rigorously.
 - A decision support database that is maintained [separately](#) from the organization's operational database
 - Support [information processing](#) by providing a solid platform of consolidated, historical data for analysis.
- “A data warehouse is a [subject-oriented](#), [integrated](#), [time-variant](#), and [nonvolatile](#) collection of data in support of management's decision-making process.”—W. H. Inmon
- Data warehousing:
 - The process of constructing and using data warehouses

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Data Warehouse—Subject-Oriented

- Organized around major subjects, such as [customer](#), [product](#), [sales](#).
- Focusing on the modeling and analysis of data for decision makers, not on daily operations or transaction processing.
- Provide [a simple and concise](#) view around particular subject issues by [excluding data that are not useful in the decision support process](#).

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Data Warehouse—Integrated

- Constructed by integrating multiple, heterogeneous data sources
 - relational databases, flat files, on-line transaction records
- Data cleaning and data integration techniques are applied.
 - Ensure consistency in naming conventions, encoding structures, attribute measures, etc. among different data sources
 - E.g., Hotel price: currency, tax, breakfast covered, etc.
 - When data is moved to the warehouse, it is converted.

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Data Warehouse—Time Variant

- The time horizon for the data warehouse is significantly longer than that of operational systems.
 - Operational database: current value data.
 - Data warehouse data: provide information from a historical perspective (e.g., past 5-10 years)
- Every key structure in the data warehouse
 - Contains an element of time, explicitly or implicitly
 - But the key of operational data may or may not contain “time element”.

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Data Warehouse—Non-Volatile

- A **physically separate store** of data transformed from the operational environment.
- Operational **update of data does not occur** in the data warehouse environment.
 - Does not require transaction processing, recovery, and concurrency control mechanisms
 - Requires only two operations in data accessing:
 - *initial loading of data* and *access of data*.

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Data Warehouse vs. Heterogeneous DBMS

- Traditional heterogeneous DB integration:
 - Build **wrappers/mediators** on top of heterogeneous databases
 - **Query driven** approach
 - When a query is posed to a client site, a meta-dictionary is used to translate the query into queries appropriate for individual heterogeneous sites involved, and the results are integrated into a global answer set
 - Complex information filtering, compete for resources
- **Data warehouse: update-driven, high performance**
 - Information from heterogeneous sources is integrated in advance and stored in warehouses for direct query and analysis

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Data Warehouse vs. Operational DBMS

- OLTP (on-line transaction processing)
 - Major task of traditional relational DBMS
 - Day-to-day operations: purchasing, inventory, banking, manufacturing, payroll, registration, accounting, etc.
- OLAP (on-line analytical processing)
 - Major task of data warehouse system
 - Data analysis and decision making
- Distinct features (OLTP vs. OLAP):
 - User and system orientation: customer vs. market
 - Data contents: current, detailed vs. historical, consolidated
 - Database design: ER + application vs. star + subject
 - View: current, local vs. evolutionary, integrated
 - Access patterns: update vs. read-only but complex queries

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OLTP vs. OLAP

	OLTP	OLAP
users	clerk, IT professional	knowledge worker
function	day to day operations	decision support
DB design	application-oriented	subject-oriented
data	current, up-to-date detailed, flat relational isolated	historical, summarized, multidimensional integrated, consolidated
usage	repetitive	ad-hoc
access	read/write index/hash on prim. key	lots of scans
unit of work	short, simple transaction	complex query
# records accessed	tens	millions
#users	thousands	hundreds
DB size	100MB-GB	100GB-TB
metric	transaction throughput	query throughput, response

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Why Separate Data Warehouse?

- High performance for both systems
 - DBMS— tuned for OLTP: access methods, indexing, concurrency control, recovery
 - Warehouse—tuned for OLAP: complex OLAP queries, multidimensional view, consolidation.
- Different functions and different data:
 - **missing data**: Decision support requires historical data which operational DBs do not typically maintain
 - **data consolidation**: DS requires consolidation (aggregation, summarization) of data from heterogeneous sources
 - **data quality**: different sources typically use inconsistent data representations, codes and formats which have to be reconciled

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Data Warehousing and OLAP Technology for Data Mining

- What is a data warehouse?
- **A multi-dimensional data model**
- Data warehouse architecture
- Data warehouse implementation
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- From data warehousing to data mining

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From Tables and Spreadsheets to Data Cubes

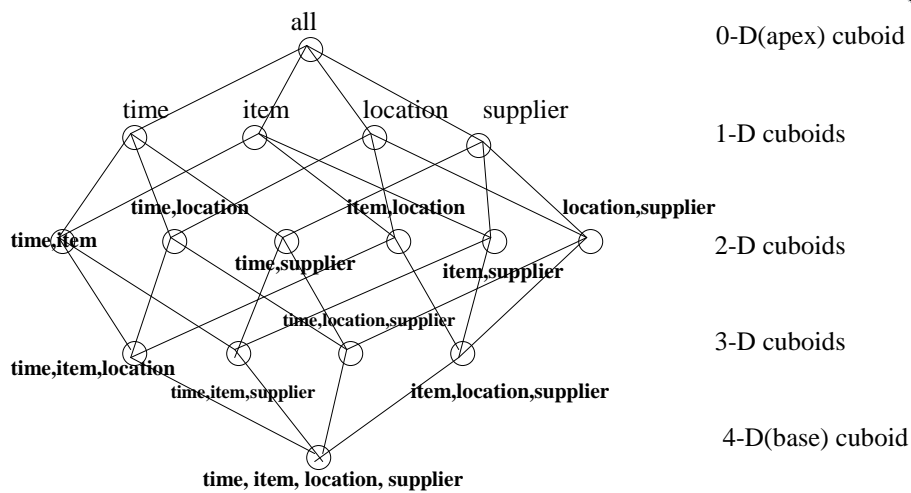
- A data warehouse is based on a **multidimensional data model** which views data in the form of a data cube
- A data cube, such as **sales**, allows data to be modeled and viewed in multiple dimensions
 - Dimension tables, such as **item** (*item_name, brand, type*), or **time**(*day, week, month, quarter, year*)
 - Fact table contains measures (such as **dollars_sold**) and keys to each of the related dimension tables
- In data warehousing literature, an n-D base cube is called a **base cuboid**. The top most 0-D cuboid, which holds the highest-level of summarization, is called the **apex cuboid**. The lattice of cuboids forms a **data cube**.

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Cube: A Lattice of Cuboids



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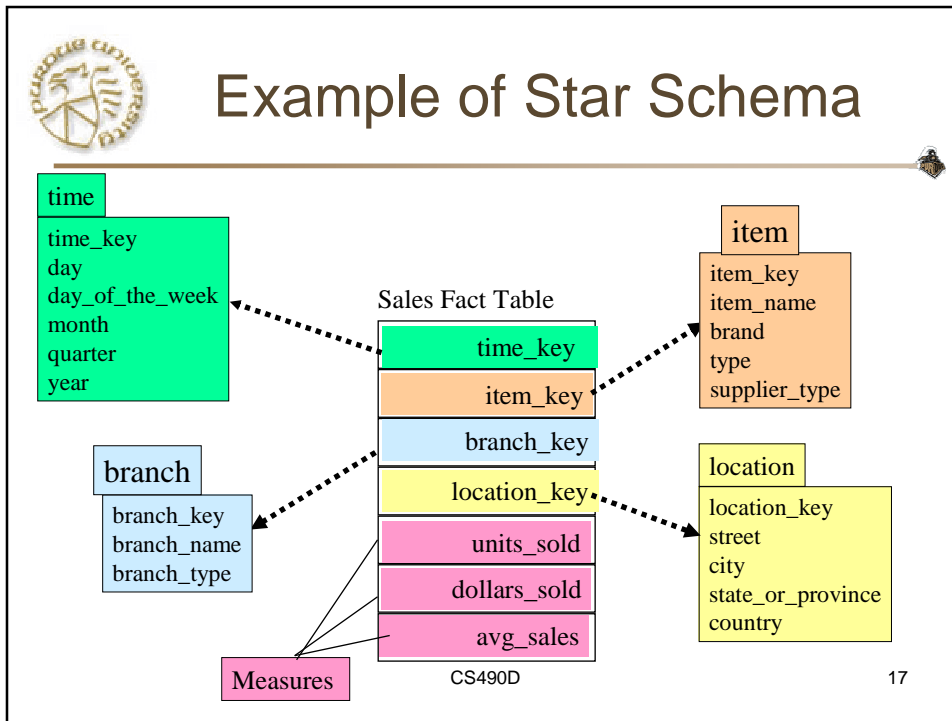


Conceptual Modeling of Data Warehouses

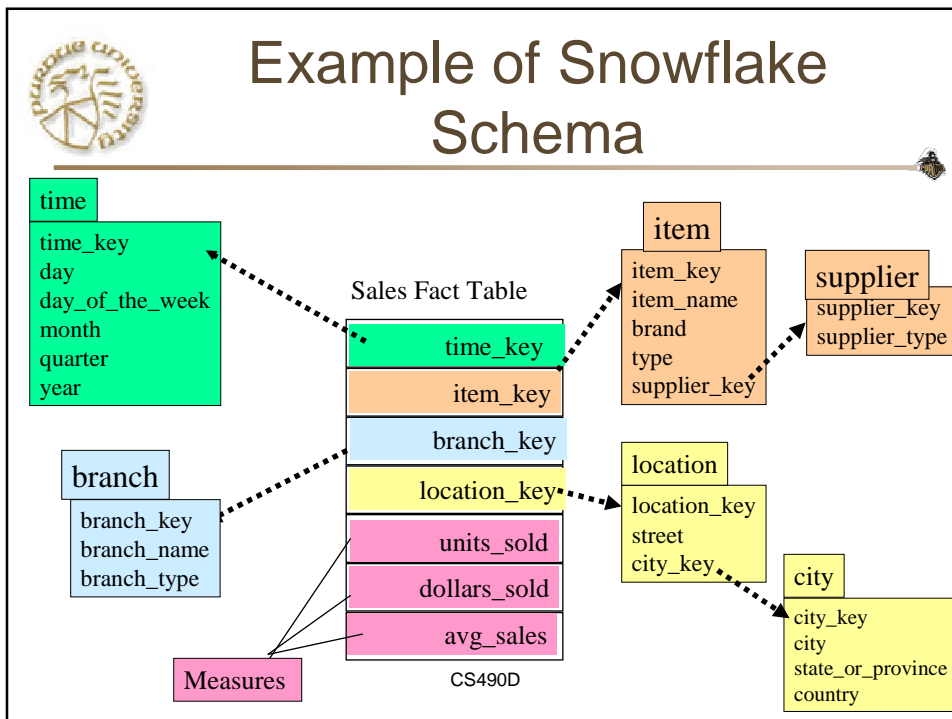
- Modeling data warehouses: dimensions & measures
 - [Star schema](#): A fact table in the middle connected to a set of dimension tables
 - [Snowflake schema](#): A refinement of star schema where some dimensional hierarchy is **normalized** into a set of smaller dimension tables, forming a shape similar to snowflake
 - [Fact constellations](#): Multiple fact tables share dimension tables, viewed as a collection of stars, therefore called **galaxy schema** or fact constellation

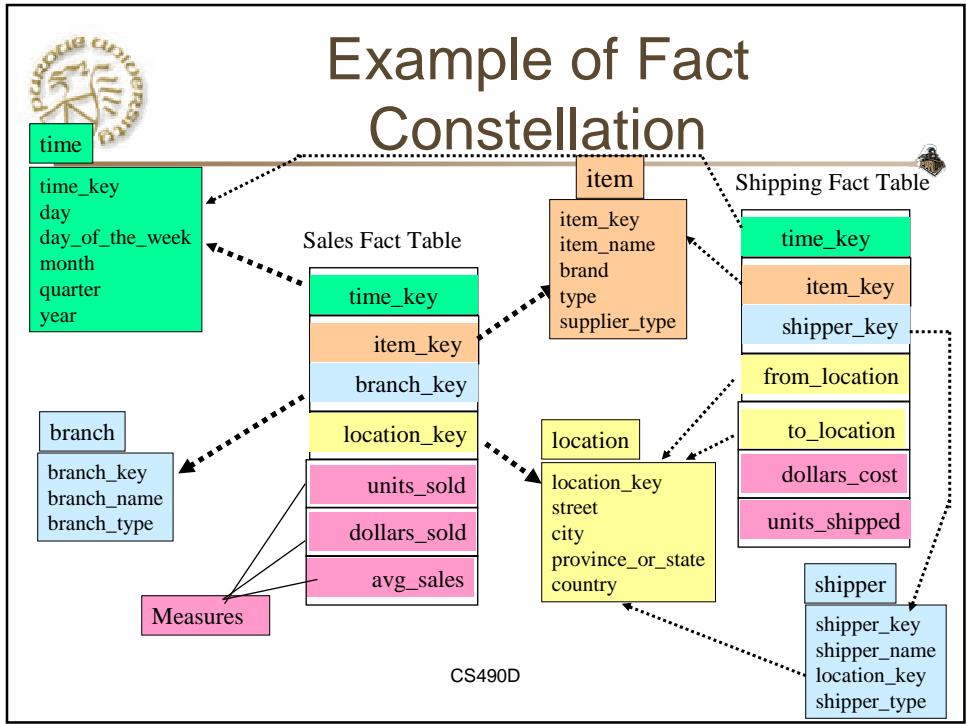


Example of Star Schema



Example of Snowflake Schema





A Data Mining Query Language: DMQL

- **Cube Definition (Fact Table)**
`define cube <cube_name> [<dimension_list>]:
 <measure_list>`
- **Dimension Definition (Dimension Table)**
`define dimension <dimension_name> as
 (<attribute_or_subdimension_list>)`
- **Special Case (Shared Dimension Tables)**
 - First time as “cube definition”
 - `define dimension <dimension_name> as
 <dimension_name_first_time> in cube
 <cube_name_first_time>`

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Defining a Star Schema in DMQL

```
define cube sales_star [time, item, branch, location]:
    dollars_sold = sum(sales_in_dollars), avg_sales =
        avg(sales_in_dollars), units_sold = count(*)
define dimension time as (time_key, day, day_of_week,
    month, quarter, year)
define dimension item as (item_key, item_name, brand,
    type, supplier_type)
define dimension branch as (branch_key, branch_name,
    branch_type)
define dimension location as (location_key, street, city,
    province_or_state, country)
```

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Defining a Snowflake Schema in DMQL

```
define cube sales_snowflake [time, item, branch, location]:
    dollars_sold = sum(sales_in_dollars), avg_sales =
        avg(sales_in_dollars), units_sold = count(*)
define dimension time as (time_key, day, day_of_week, month,
    quarter, year)
define dimension item as (item_key, item_name, brand, type,
    supplier(supplier_key, supplier_type))
define dimension branch as (branch_key, branch_name, branch_type)
define dimension location as (location_key, street, city(city_key,
    province_or_state, country))
```

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Defining a Fact Constellation in DMQL

```
define cube sales [time, item, branch, location]:
    dollars_sold = sum(sales_in_dollars), avg_sales = avg(sales_in_dollars),
    units_sold = count(*)
define dimension time as (time_key, day, day_of_week, month, quarter, year)
define dimension item as (item_key, item_name, brand, type, supplier_type)
define dimension branch as (branch_key, branch_name, branch_type)
define dimension location as (location_key, street, city, province_or_state,
    country)
define cube shipping [time, item, shipper, from_location, to_location]:
    dollar_cost = sum(cost_in_dollars), unit_shipped = count(*)
define dimension time as time in cube sales
define dimension item as item in cube sales
define dimension shipper as (shipper_key, shipper_name, location as location
    in cube sales, shipper_type)
define dimension from_location as location in cube sales
define dimension to_location as location in cube sales
```

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Measures: Three Categories

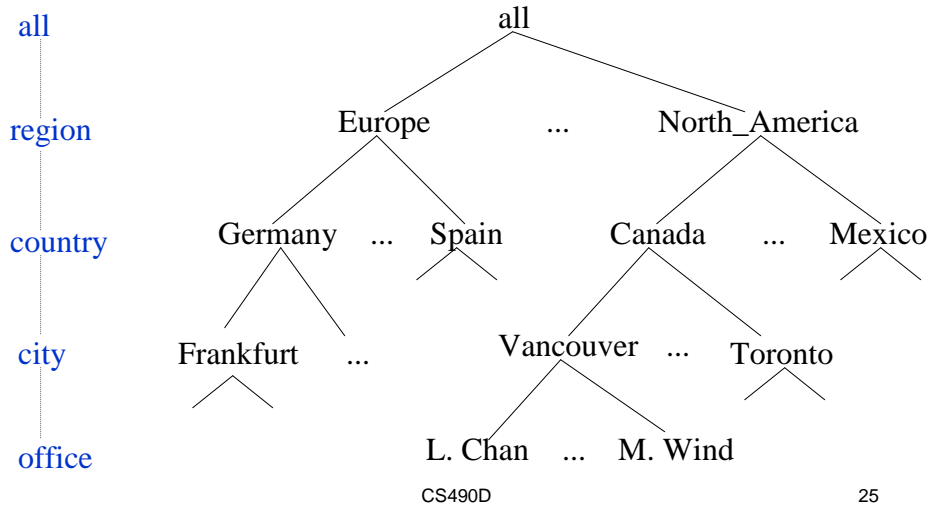
- **distributive**: if the result derived by applying the function to n aggregate values is the same as that derived by applying the function on all the data without partitioning.
 - E.g., count(), sum(), min(), max().
- **algebraic**: if it can be computed by an algebraic function with M arguments (where M is a bounded integer), each of which is obtained by applying a distributive aggregate function.
 - E.g., avg(), min_N(), standard_deviation().
- **holistic**: if there is no constant bound on the storage size needed to describe a subaggregate.
 - E.g., median(), mode(), rank().

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A Concept Hierarchy: Dimension (location)



View of Warehouses and Hierarchies

Warehouse

Level Name

Warehouse

Description

ANY

Europe

Belgium

France

Germany

Essen

Frankfurt

Spain

Sweden

United Kingdom

Far East

North America

Canada

Montreal

Toronto

Vancouver

Charles Loo Nam

Hari Krain

Kaley Gregson

Lee Chan

Malcom Young

Marthe White duck

Torey Wandiko

Mexico

United States

Specification of hierarchies

- Schema hierarchy
day < {month < quarter;
week} < year
- Set_grouping hierarchy
{1..10} < inexpensive

NUM

For Help, press F1

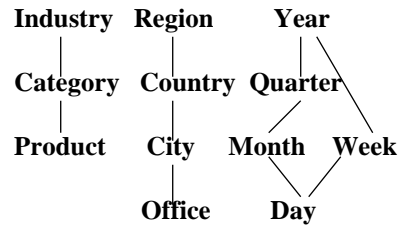
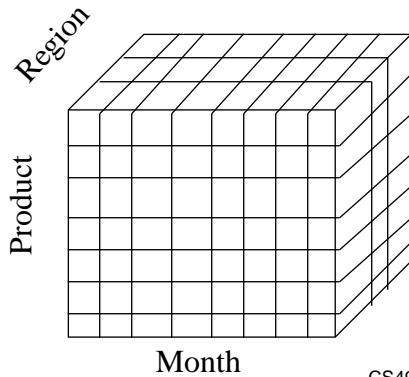
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Multidimensional Data

- Sales volume as a function of product, month, and region

Dimensions: Product, Location, Time
Hierarchical summarization paths

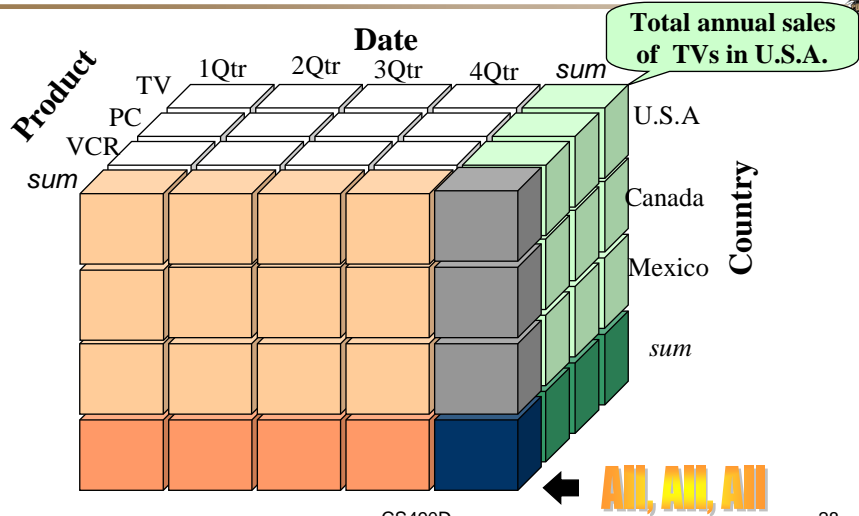


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A Sample Data Cube

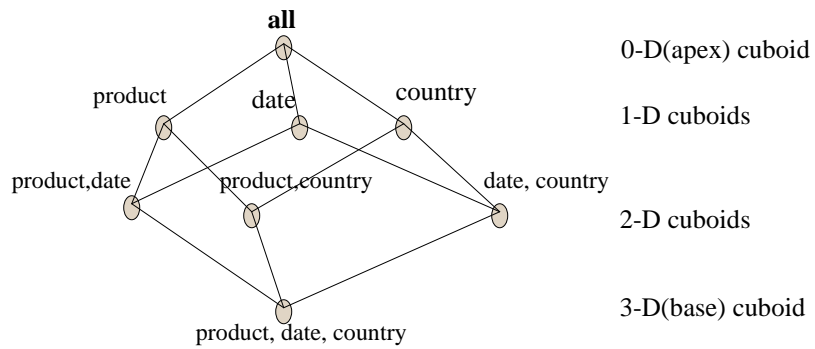


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Cuboids Corresponding to the Cube

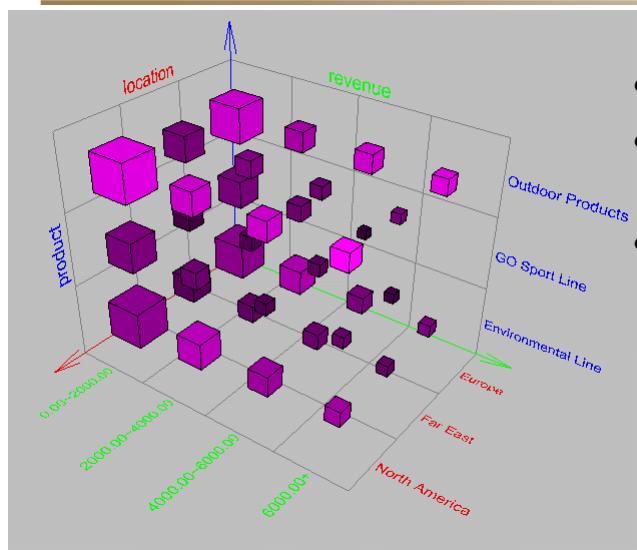


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Browsing a Data Cube



- Visualization
- OLAP capabilities
- Interactive manipulation

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Typical OLAP Operations

- Roll up (drill-up): summarize data
 - *by climbing up hierarchy or by dimension reduction*
- Drill down (roll down): reverse of roll-up
 - *from higher level summary to lower level summary or detailed data, or introducing new dimensions*
- Slice and dice:
 - *project and select*
- Pivot (rotate):
 - *reorient the cube, visualization, 3D to series of 2D planes.*
- Other operations
 - *drill across: involving (across) more than one fact table*
 - *drill through: through the bottom level of the cube to its back-end relational tables (using SQL)*

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Data Warehousing and OLAP Technology for Data Mining

- What is a data warehouse?
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Efficient Data Cube Computation

- Data cube can be viewed as a lattice of cuboids
 - The bottom-most cuboid is the base cuboid
 - The top-most cuboid (apex) contains only one cell
 - How many cuboids in an n-dimensional cube with L levels?

$$T = \prod_{i=1}^n (L_i + 1)$$

- Materialization of data cube
 - Materialize every (cuboid) (full materialization), none (no materialization), or some (partial materialization)
 - Selection of which cuboids to materialize
 - Based on size, sharing, access frequency, etc.



Cube Operation

- Cube definition and computation in DMQL

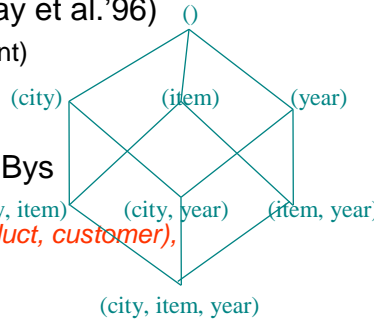

```
define cube sales[item, city, year]: sum(sales_in_dollars)
compute cube sales
```

- Transform it into a SQL-like language (with a new operator **cube by**, introduced by Gray et al.'96)

```
SELECT item, city, year, SUM (amount)
FROM SALES
CUBE BY item, city, year
```

- Need compute the following Group-Bys

```
(date, product, customer),
(date,product),(date, customer), (product, customer),
(date), (product), (customer)
()
```





Cube Computation: ROLAP-Based Method

- Efficient cube computation methods
 - ROLAP-based cubing algorithms (Agarwal et al'96)
 - Array-based cubing algorithm (Zhao et al'97)
 - Bottom-up computation method (Beyer & Ramarkrishnan'99)
 - H-cubing technique (Han, Pei, Dong & Wang:SIGMOD'01)
- ROLAP-based cubing algorithms
 - Sorting, hashing, and grouping operations are applied to the dimension attributes in order to reorder and cluster related tuples
 - Grouping is performed on some sub-aggregates as a “partial grouping step”
 - Aggregates may be computed from previously computed aggregates, rather than from the base fact table

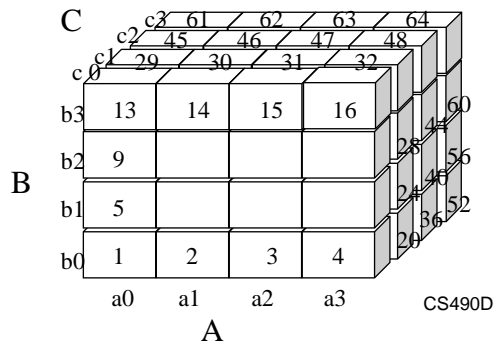
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Multi-way Array Aggregation for Cube Computation

- Partition arrays into chunks (a small subcube which fits in memory).
- Compressed sparse array addressing: (chunk_id, offset)
- Compute aggregates in “multiway” by visiting cube cells in the order which minimizes the # of times to visit each cell, and reduces memory access and storage cost.



What is the best traversing order to do multi-way aggregation?

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Multi-Way Array Aggregation for Cube Computation (Cont.)

- Method: the planes should be sorted and computed according to their size in ascending order.
 - See the details of Example 2.12 (pp. 75-78)
 - Idea: keep the smallest plane in the main memory, fetch and compute only one chunk at a time for the largest plane
- Limitation of the method: computing well only for a small number of dimensions
 - If there are a large number of dimensions, “bottom-up computation” and iceberg cube computation methods can be explored

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Data Warehousing and OLAP Technology for Data Mining

- What is a data warehouse?
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Data Warehouse Usage

- Three kinds of data warehouse applications
 - Information processing
 - supports querying, basic statistical analysis, and reporting using crosstabs, tables, charts and graphs
 - Analytical processing
 - multidimensional analysis of data warehouse data
 - supports basic OLAP operations, slice-dice, drilling, pivoting
 - Data mining
 - knowledge discovery from hidden patterns
 - supports associations, constructing analytical models, performing classification and prediction, and presenting the mining results using visualization tools.
- Differences among the three tasks

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From On-Line Analytical Processing to On Line Analytical Mining (OLAM)

- Why online analytical mining?
 - High quality of data in data warehouses
 - DW contains integrated, consistent, cleaned data
 - Available information processing structure surrounding data warehouses
 - ODBC, OLEDB, Web accessing, service facilities, reporting and OLAP tools
 - OLAP-based exploratory data analysis
 - mining with drilling, dicing, pivoting, etc.
 - On-line selection of data mining functions
 - integration and swapping of multiple mining functions, algorithms, and tasks.

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Discovery-Driven Exploration of Data Cubes

- Hypothesis-driven
 - exploration by user, huge search space
- Discovery-driven (Sarawagi, et al.'98)
 - Effective navigation of large OLAP data cubes
 - pre-compute measures indicating exceptions, guide user in the data analysis, at all levels of aggregation
 - Exception: significantly different from the value anticipated, based on a statistical model
 - Visual cues such as background color are used to reflect the degree of exception of each cell

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Examples: Discovery-Driven Data Cubes

item	all
region	all

Sum of sales	month											
	Jan	Feb	Mar	Apr	May	Jun	Jul	Aug	Sep	Oct	Nov	Dec
Total		1%	-1%	0%	1%	3%	-1%	-9%	-1%	2%	-4%	3%

Avg sales	item	month											
		Jan	Feb	Mar	Apr	May	Jun	Jul	Aug	Sep	Oct	Nov	Dec
	Sony b/w printer	9%	-8%	2%	-5%	14%	-4%	0%	41%	-13%	-15%	-11%	
	Sony color printer	0%	0%	3%	2%	4%	-10%	-13%	0%	4%	-6%	4%	
	HP b/w printer	-2%	1%	2%	3%	8%	0%	-12%	-9%	3%	-3%	6%	
	HP color printer	0%	0%	-2%	1%	0%	-1%	-7%	-2%	1%	-5%	1%	
	IBM home computer	1%	-2%	-2%	1%	3%	3%	-10%	4%	1%	-4%	-1%	
	IBM laptop computer	0%	0%	-1%	-1%	3%	4%	2%	-10%	-2%	0%	-9%	
	Toshiba home computer	-2%	-5%	1%	1%	-1%	1%	5%	-3%	-5%	-1%	-1%	
	Toshiba laptop computer	1%	0%	3%	0%	-2%	-2%	-5%	3%	2%	-1%	0%	
	Logitech mouse	3%	-2%	-1%	0%	4%	6%	-11%	2%	1%	-4%	0%	
	Ergo-way mouse	0%	0%	2%	3%	1%	-2%	-2%	-5%	0%	-5%	8%	

item	IBM home computer
------	-------------------

Avg sales	region	month											
		Jan	Feb	Mar	Apr	May	Jun	Jul	Aug	Sep	Oct	Nov	Dec
	North	-1%	-3%	-1%	0%	3%	4%	-7%	1%	0%	-3%	-3%	
	South	-1%	1%	-9%	6%	-1%	-39%	9%	-34%	4%	1%	7%	
	East	-1%	-2%	2%	-3%	1%	18%	-2%	11%	-3%	-2%	-1%	
	West	4%	0%	-1%	-3%	5%	1%	-18%	8%	5%	-8%	1%	



Summary

- **Data warehouse**
- A **multi-dimensional model** of a data warehouse
 - Star schema, snowflake schema, fact constellations
 - A data cube consists of dimensions & measures
- **OLAP operations**: drilling, rolling, slicing, dicing and pivoting
- **OLAP servers**: ROLAP, MOLAP, HOLAP
- **Efficient computation of data cubes**
 - Partial vs. full vs. no materialization
 - Multiway array aggregation
 - Bitmap index and join index implementations
- **Further development of data cube technology**
 - Discovery-drive and multi-feature cubes
 - From OLAP to OLAM (on-line analytical mining)

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